

# Structural Programming and Data Structures

Winter 2000

## CMPUT 102: Sorting

Dr. Osmar R. Zaiane



University of Alberta

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## Course Content

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## Objectives of Lecture 23

### Sorting

- Introduce the problem of sorting collections;
- Learn how to sort using a bubble sort algorithm;
- Learn how to sort with the selection algorithm.

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## Outline of Lecture 23



- The sorting problem
- Simple methods like bubble sort
- Selection sort example
- Selection sort code
- Complexity of selection sort

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## The Sort Problem

- Given a container, with elements that can be compared, put it in increasing or decreasing order.

|    |    |    |    |    |    |    |    |    |    |
|----|----|----|----|----|----|----|----|----|----|
| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  |
| 25 | 50 | 10 | 95 | 75 | 30 | 70 | 55 | 60 | 80 |



|    |    |    |    |    |    |    |    |    |    |
|----|----|----|----|----|----|----|----|----|----|
| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  |
| 10 | 25 | 30 | 50 | 55 | 60 | 70 | 75 | 80 | 95 |

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## Sorting Problem (con't)

- Given a container of  $n$  elements  $A[0..n-1]$  such that any elements  $x$  and  $y$  in the container  $A$  can be compared directly, either  $x < y$ , or  $x = y$ , or  $x > y$ .
- We want to permute the elements of  $A$  so that at the end  $A[0] \leq A[1] \leq \dots \leq A[n-1]$  (monotone non-decreasing), or  $A[0] \geq A[1] \geq \dots \geq A[n-1]$  (monotone decreasing)

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## The Order of Things

- Numbers
  - $-99 < -34 < -6 < 0 < 1 < 9 < 23 < 999$
- Characters
  - $A < B < C < D < E < F < \dots < X < Y < Z$
  - $a < b < c < d < e < f < \dots < x < y < z$
  - $a < z < A < Z$
- Strings
  - "Abacus" < "Alpha" < "Hello" < "Memorization" < "Memorize" < "Memory" < "Zebra"

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## Sorting

- There is often a need to put data in order.
- Sorting is among the most basic and universal of computational problems.
- There are hundreds of algorithms and variations on algorithms.
- Variety of sorting methods: internal vs. external, sorting in place vs. sorting with auxiliary structures, etc.

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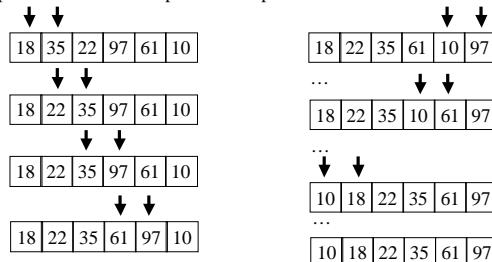
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### One simple sorting method

Given a list:

35 18 22 97 61 10

Iterate over the collection and permute neighbours if necessary repeat iteration until no permutation possible.



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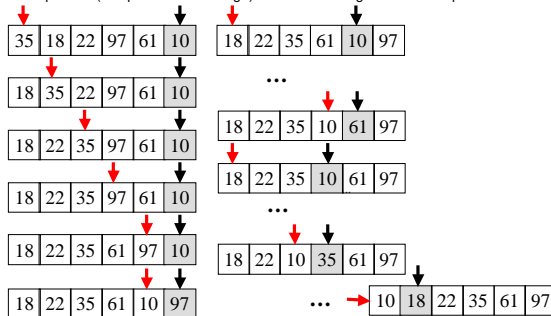
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### The Bubble Sort

Given a list: 35 18 22 97 61 10

First pass of (compare and exchange) will send the largest to the last position.



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```
public static void bubble_sort( int data[] ) {
    // Sort the given Array with selection sort method
    //(Ascending order)
```

