



PREDICTING THE EFFECTIVENESS OF BIDIRECTIONAL HEURISTIC SEARCH

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Jingwei Chen, University of Alberta



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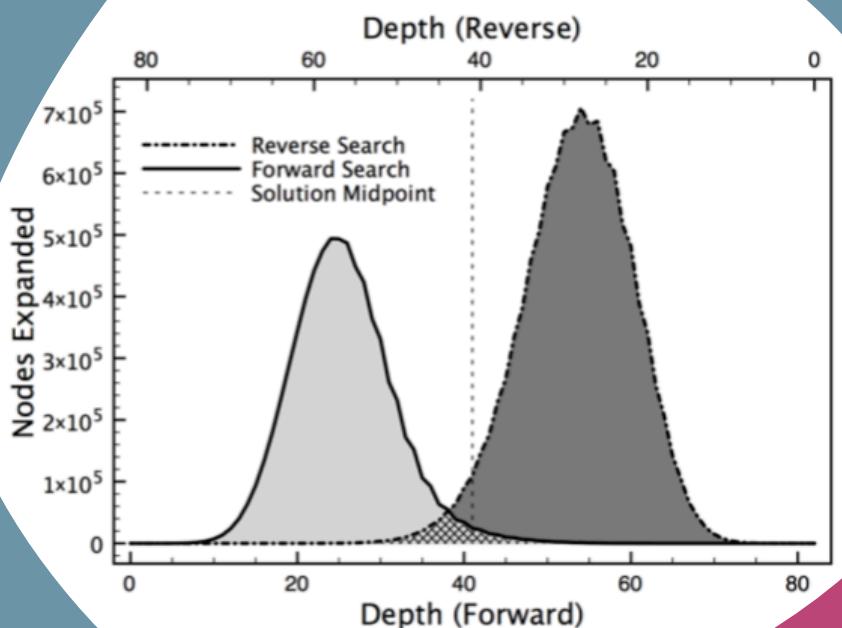
BIDIRECTIONAL SEARCH

Q: When does bidirectional (heuristic) search perform well?

A: Performance of bidirectional search is positively correlated with the number of states that have heuristics that are both low and inaccurate.

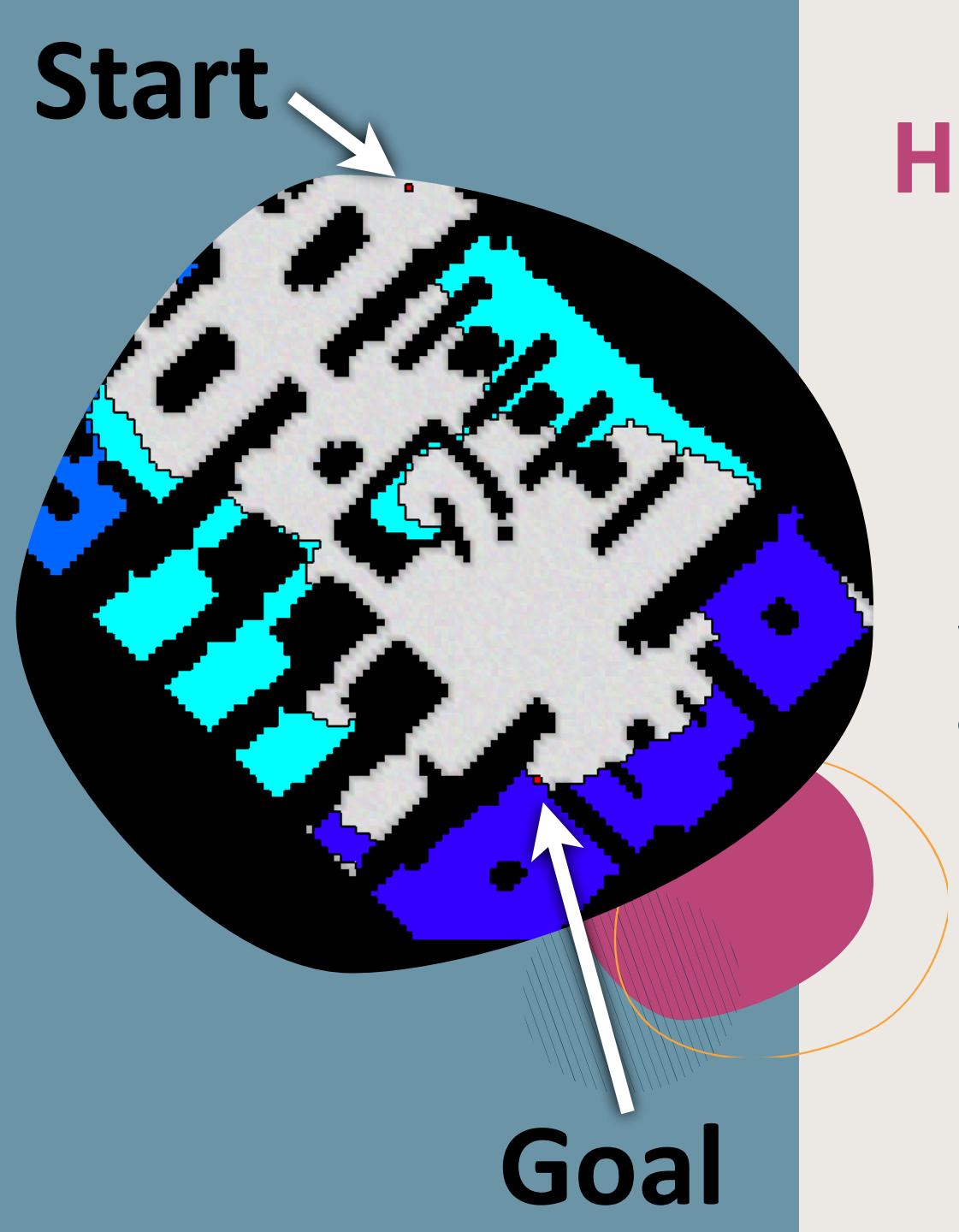


BARKER AND KORF (2015)



A strong heuristic expands a majority of the states in the first half of the search.

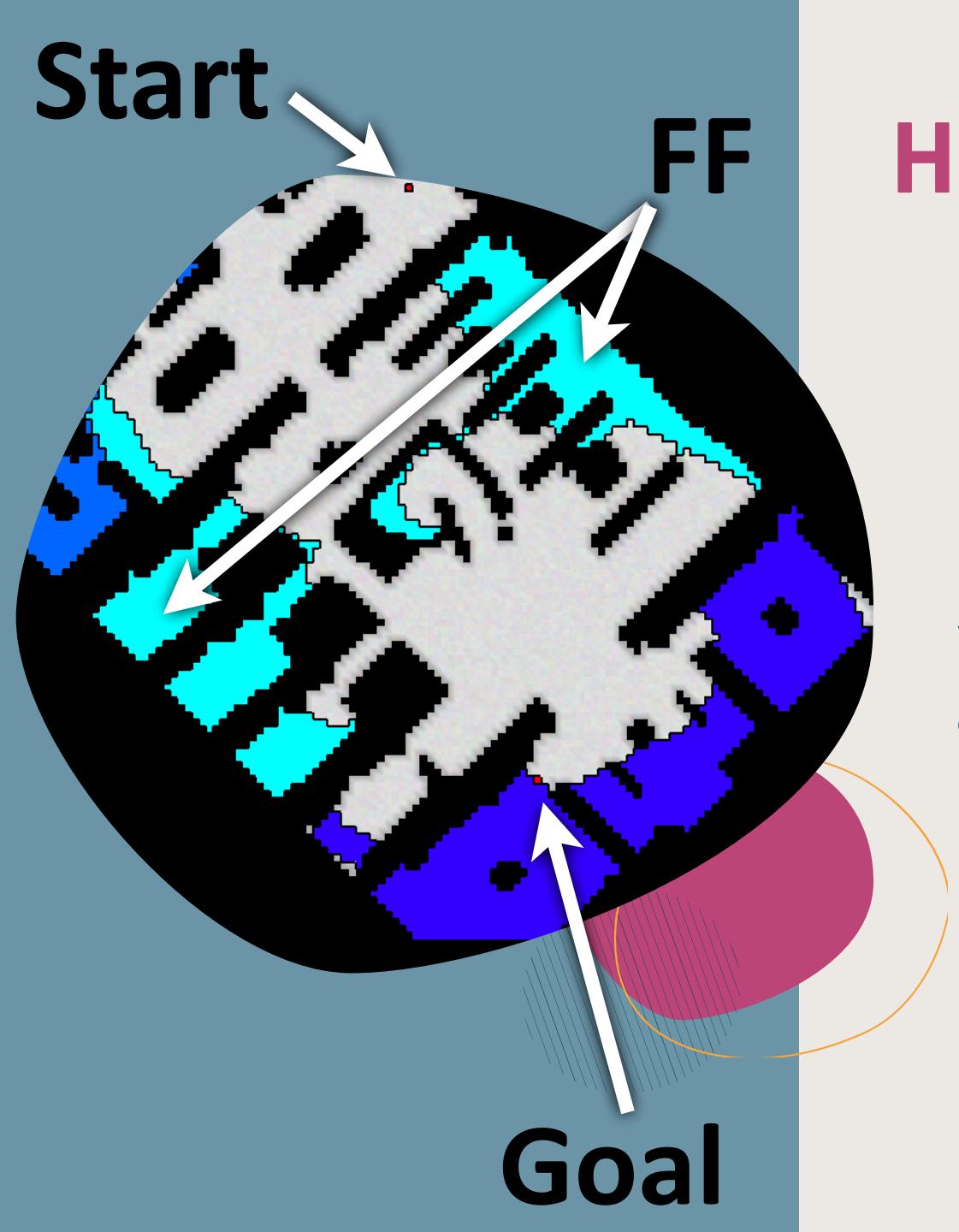
BK2: Unidirectional search outperforms bidirectional search with a strong heuristic.



HOLTE ET AL (2017)

MM is guaranteed to meet in the middle.

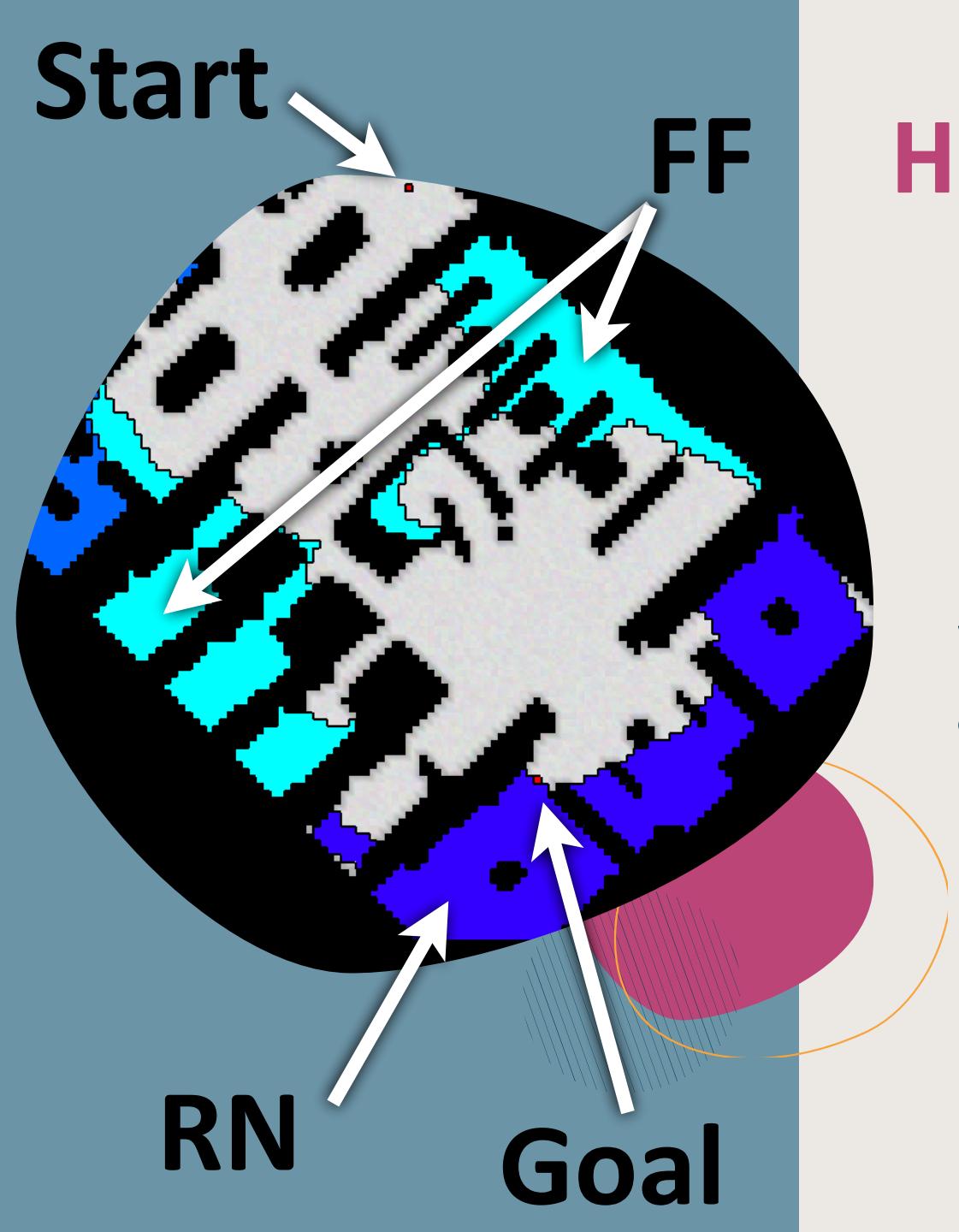
If $|FF| > |RN|$, A* will expand more states than MM if the heuristic is weak, fewer if the heuristic is accurate.



