"STRIPS" Planning

• Set of *operators*, where each operator has

- Set of *parameters*
- Set of *preconditions*

 Set of *effects*, consisting of add effects and delete effects.

- Set of *objects* to instantiate operator's parameters fully instantiated operator = *action*
- Set of propositions representing *initial state*

1

• Set of propositions representing *goals*

Planning problem: Find sequence of actions that, starting in initial state, achieve all the goals

Approaches to STRIPS planning

- Search through space of *world states*
 - forward search,
 - regression search
 - bi-directional search
 - means-ends analysis
 - . . .
- Search through space of *plans*
 - total order planning
 - partial order planning
- Search through *planning graph*

 2

GraphPlan Approach

 Construct a "PlanGraph" that contains all valid plans

 + other stuff (invalid plans)
 up to a maximum depth

3

2. Search PlanGraph for valid plan ... then return that plan

Simple Cake-Eating Domain

- Initial: HaveCake $\land \neg$ EatenCake
- Goal: HaveCake \land EatenCake



• PlanGraph



 4

 $\operatorname{Graph-Plan}$

Parts of a PlanGraph



"2-leveled" Graph $\langle \mathcal{S}_0, \mathcal{A}_0, \mathcal{S}_1, \mathcal{A}_1, \ldots \rangle$

- S_0 : propositions in initial state
- A_i : each action whose preconditions all occur in level S_{i-1}
- *S_i*: each prop'n that is ADDed/DELETEd by
 ★ an action in level *A_i* ★ a "No-Op" (persistance)

5

- Mutex links
 - \star between actions in level A_i
 - \star between propositions in level S_i "mutually exclusive"
 - "cannot occur in same plan"



Mutex Conditions#2: Propositions



Between 2 propositions ρ_1 and ρ_2 , same level S_i :

- Negation $\rho_1 = \neg \rho_2$
- Inconsistent Support
 Every action achieving ρ₁ (from S_{i-1}) is mutex with every action achieving ρ₂
 In S₁: HaveCake mutex EatenCake as only way to achieve HaveCake:
 NoOp(HaveCake)
 is mutex with only way to achieve EatenCake:
 Eat
 N.b.: Not mutex at S₂ !
 Note:
 Note:

Planning Graphs

- A valid plan is "2-leveled" graph
 two kinds of nodes
 - (propositions, actions) alternates: proposition level, action level
 - 5 kinds of edges
 - $\begin{array}{ll} \star \text{ precondition} & (S_i \to A_i) \\ \star \text{ add effect} & (A_i \to S_{i+1}) \\ \star \text{ delete effect} & (A_i \to S_{i+1}) \\ \star \text{ mutex-action} & (A_i \leftrightarrow A_i) \\ \star \text{ mutex-prop} & (S_i \leftrightarrow S_i) \end{array}$
 - Include action O at action-level A_i if all preconditions at proposition-level S_i
 - Include proposition ρ at proposition-level S_i if it is add/delete effect of action $O \in A_{i-1}$ (including *no-op* actions)

Restriction:

Allow actions O_1 , O_2 at same time tONLY if don't interfere with each other

• $PlanningGraph \approx valid plan$ but without no-interfere restriction

GraphPlan Algorithm

function Graphplan(problem) returns solution or failure
graph ← Initial-Planning-Graph(problem)
goals ← Goals[problem]
loop do
if goals all non-mutex in last level of graph then do
solution ← Extract-Solution(graph, goals, Length(graph))
if solution ≠ failure then return solution
else if No-Solution-Possible(graph) then return failure
graph ← Expand-Graph(graph, problem)