```
    int current, last;
```

```
    for ( last = data.length-1; last >= 1; last--)
        for ( current = 0; current < last; current++)
            if ( data[current] > data[current+1] )
                this.exchange( data, current, current+1 );
}
```

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## Outline of Lecture 23



- The sorting problem
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## Selection Sort

- Look for the smallest element and exchange it with the element whose index is 0.

| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  |
|----|----|----|----|----|----|----|----|----|----|
| 25 | 50 | 10 | 95 | 75 | 30 | 70 | 55 | 60 | 80 |

| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  |
|----|----|----|----|----|----|----|----|----|----|
| 10 | 50 | 25 | 95 | 75 | 30 | 70 | 55 | 60 | 80 |

## Selection Sort (con't)

- Look for the smallest element whose index is greater than or equal to 1 and exchange it with the element whose index is 1.

| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  |
|----|----|----|----|----|----|----|----|----|----|
| 10 | 50 | 25 | 95 | 75 | 30 | 70 | 55 | 60 | 80 |

| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  |
|----|----|----|----|----|----|----|----|----|----|
| 10 | 25 | 50 | 95 | 75 | 30 | 70 | 55 | 60 | 80 |

## Selection Sort (con't)

- Look for the smallest element whose index is greater than or equal to 2 and exchange it with the element whose index is 2.

| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  |
|----|----|----|----|----|----|----|----|----|----|
| 10 | 25 | 50 | 95 | 75 | 30 | 70 | 55 | 60 | 80 |

| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  |
|----|----|----|----|----|----|----|----|----|----|
| 10 | 25 | 30 | 95 | 75 | 50 | 70 | 55 | 60 | 80 |

## Selection Sort (con't)

- Look for the smallest element whose index is greater than or equal to  $k$  and exchange it with the element whose index is  $k$  (for  $k = 3, 4, \dots, n-1$ )

| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  |
|----|----|----|----|----|----|----|----|----|----|
| 10 | 25 | 30 | 95 | 75 | 50 | 70 | 55 | 60 | 80 |

| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  |
|----|----|----|----|----|----|----|----|----|----|
| 10 | 25 | 30 | 50 | 75 | 95 | 70 | 55 | 60 | 80 |

## Selection Sort (con't)

| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  |
|----|----|----|----|----|----|----|----|----|----|
| 10 | 25 | 30 | 50 | 75 | 95 | 70 | 55 | 60 | 80 |

| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  |
|----|----|----|----|----|----|----|----|----|----|
| 10 | 25 | 30 | 50 | 55 | 95 | 70 | 75 | 60 | 80 |

| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  |
|----|----|----|----|----|----|----|----|----|----|
| 10 | 25 | 30 | 50 | 55 | 60 | 70 | 75 | 95 | 80 |

## Selection Sort (con't)

|    |    |    |    |    |    |    |    |    |    |
|----|----|----|----|----|----|----|----|----|----|
| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  |
| 10 | 25 | 30 | 50 | 55 | 60 | 70 | 75 | 95 | 80 |

↓

|    |    |    |    |    |    |    |    |    |    |
|----|----|----|----|----|----|----|----|----|----|
| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  |
| 10 | 25 | 30 | 50 | 55 | 60 | 70 | 75 | 95 | 80 |

↓ ↓

|    |    |    |    |    |    |    |    |    |    |
|----|----|----|----|----|----|----|----|----|----|
| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  |
| 10 | 25 | 30 | 50 | 55 | 60 | 70 | 75 | 80 | 95 |

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## Selection Sort Algorithm

**INPUT :** data: an array of int  
**OUTPUT:** data: sorted in ascending order

**Method:**

```
for ( first = 1; first < length - 1; first ++ ) {  
    find Smallest such that data[Smallest] is the  
    smallest between data[first] and data[length-1];  
    permute Data[first] and Data[Smallest];  
}
```

## Selection Sort Code

```
private void selectionSort(int anArray[] ) {  
    // Sort the given Array with selection sort method (Ascending order)  
  
    int index;  
    int smallestIndex;  
  
    for (index = 0; index < anArray.length - 1; index++) {  
        smallestIndex = this.getSmallest(anArray, index);  
        this.exchange(anArray, index, smallestIndex);  
    }  
}
```

## Code for method: exchange

```
private void exchange(int anArray[], int i, int j) {  
    // Exchange the elements of the array with  
    // the given two indexes.  
  
    int temp;  
  
    temp = anArray[i];  
    anArray[i] = anArray[j];  
    anArray[j] = temp;  
}
```

## Code for method: getSmallest

```
private int getSmallest(int anArray[], int start) {  
    // Return the index of the smallest element  
    // of the given array whose index is greater  
    // than or equal to the given start index.  
    int smallestIndex;  
    int index;  
  
    smallestIndex = start;  
    for (index = start + 1; index < anArray.length; index++)  
        if (anArray[index] < anArray[smallestIndex])  
            smallestIndex = index;  
    return smallestIndex;  
}
```

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## Complexity of Selection Sort

- How many comparison operations are required for a selection sort of an  $n$ -element container?
- The sort method executes **getSmallest** for the indexes:  $0, 1, \dots, n-2$ .
- Each time **getSmallest** is executed for an index, it does:  $(n - \text{index})$  comparisons.
- The total number of comparisons is:

$$(n-0) + (n-1) + \dots + (n-(n-2)) = (1 + 2 + \dots + n) - 1 = \frac{n(n+1)}{2} - 1 \approx n^2 \text{ for large } n.$$

$O(n^2)$  → Quadratic time complexity