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ECKERLE ET AL (2017)

Front-to-end bidirectional search

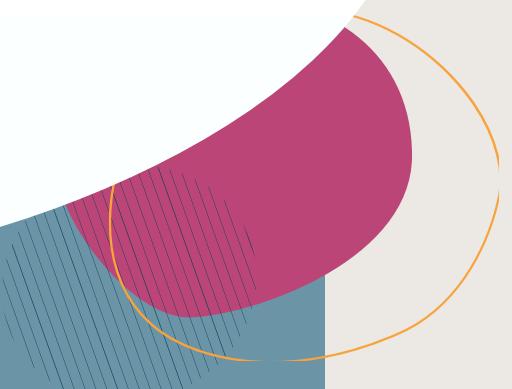
- admissible heuristic
- not necessarily consistent

Corresponds to finding a minimum vertex cover on a bipartite graph

Theorem 6. Let $I = (G, h_F, h_B) \in I_{CON}$ have an optimal solution cost of C^* . If U is an optimal forward path and V is an optimal backward path such that $U_0 = \text{start}$, $V_0 = \text{goal}$, and:

- (1) $f_F(U) < C^*$
- (2) $f_B(V) < C^*$
- (3) $c(U) + c(V) < C^*$

then, in solving problem instance I , any admissible DXBB bidirectional front-to-end search algorithm must expand $(\text{end}(U), \text{end}(V))$.





ECKERLE ET AL (2017)

Necessary expansions for a pair of states:

$$f_f(a) < C^*$$

$$f_b(b) < C^*$$

$$g_f(a) + g_b(b) < C^*$$

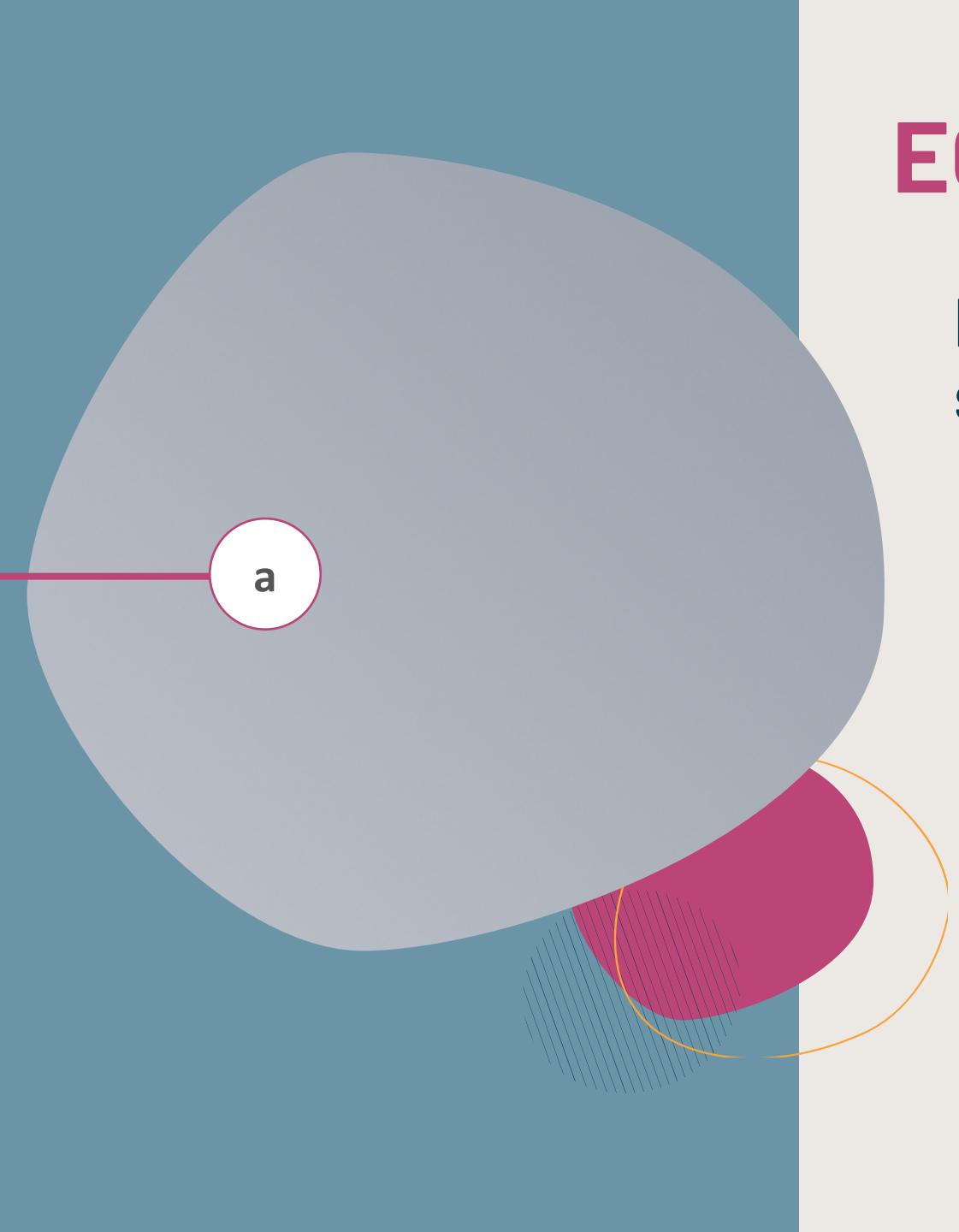
ECKERLE ET AL (2017)

Necessary expansions for a pair of states:

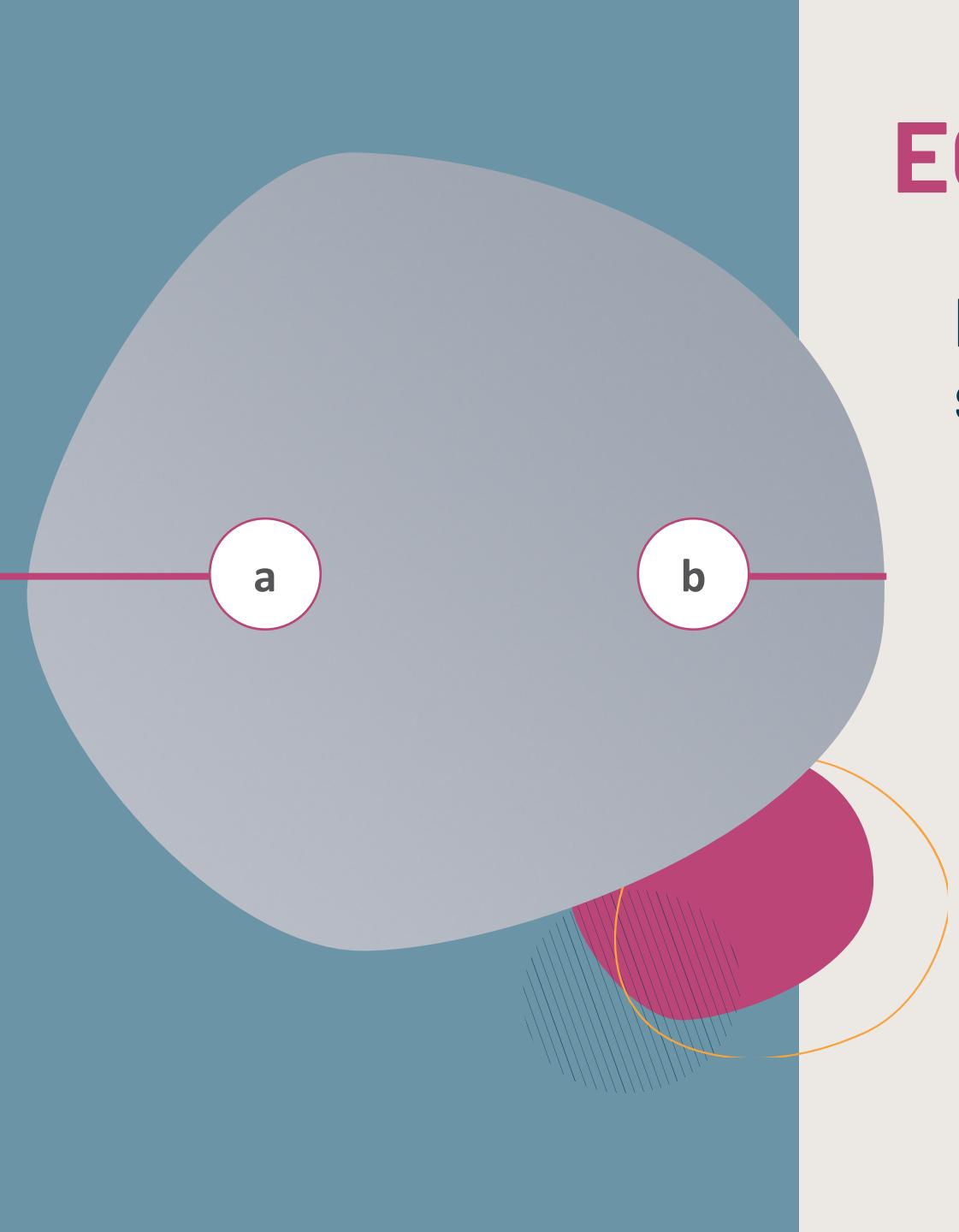
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ECKERLE ET AL (2017)



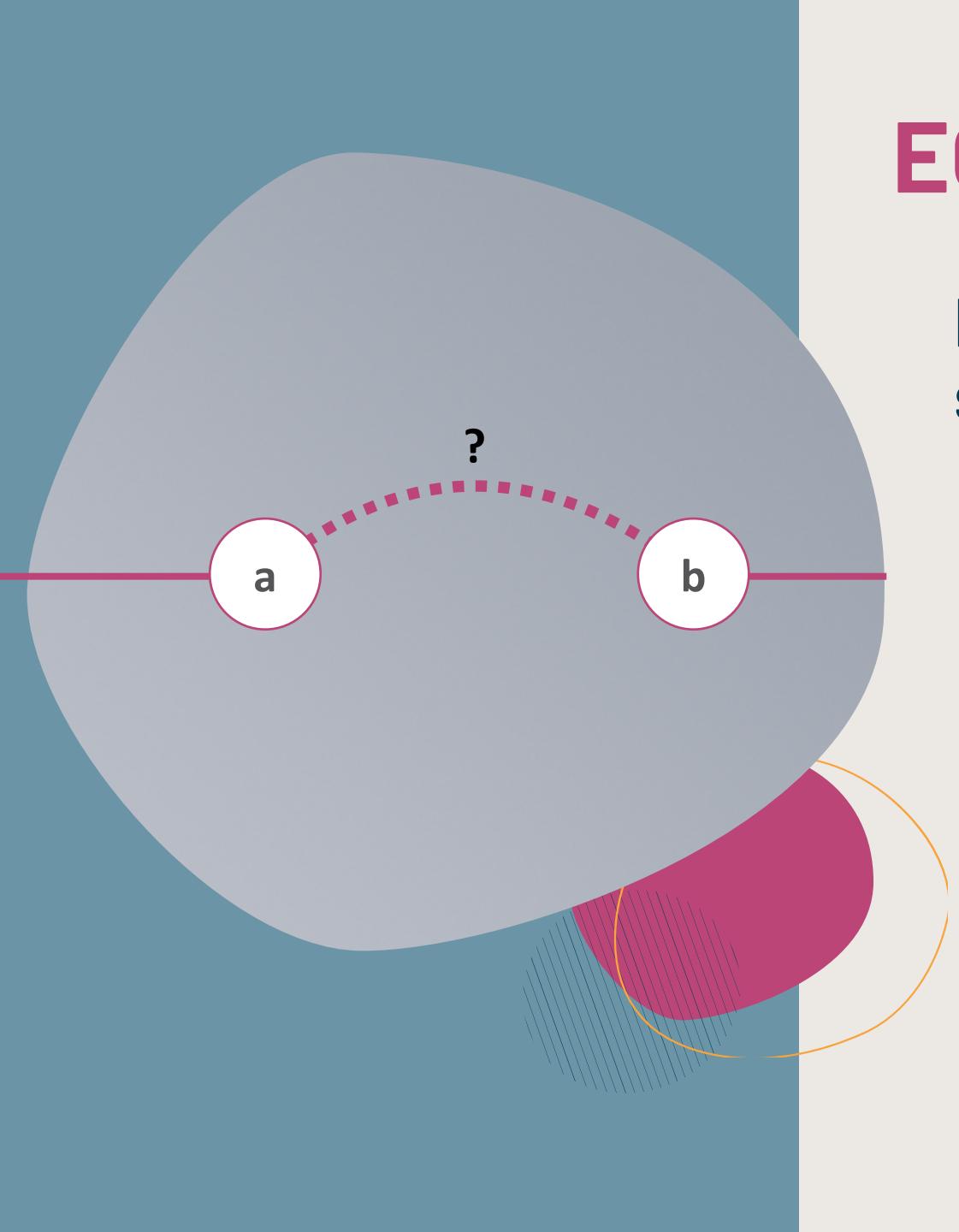
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ECKERLE ET AL (2017)



Necessary expansions for a pair of states:

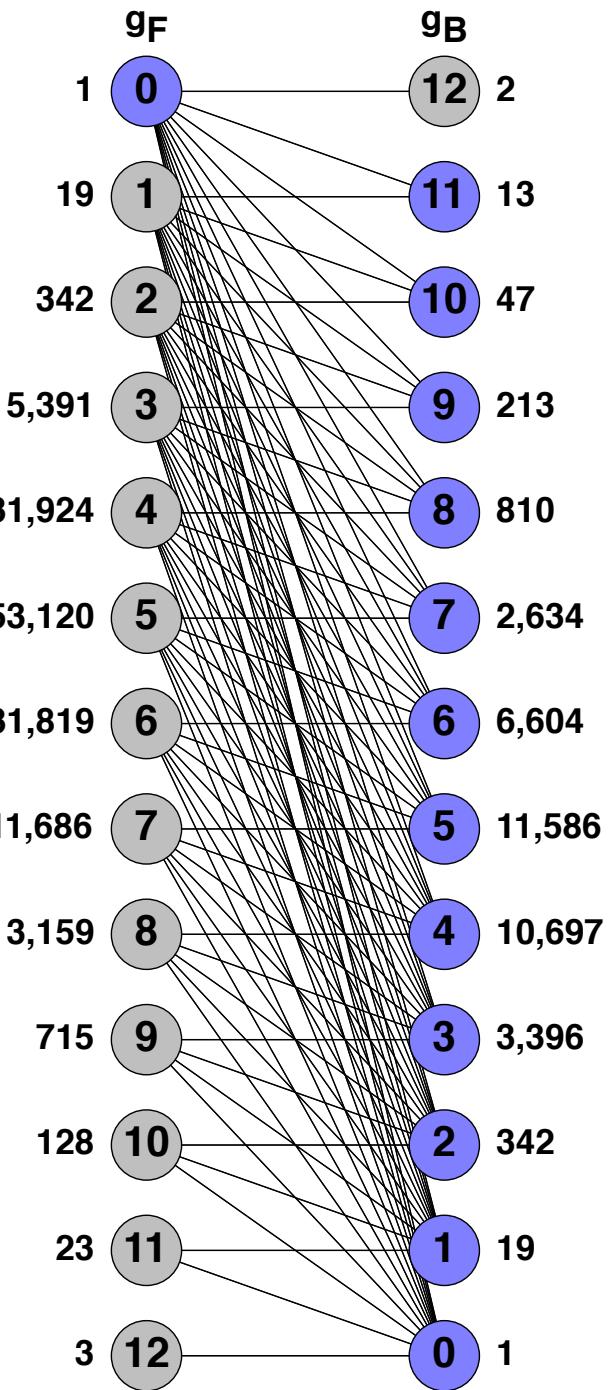
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EXPLANATION

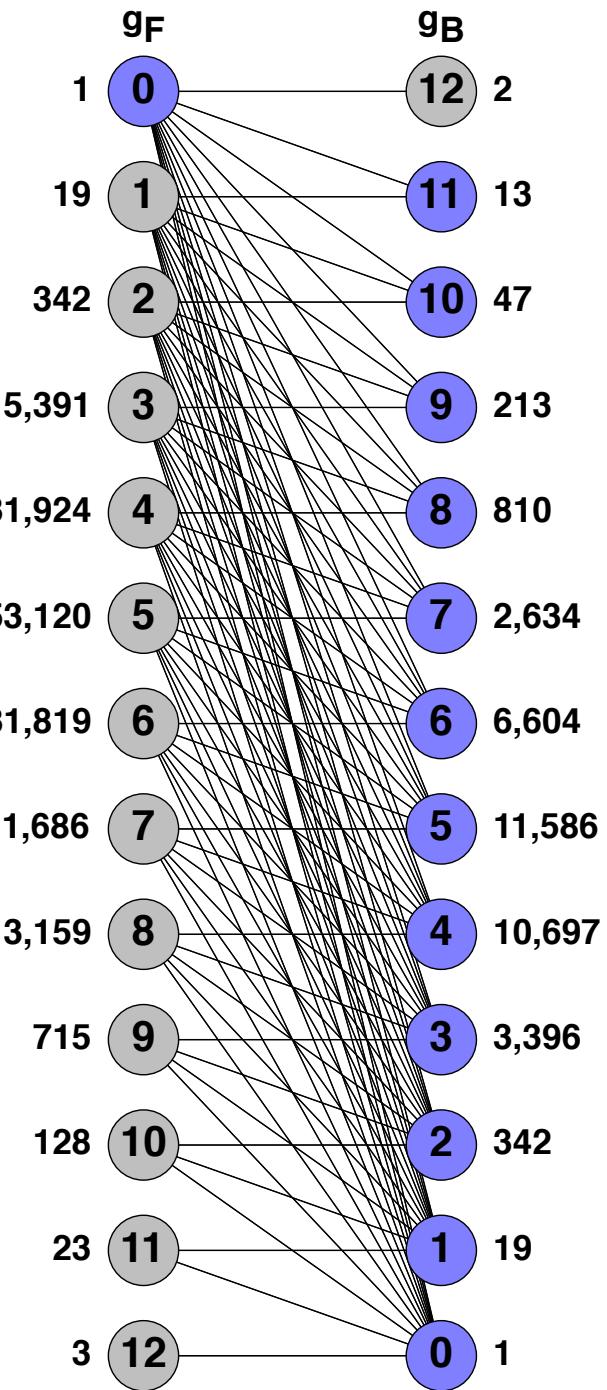
$C^* = 13$



EXPLANATION

$f_f < C^*$

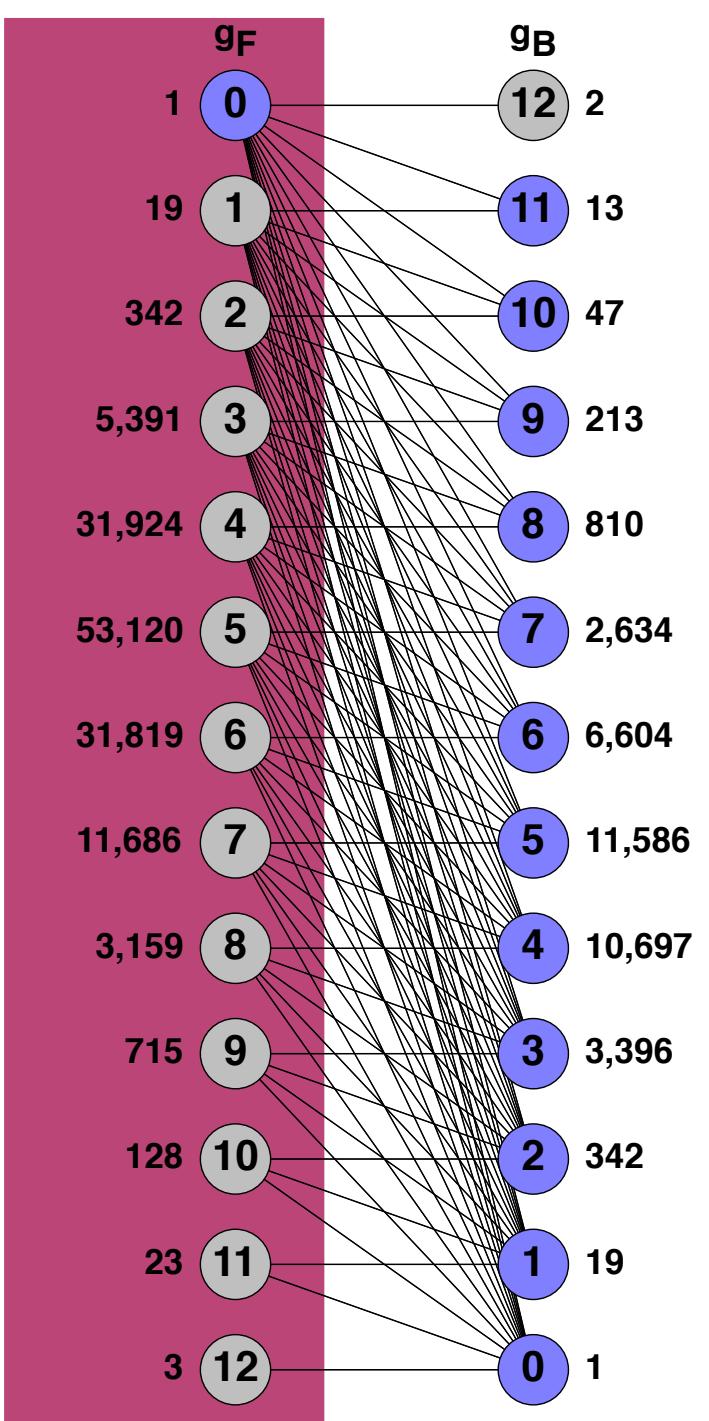
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EXPLANATION

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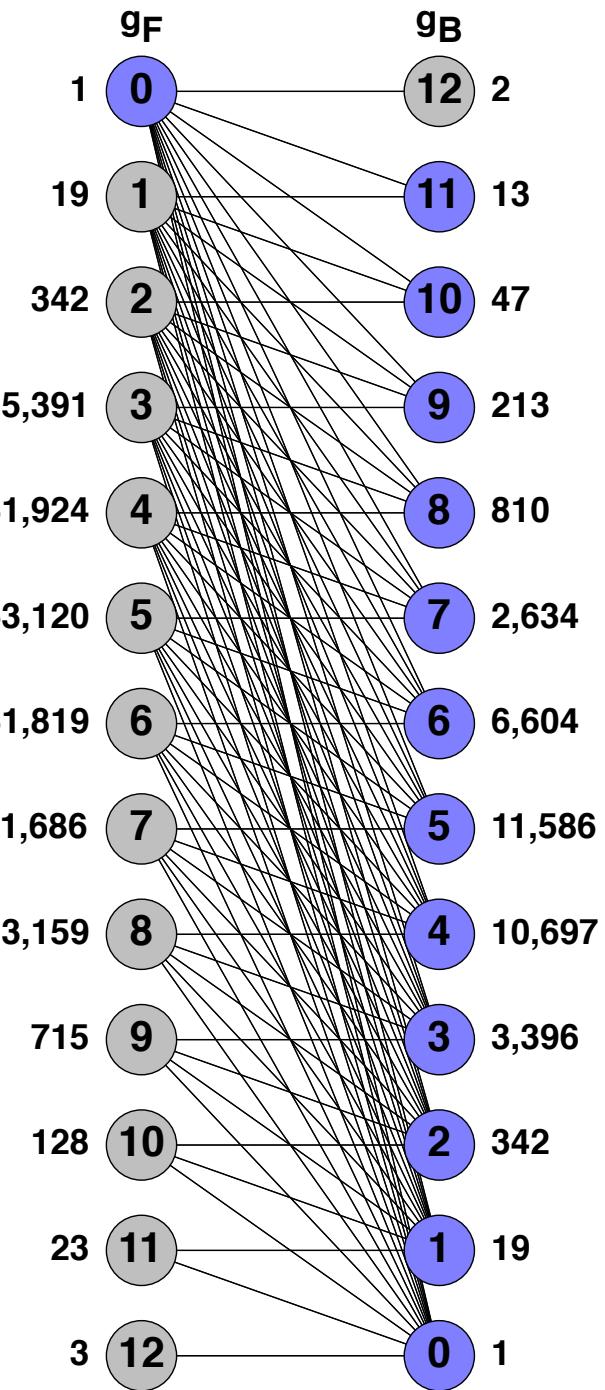
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EXPLANATION

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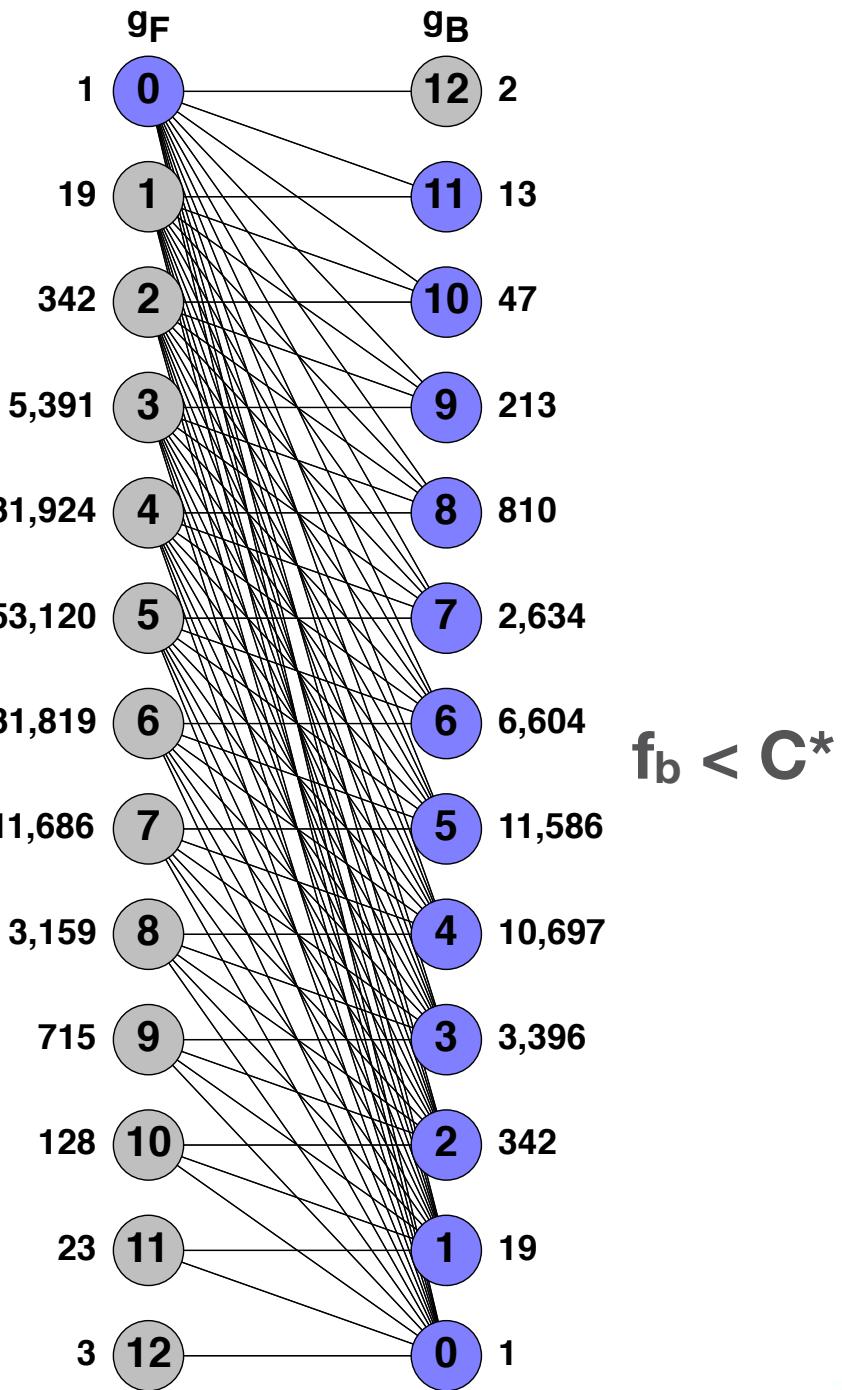
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EXPLANATION