9

end

Flat-Tire Domain

FI= Flat; Sp= Spare; Ax= Axel; Tr= Trunk; Gr= Ground

- Initial: At(Fl, Ax) \land At(Sp, Tr)
- Goal: At(Sp, Ax)
- Actions:

$$\begin{array}{c} \mathsf{Op} \left(\begin{array}{ccc} \mathsf{TakeOutSpare} \\ \mathsf{PreC:} & \mathsf{At}(\mathsf{Sp}, \mathsf{Tr}) \\ \mathsf{Eff:} & \neg \mathsf{At}(\mathsf{Sp}, \mathsf{Tr}) \land \mathsf{At}(\mathsf{Sp}, \mathsf{Gr}) \end{array} \right) \\ \\ \mathsf{Op} \left(\begin{array}{ccc} \mathsf{RemoveFlat} \\ \mathsf{PreC:} & \mathsf{At}(\mathsf{Fl}, \mathsf{Ax}) \\ \mathsf{Eff:} & \neg \mathsf{At}(\mathsf{Fl}, \mathsf{Ax}) \land \mathsf{At}(\mathsf{Fl}, \mathsf{Gr}) \end{array} \right) \\ \\ \mathsf{Op} \left(\begin{array}{ccc} \mathsf{PutOnSpare} \\ \mathsf{PreC:} & \mathsf{At}(\mathsf{Sp}, \mathsf{Gr}) \land \neg \mathsf{At}(\mathsf{Fl}, \mathsf{Ax}) \\ \mathsf{Eff:} & \neg \mathsf{At}(\mathsf{Sp}, \mathsf{Gr}) \land \mathsf{At}(\mathsf{Sp}, \mathsf{Ax}) \end{array} \right) \\ \\ \mathsf{Op} \left(\begin{array}{ccc} \mathsf{LeaveOverNight} \\ \mathsf{PreC:} & \{\} \\ \mathsf{Eff:} & \neg \mathsf{At}(\mathsf{Sp}, \mathsf{Gr}) \land \neg \mathsf{At}(\mathsf{Sp}, \mathsf{Ax}) \land \neg \mathsf{At}(\mathsf{Sp}, \mathsf{Tr}) \\ & \land \neg \mathsf{At}(\mathsf{Fl}, \mathsf{Gr}) \land \neg \mathsf{At}(\mathsf{Fl}, \mathsf{Ax}) \end{array} \right) \end{array} \right) \\ \end{array}$$



Trace of GraphPlan Algorithm #1

- S_0 : initial facts (include \neg facts)
- As $At(Sp, Ax) \notin S_0$ do not call Extract-Solution
- Expand-Graph forms A₀ with

 * 3 "real" actions
 * 5 no-op actions;

 S₁ is effects

 Expand-Graph then finds

 * 4 action-mutex within A₀
 * 4 prop-mutex within S₁

 As At(Sp,Ax) ∉ S₁ do not call Extract-Solution
- Expand-Graph forms A₁ with

 * 4 "real" actions
 * 7 no-op actions

 S₂ is effects

Mutex wrt FlatTire



Trace of GraphPlan Algorithm #2

- "All" goal literals, At(Sp, Ax), in S_2 none are mutex ...
- So there MAY be solution
 - ... call Extract-Solution

```
Extract-Solution(...)

Let G_n be the GOAL at last level, S_n

For each i = n..1

* Let H_i be a conflict-free subset of A_{i-1},

that covers G_i (in S_i)

* Let G_{i-1} be preconditions of H_i

... until reach state in S_0 satisfying all goals

Action-set H is "conflict-free"

\equiv

no pair of H are mutex, and

no pair of preconditions (in G) are mutex
```

Trace of Extract-Solution

•
$$G_2 = \{ \operatorname{At}(\operatorname{Sp}, \operatorname{Ax}) \}$$

 $H_2 = \{ \operatorname{PutOnSpare} \}$

 \Rightarrow $H_1 = \{$ TakeOutSpare, RemoveFlat $\}$

•
$$G_0 = \{ \operatorname{At}(\operatorname{Sp},\operatorname{Tr}), \operatorname{At}(\operatorname{Fl},\operatorname{Ax}) \}$$

As in $G_0 \subset S_0$, DONE!

Extending PlanGraph

Add action level A_i : **ForEach** action^(*) OIf O's preconditions all true in prop-level S_i , and NOT mut-ex, **Then** add O to level A_i include precondition-links create mutex (O:actions-I-am-exclusive-of) Add prop-level S_{i+1} : **ForEach** effect ρ of each action in action-level A_i Add ρ to prop-level S_{i+1} Add $S \leftarrow \rho$ add- or delete- links Mark ρ_1, ρ_2 as mutex if each way of generating ρ_1 is mutex to each way of generating ρ_2

 $^{(\ast)}$ each instantiation of each operator; including ''no-op''s

Correctness

Graphplan is sound and complete:

- * any plan Graphplan finds is a legal plan
- * if \exists legal plan then Graphplan will find one.