$f_f < C^*$

$C^* = 13$

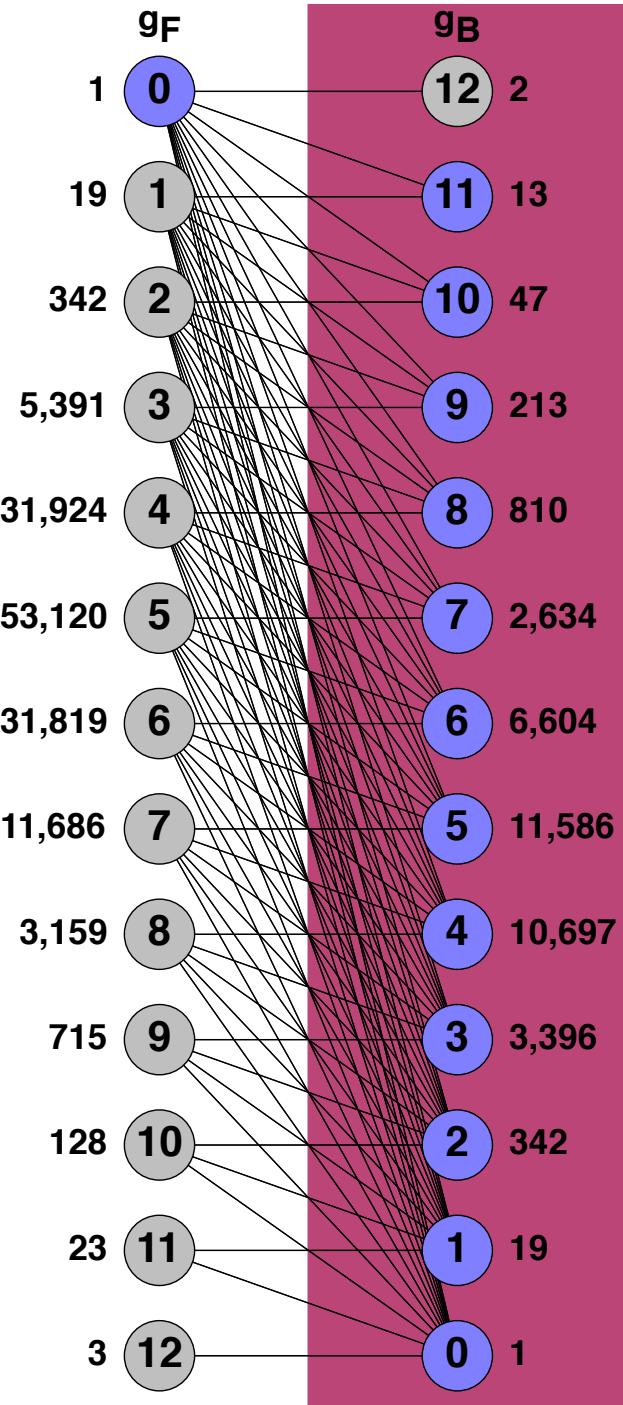


$f_b < C^*$

EXPLANATION

$f_f < C^*$

$C^* = 13$

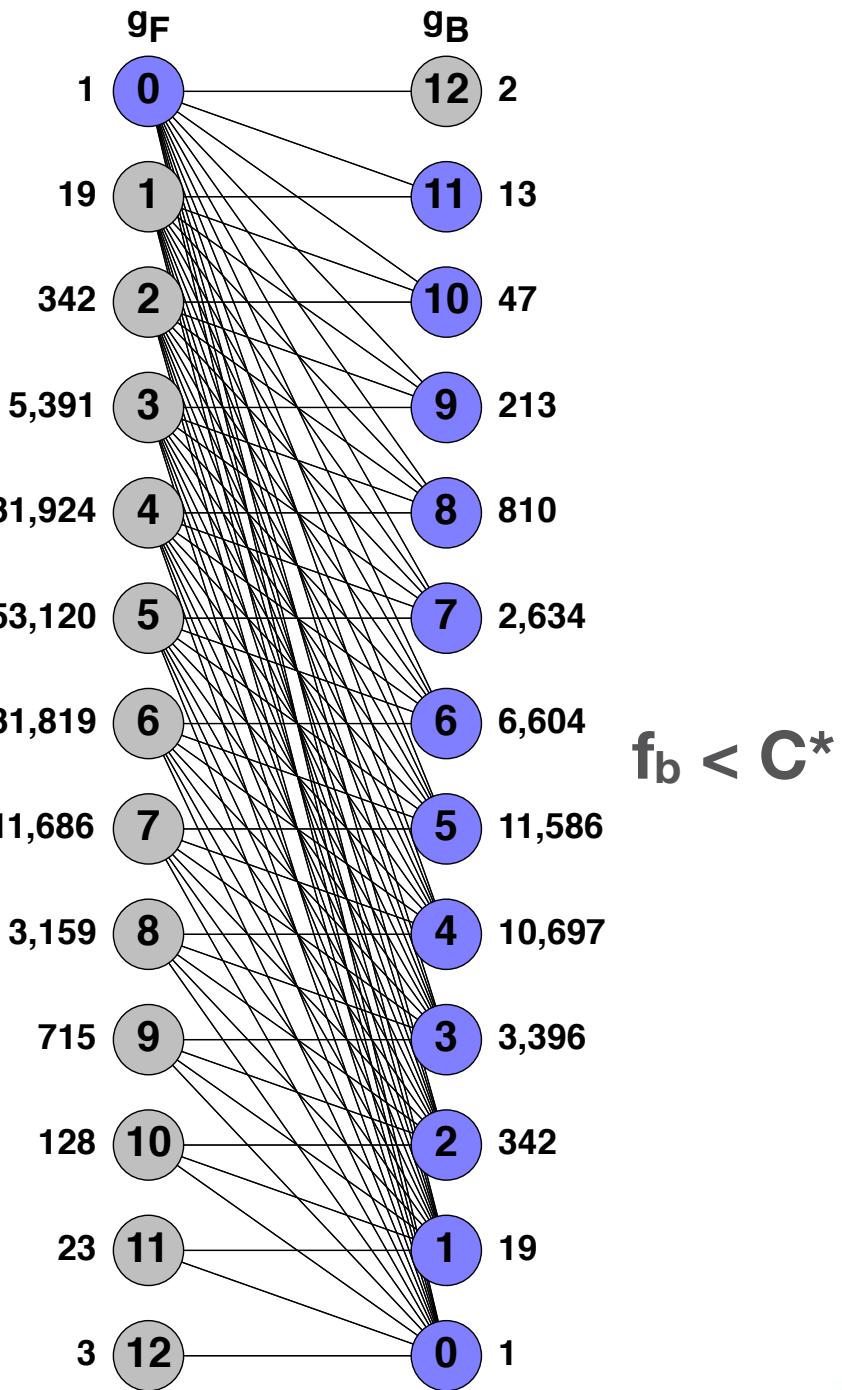


$f_b < C^*$

EXPLANATION

$f_f < C^*$

$C^* = 13$

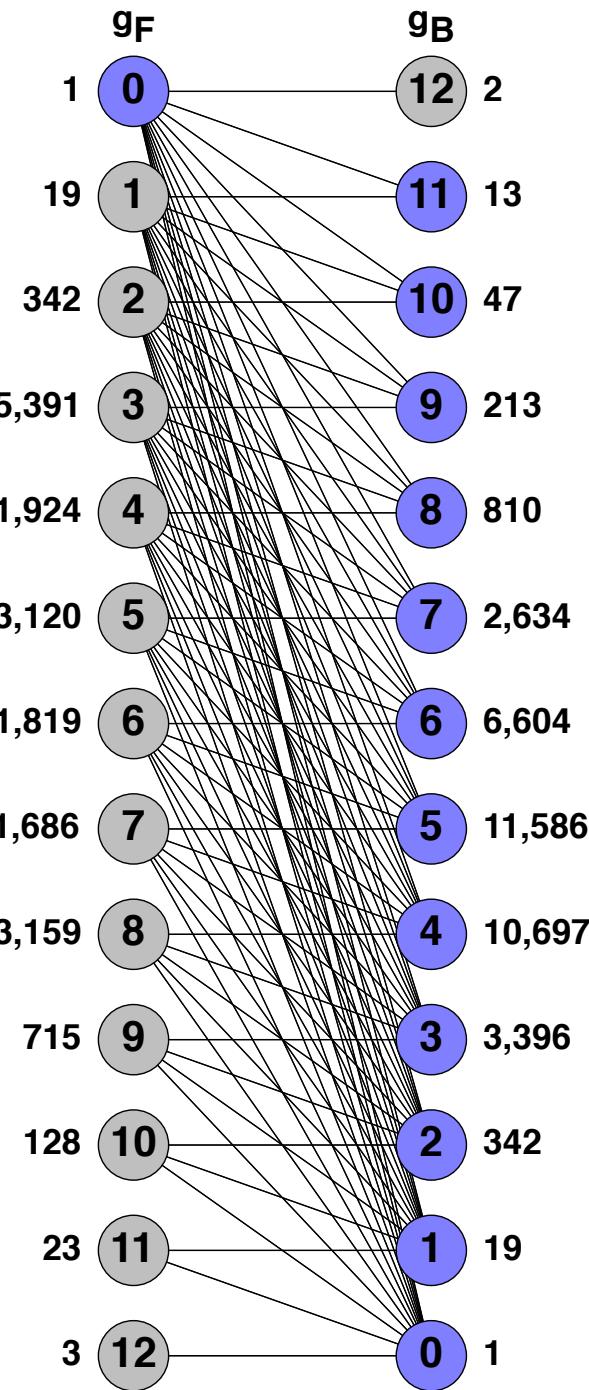


$f_b < C^*$

EXPLANATION

$$f_f < C^*$$

$$C^* = 13$$



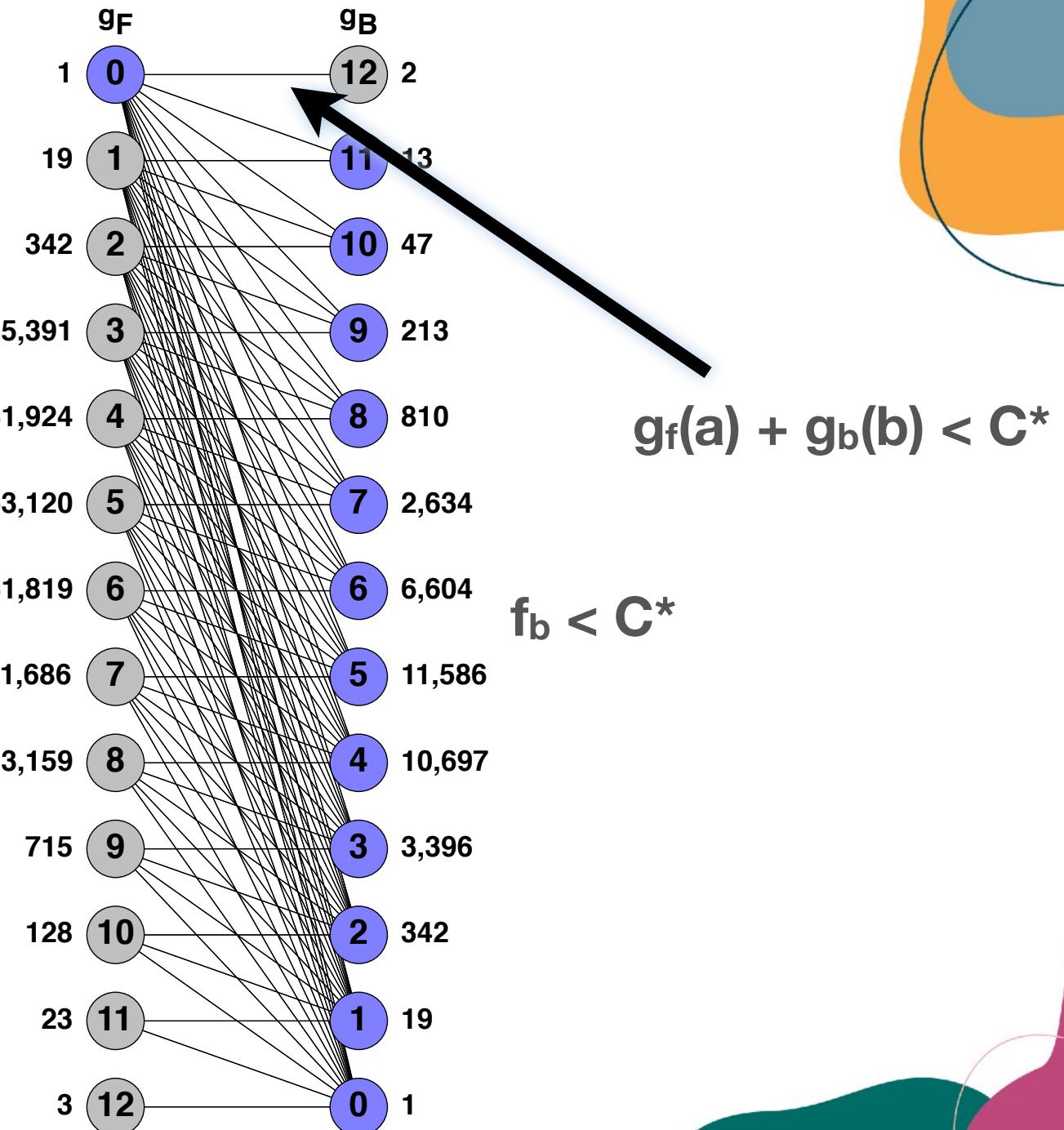
$$g_f(a) + g_b(b) < C^*$$

$$f_b < C^*$$

EXPLANATION

$$f_f < C^*$$

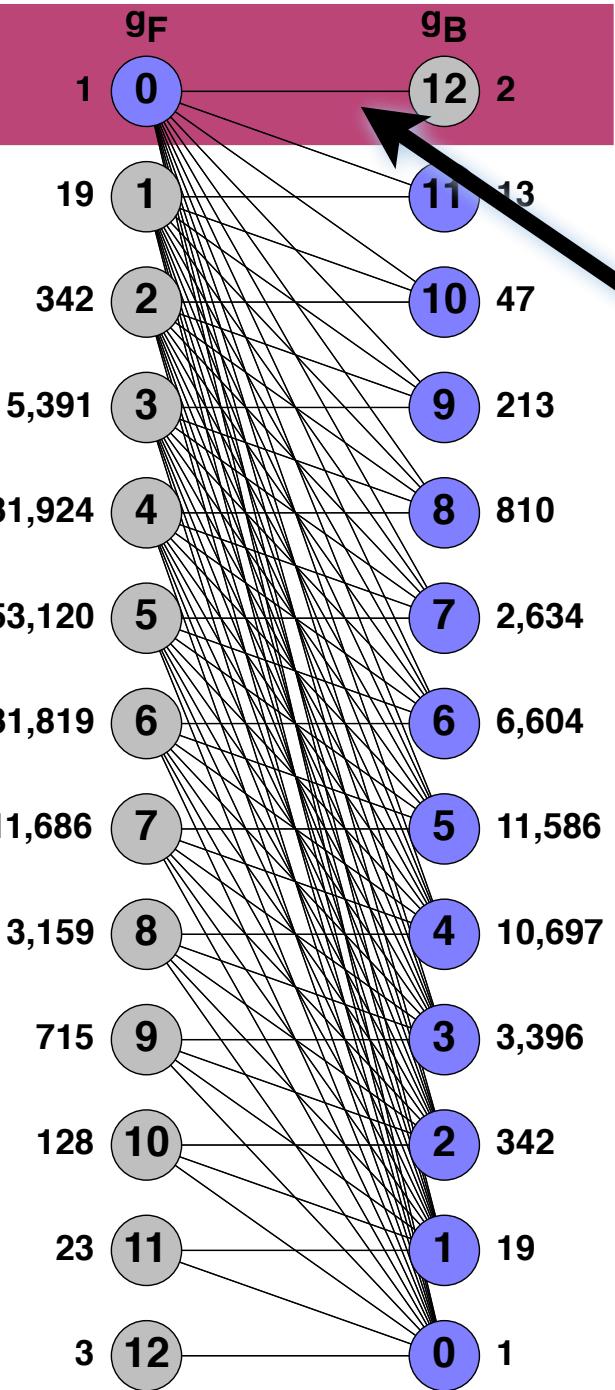
$$C^* = 13$$



EXPLANATION

$$f_f < C^*$$

$$C^* = 13$$



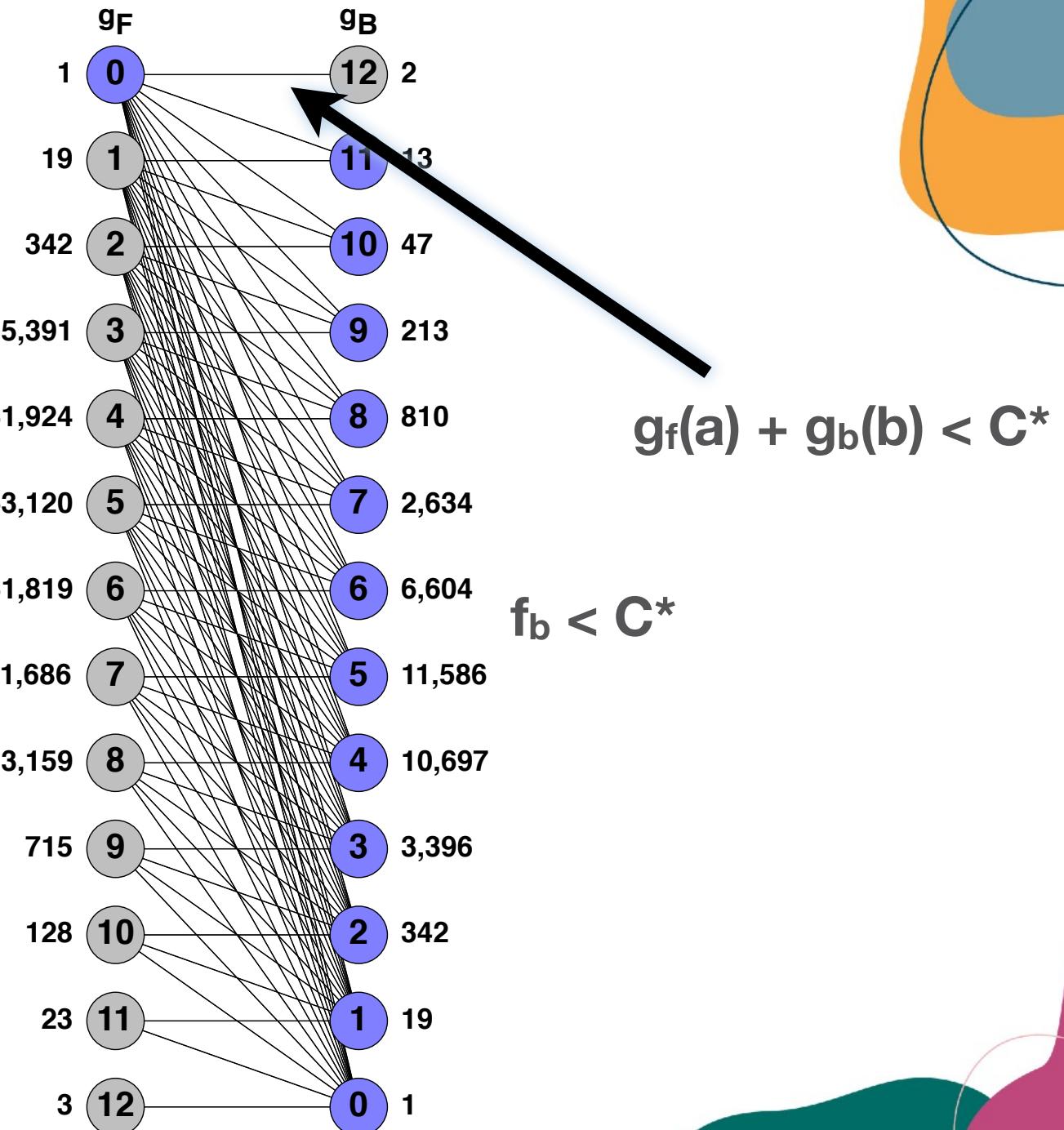
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EXPLANATION

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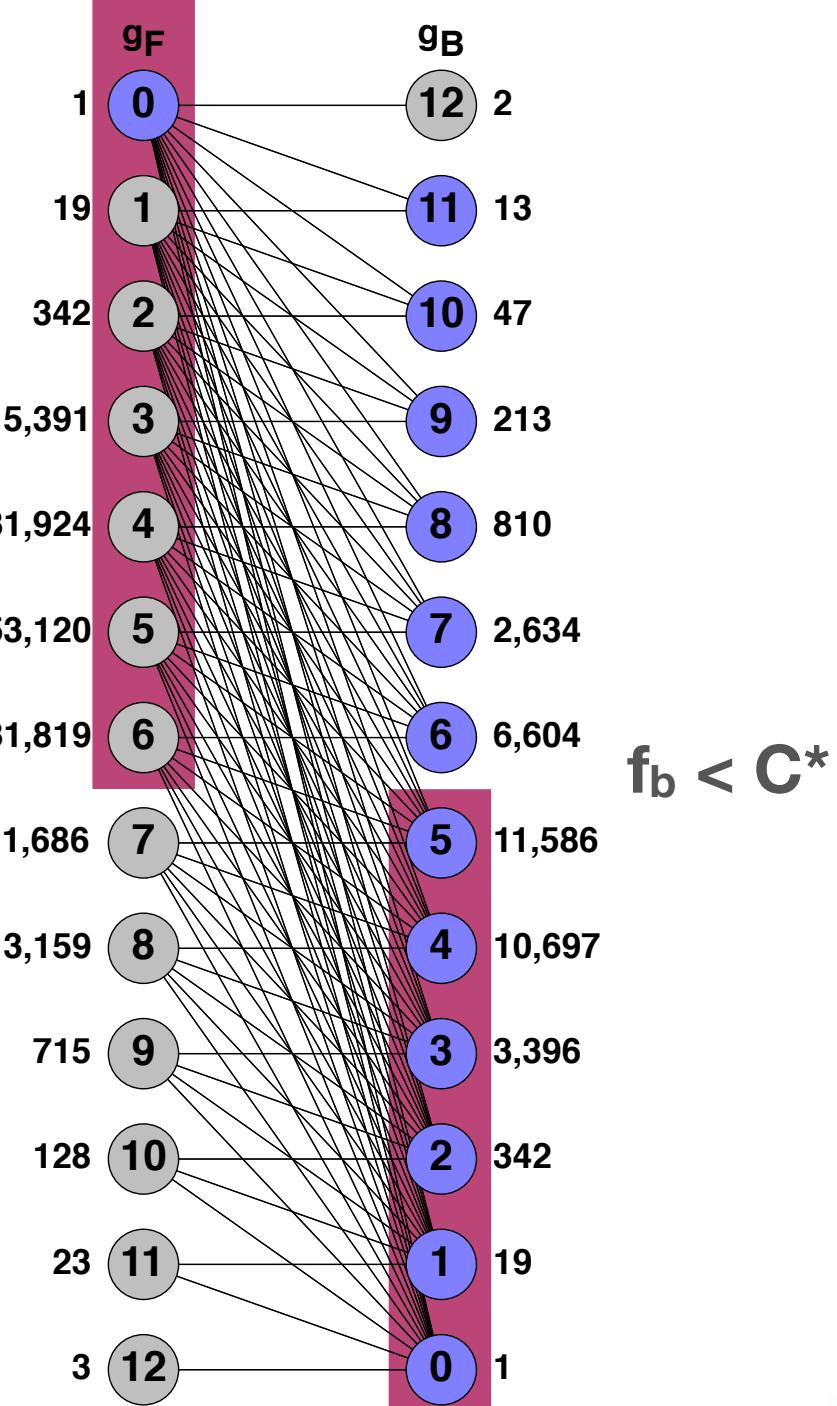
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EXPLANATION

$f_f < C^*$

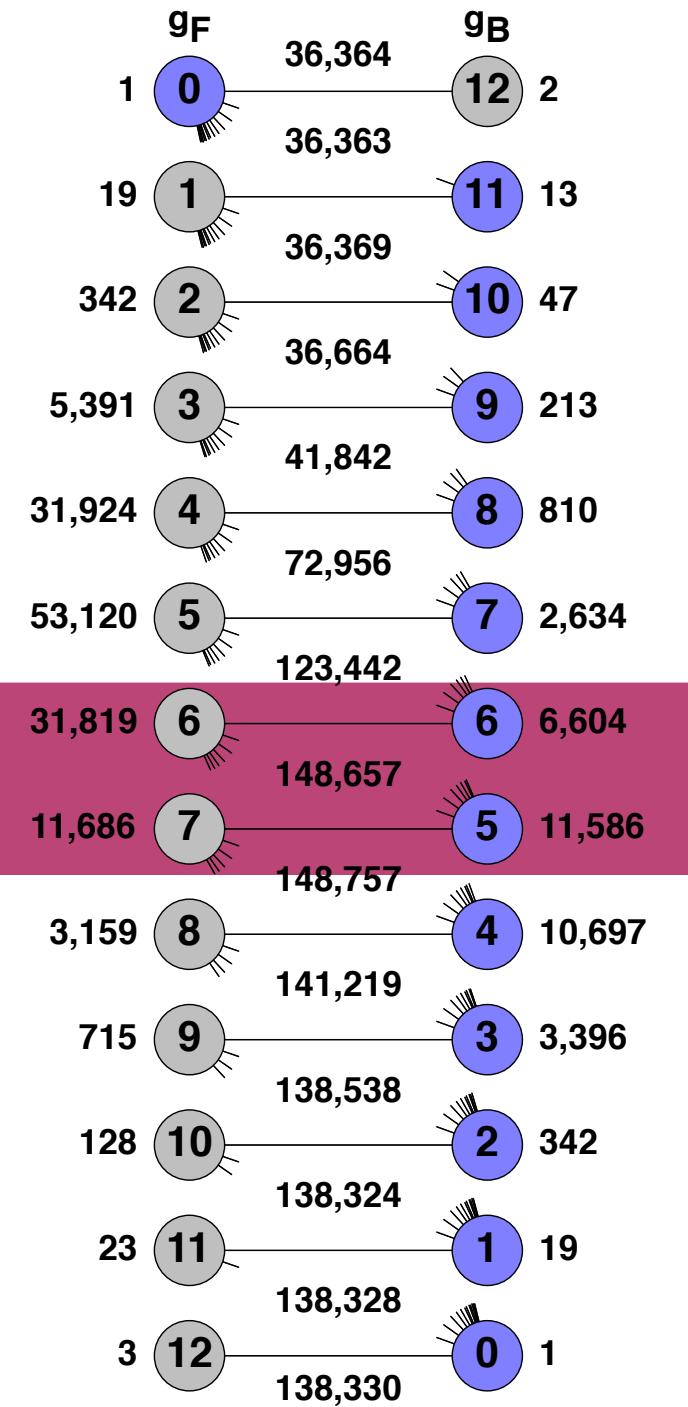
$C^* = 13$



$f_b < C^*$

EXPLANATION

$C^* = 13$



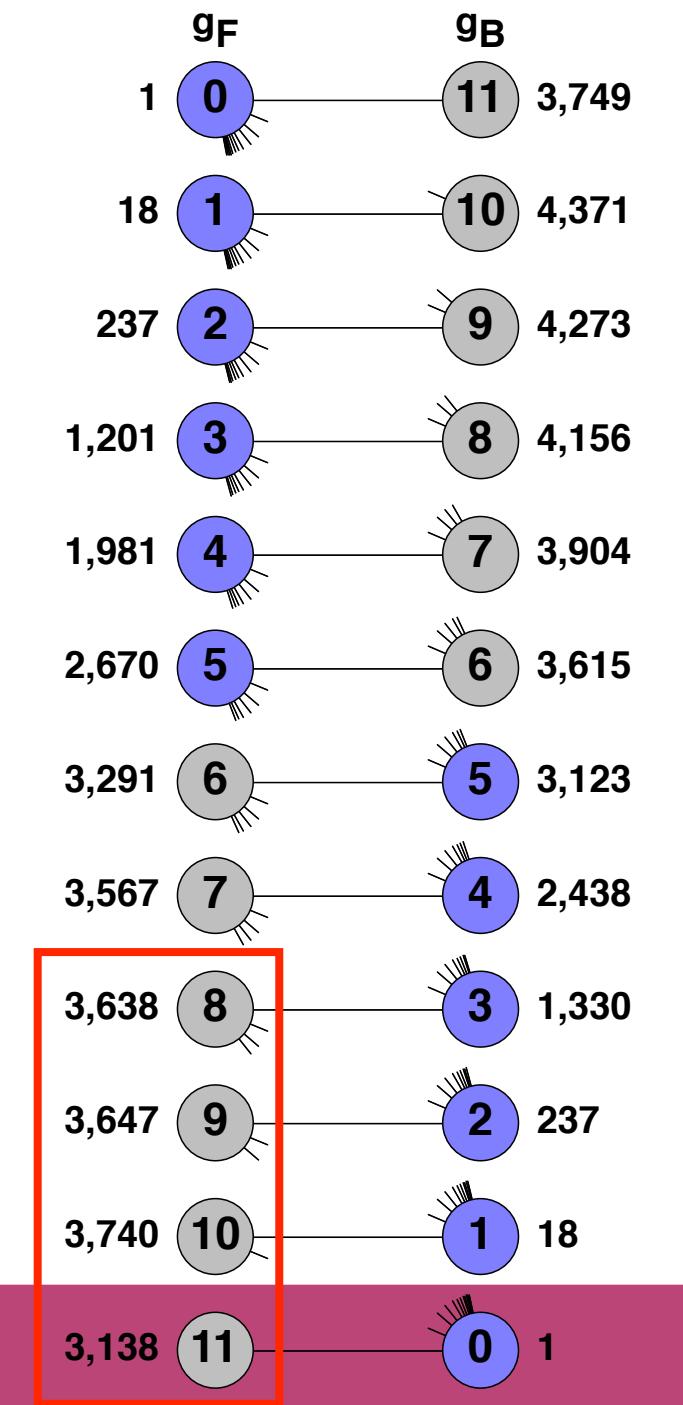
RUBIK'S CUBE 7-TILE PDB

$C^* = 12$

	g_F	g_B	
1	0	11	3,749
18	1	10	4,371
237	2	9	4,273
1,201	3	8	4,156
1,981	4	7	3,904
2,670	5	6	3,615
3,291	6	5	3,123
3,567	7	4	2,438
3,638	8	3	1,330
3,647	9	2	237
3,740	10	1	18
3,138	11	0	1

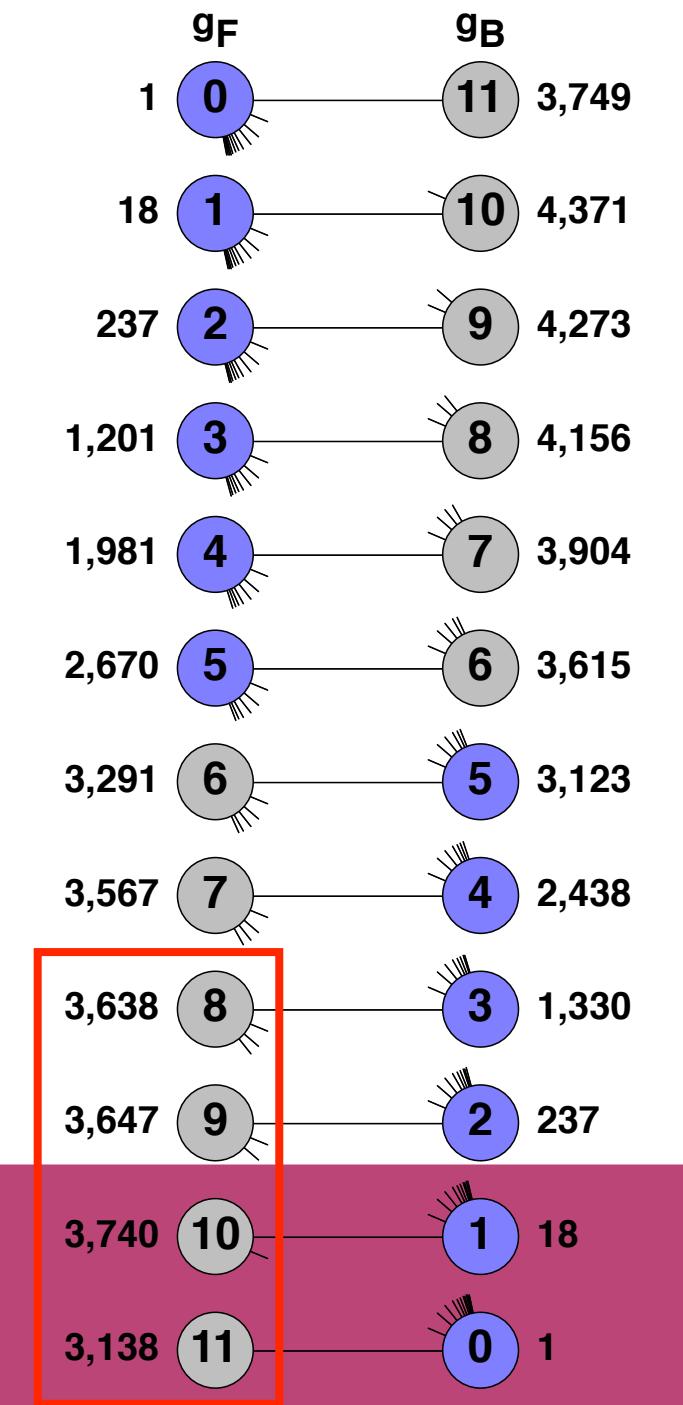