Theorem:	If \exists valid plan using $\leq t$ time steps,
	then plan is subgraph of (depth- t)
	Planning Graph.

+

If Goals not satisfiable by any valid plan, then GraphPlan will halt, w/failure, in finite time.

(extends most partial-order planners)

Leveling Off

- GraphPlan \approx Iterative deepening When to stop??
- Lemma: If no valid plan exists, then \exists a prop-level S_n s.t. all future proposition levels are identical to S_n

– Identical \equiv same propositions, mutual exclusions

- graph has "leveled off after S_n "
- Corollary: No solution exists if
 - a goal does not appear in S_n or
 - $-S_n$ has mutually exclusive goals

• Subtlety:

 $\{ on(A,B), on(B,C), on(C,A) \}$

Termination Condition

- Let S_i^t denote set of memoized goal sets at level i after an unsuccessful stage t
- Theorem: If the graph has leveled off at level n and stage t has passed in which $|S_n^{t-1}| = |S_n^t|$, then no valid plan exists

Termination Proof

- As PlanGraph gets deeper...
 - Literals increase monotonically
 - Actions increase monotonically
 - Mutex decrease monotonically
 - * If O_1 and O_2 are mutex in A_k , then mutex in A_i i = 1..k provided $O_1, O_2 \in A_i$
 - * If ρ_1 and ρ_2 are mutex in S_k , then mutex in S_i i = 1..k provided $\rho_1, \rho_2 \in S_i$
- Only finite # of actions/literals, planning graph must eventually "level off"

Experimental Results





"Link-Repeat Problem"

Graph-Plan

Accounting for Graphplan's Efficiency

Mutual exclusions

 (Most constraints are pair-wise mut-ex's;
 Propagating constraints prunes large part of space.)

Consideration of parallel plans
 (Valid parallel plans are short, wrt total plan
 ⇒ reduces cost of constructing pgraph, search)

Memoizing

 (Many goal-sets appear > 1)

 Low-level costs (Graphplan avoid cost of instantiation during search)

Efficiency

Size of Planning Graph

Theorem: Consider planning problem with

- $n \, \operatorname{objects}$,
- \boldsymbol{p} propositions in initial state,
- \boldsymbol{m} operators,

each w/constant number of parameters
Let *l* be length of longest add list.
Then size of a *t*-level planning graph, and time needed to create the graph, are polynomial in *n*, *m*, *p*, *l*, and *t*.

• Empirically: exclusion relations most expensive part of graph creation

Graph creation only significant in simple problems

⇒ As graph is small, "finding mut-ex" is hard as planning... PSpace-hard

Comments

• PlanGraph \neq StateGraph plan \equiv path in StateGraph but plan \equiv flow in PlanGraph

• Like "Traditional TotalOrder Planner": considers action at *FIXED* time

Like "Partial Order Planner" generates partially-ordered plans

• Parallel Plan: can execute many actions at once

if no conflicts

(eg, load all items at once)

- Guaranteed to find **SHORTEST** plan
- \approx *Not* sensitive to given order of goals

Final Comments

• Planning \equiv Searching

 \Rightarrow GraphPlan

...a new approach to Planning

Future work

 Learning (from one plan to next)
 Two-way search (fact→goal, goal→fact)
 beyond "Strips"-like domains
 creating objects, ∀, ...
 incorporating other types of constraints
 Why guarantee SHORTEST path?

• http://www.cs.cmu.edu/~avrim/graphplan.html

SatPlan

Convert plan-situation

 (Operators, Initial/Final Conditions, ...)
 to SAT

(Up to fixed length)

• Run WalkSat to find satisficing assignment \equiv plan...

... iterative deepening

 Plays to SatPlan's strength, as ∃ satisfying assignment...