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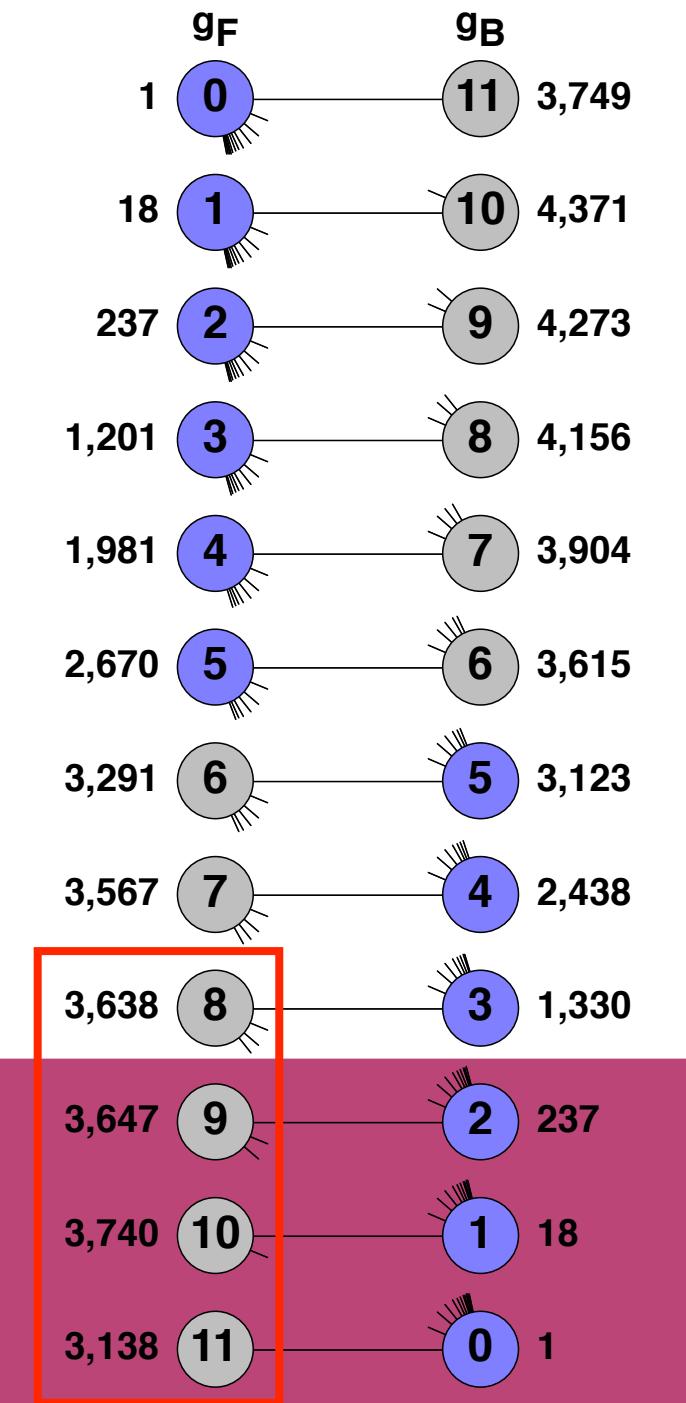
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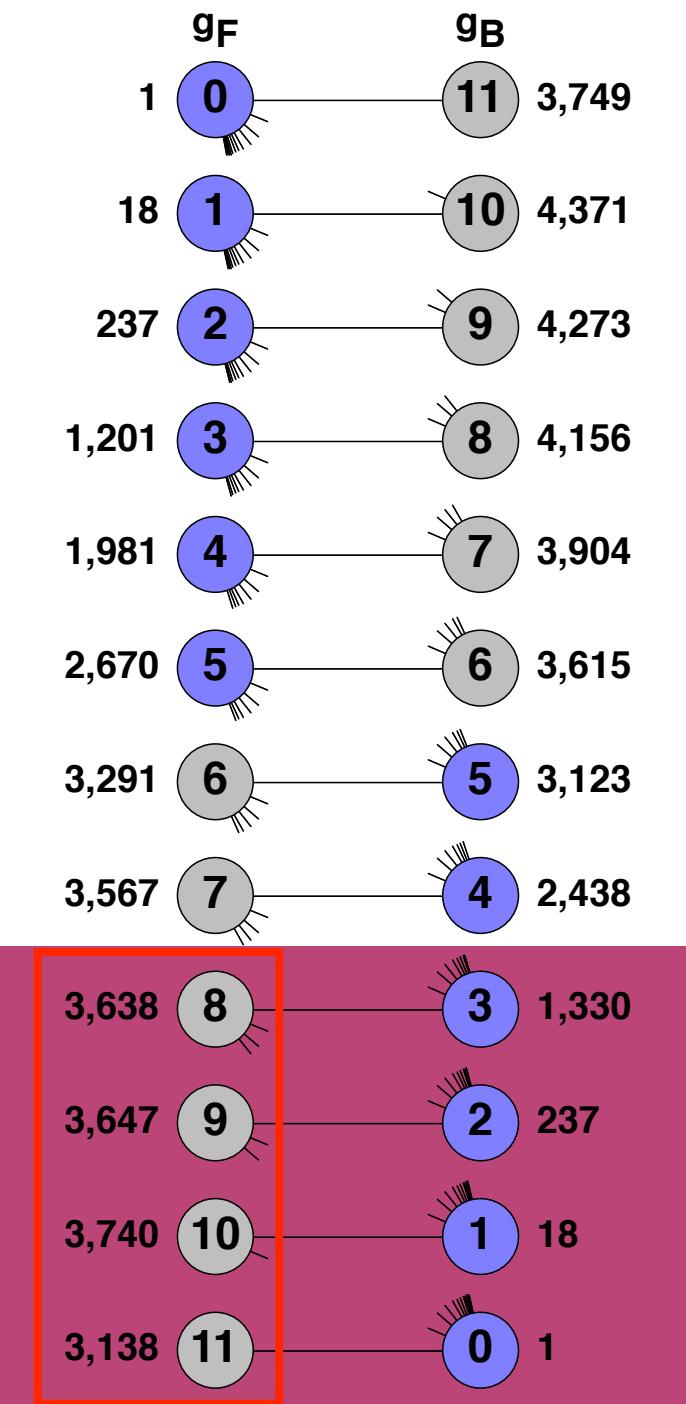
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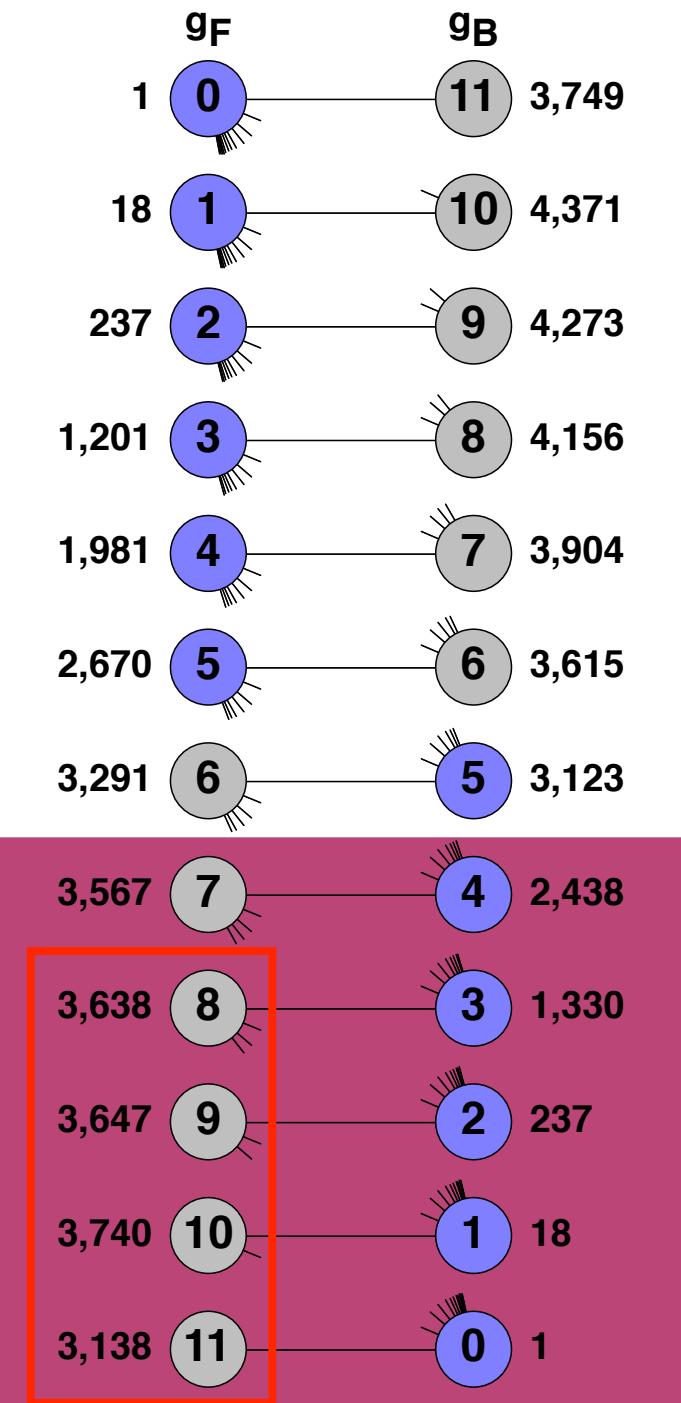
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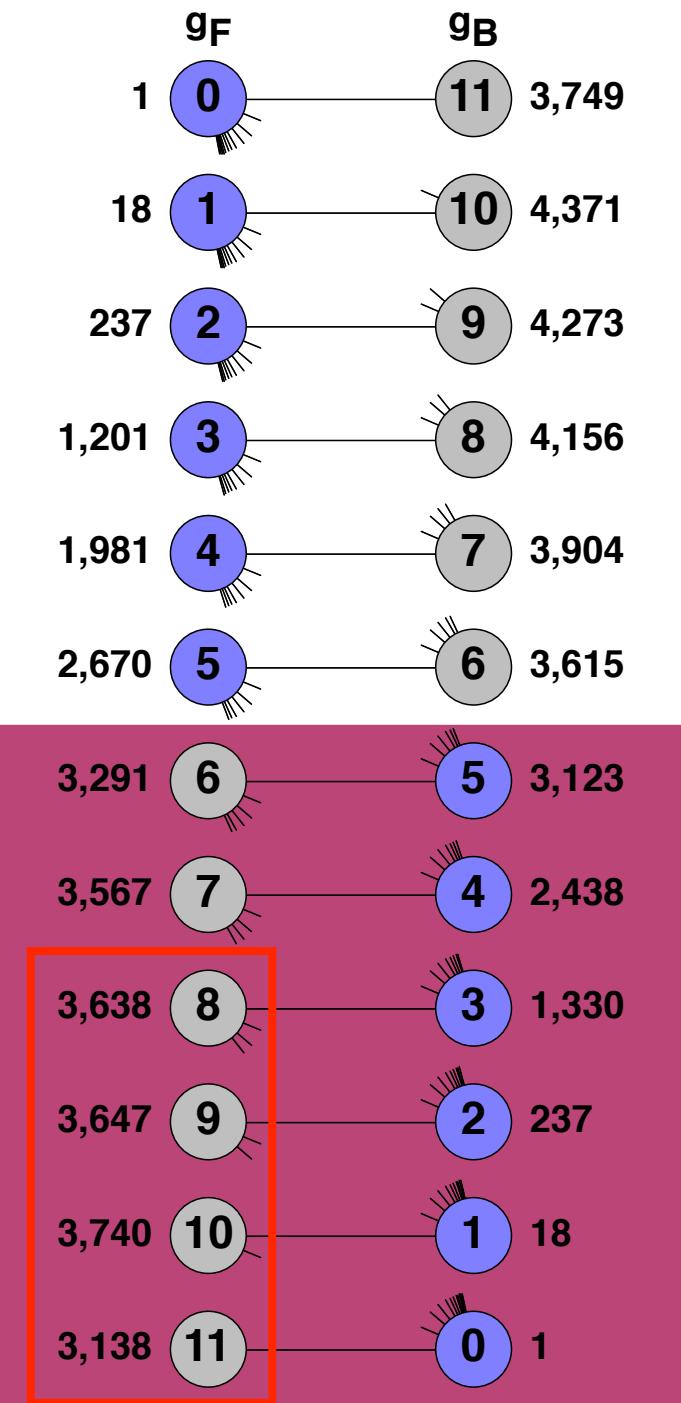
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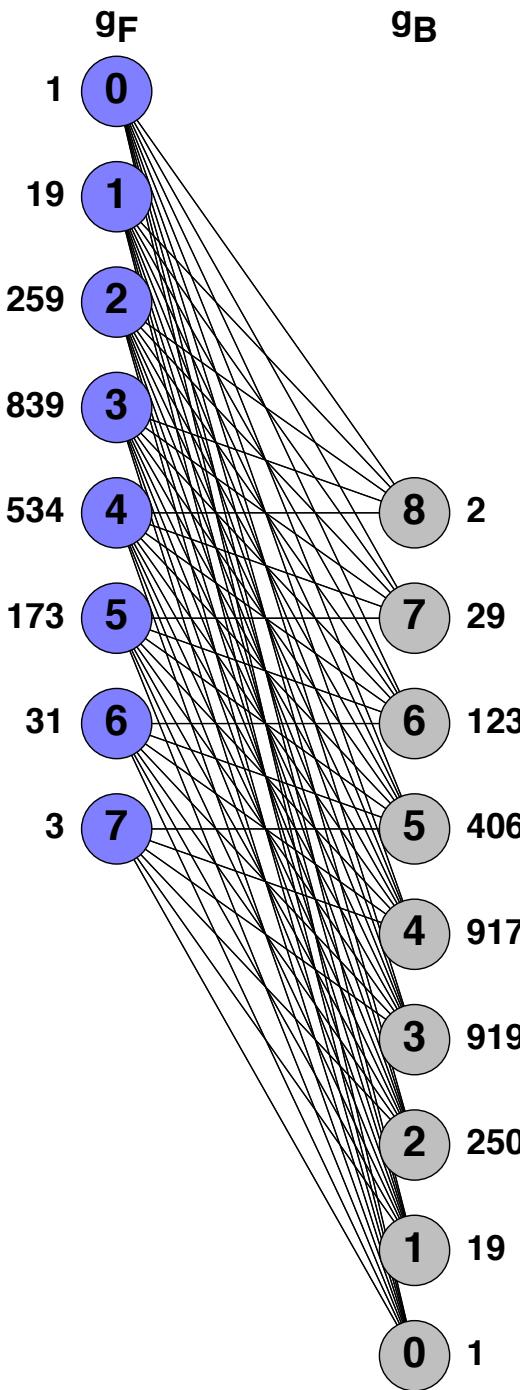
RUBIK'S CUBE 7-TILE PDB

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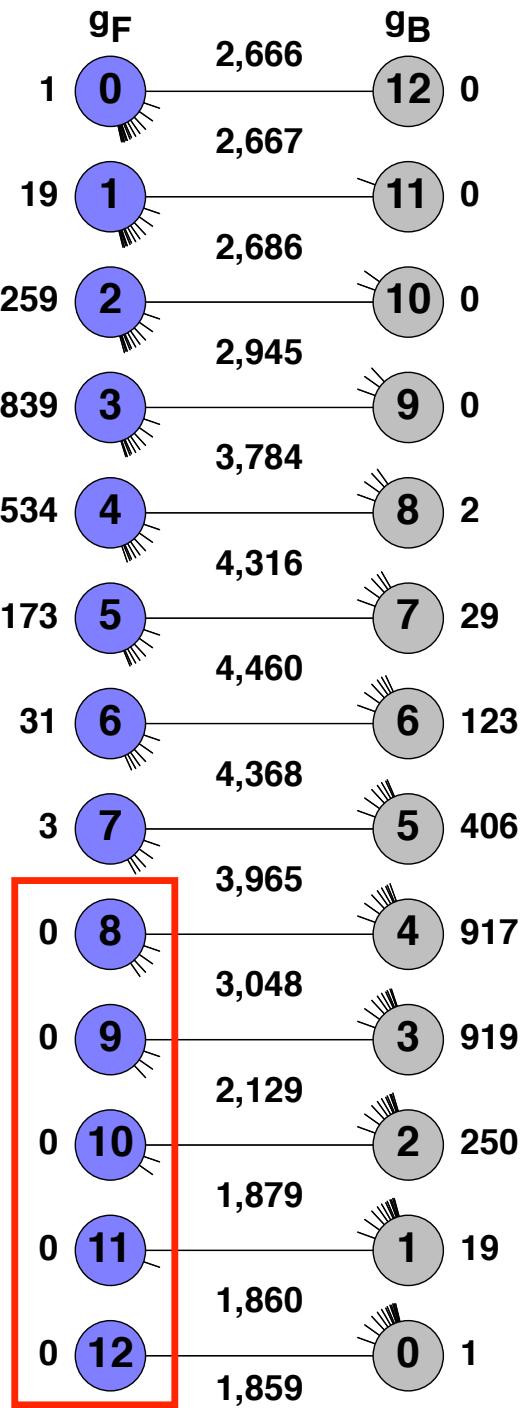
PANCAKE GAP HEURISTIC

$C^* = 13$



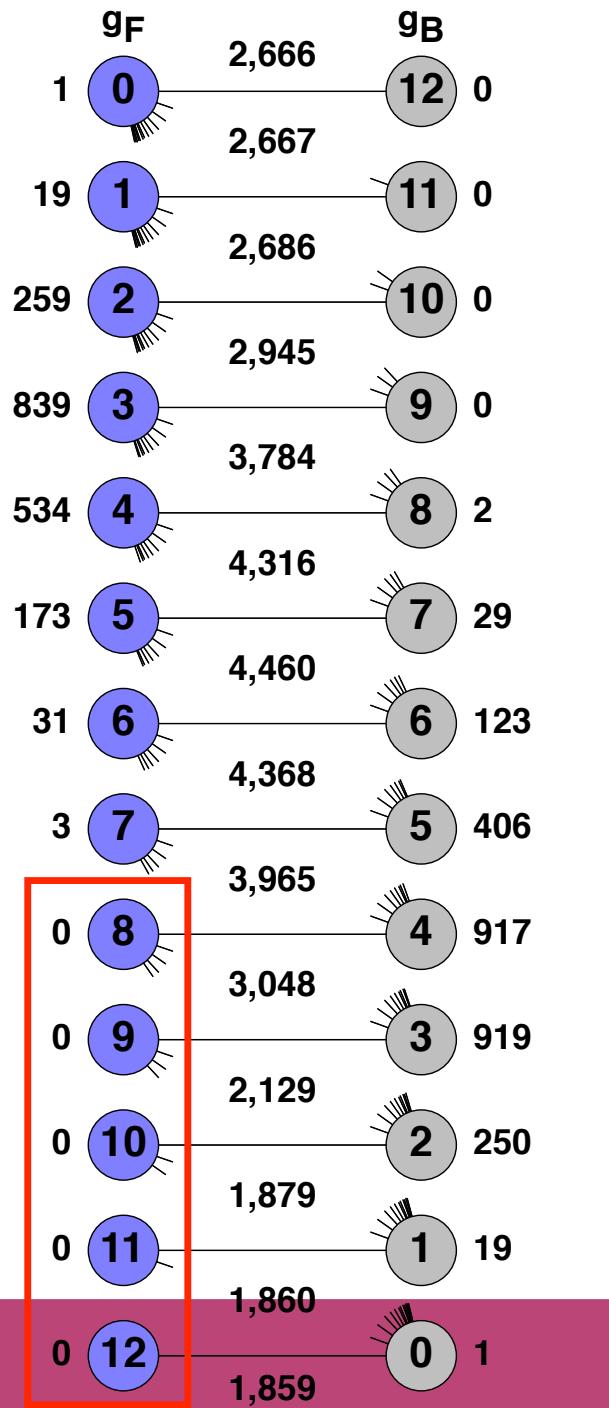
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$C^* = 13$



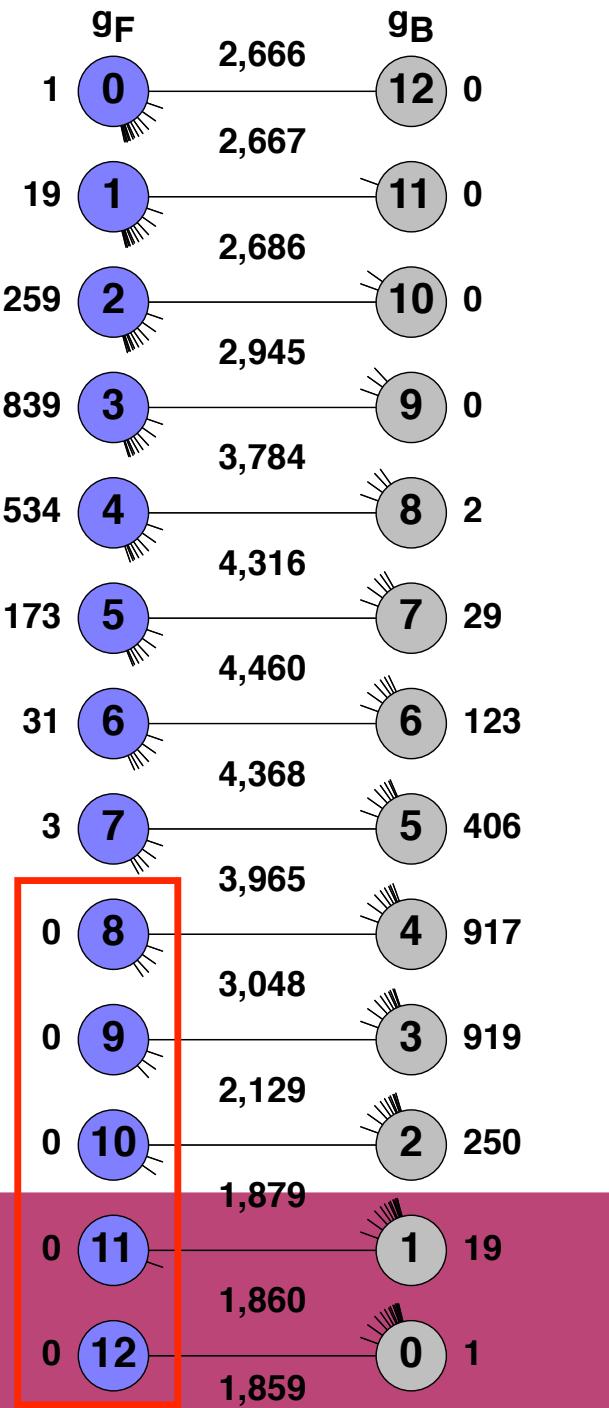
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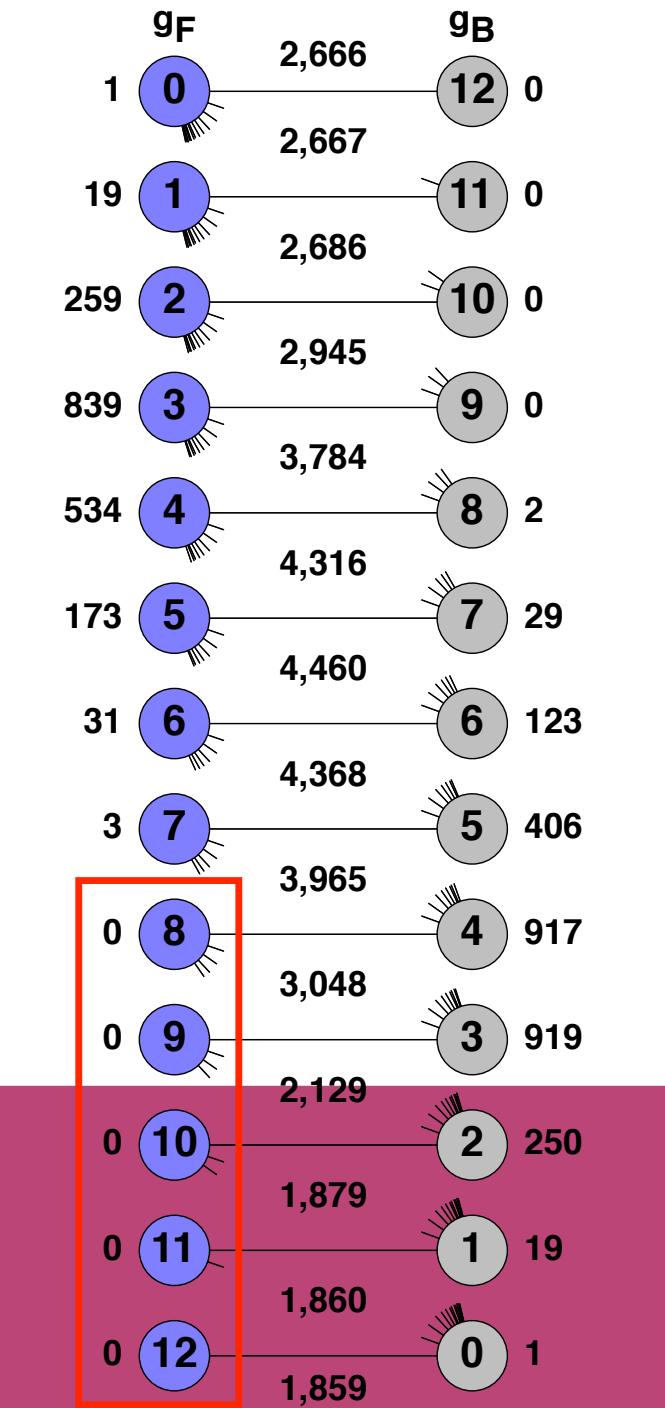
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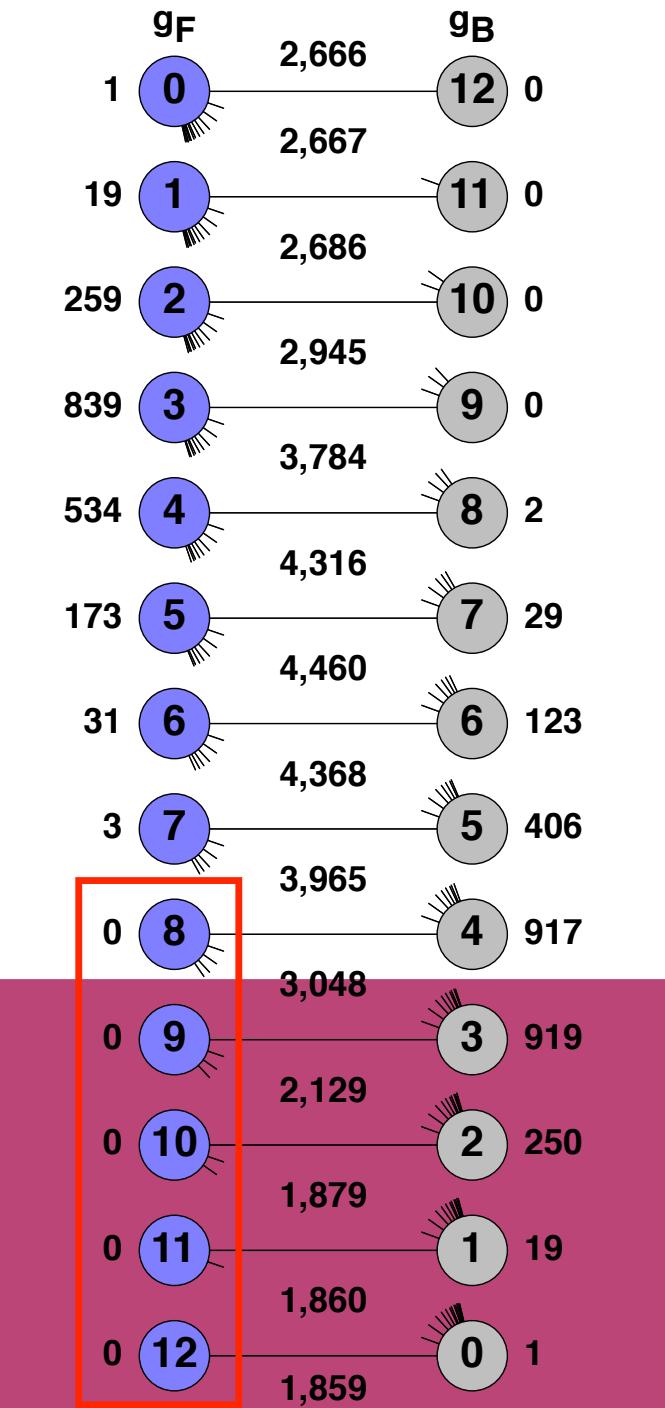
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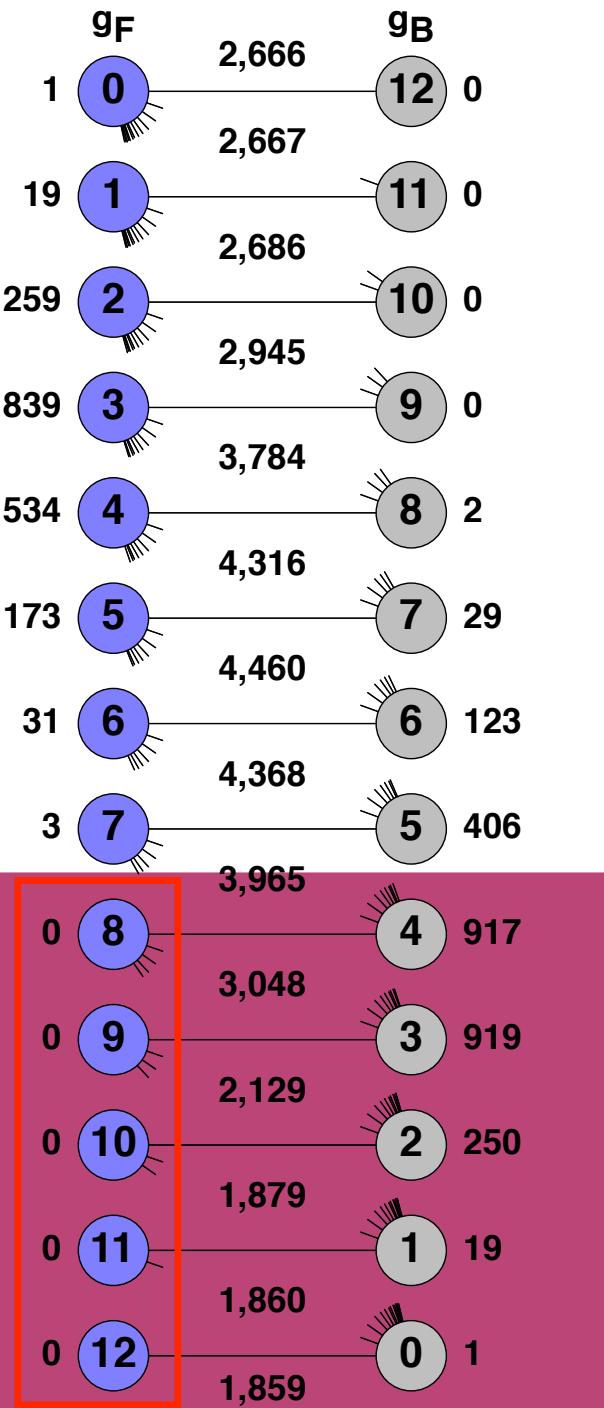
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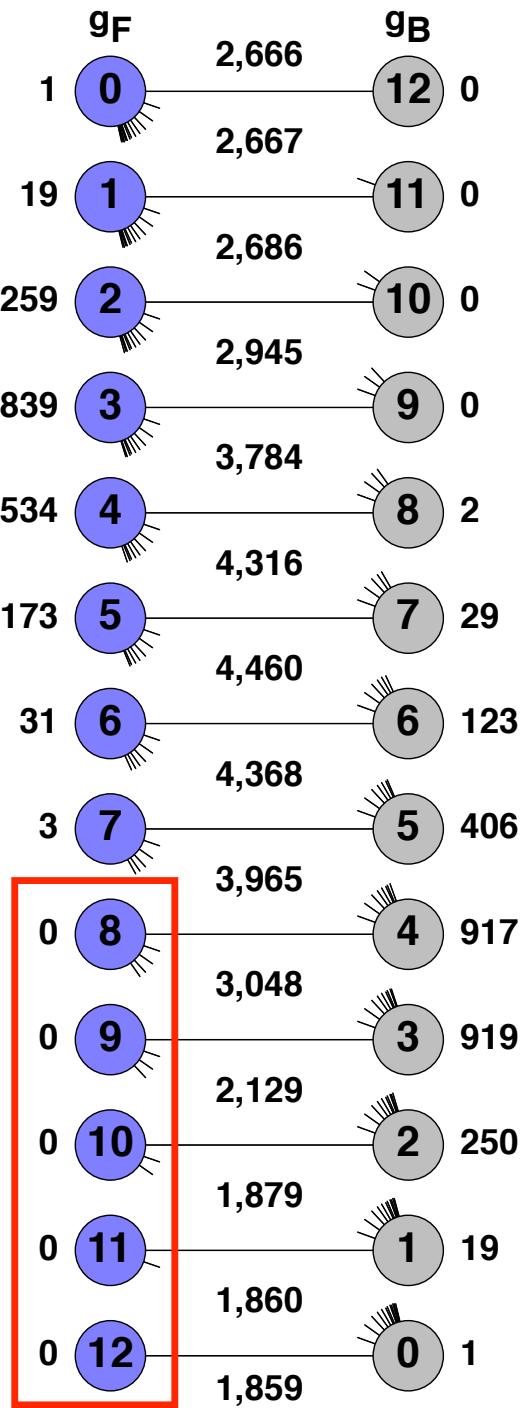
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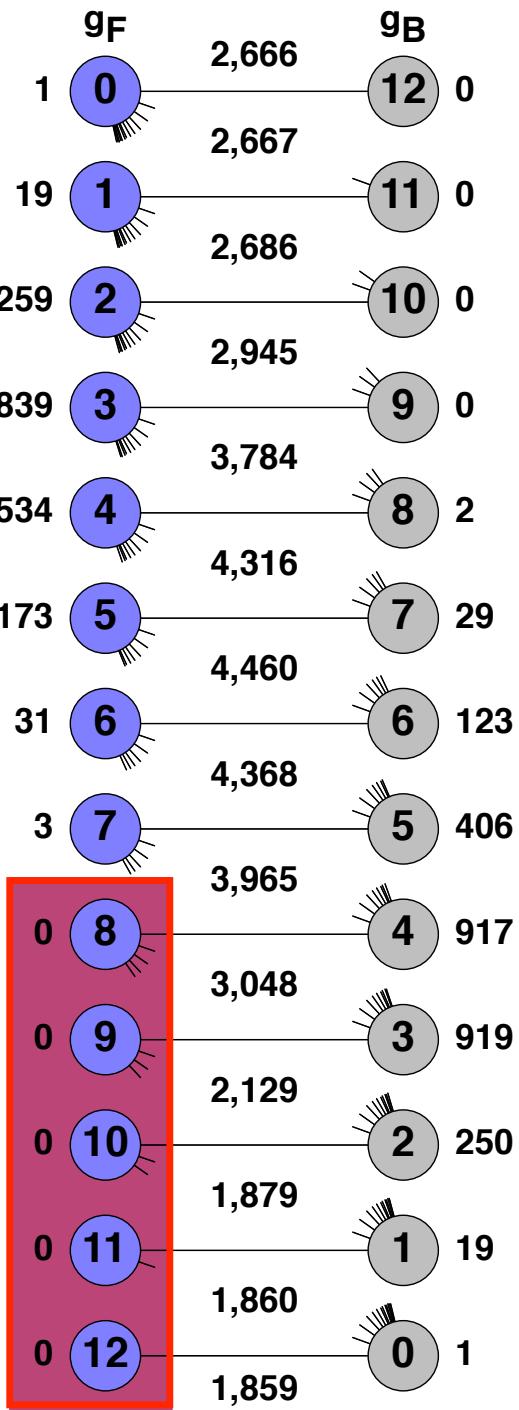
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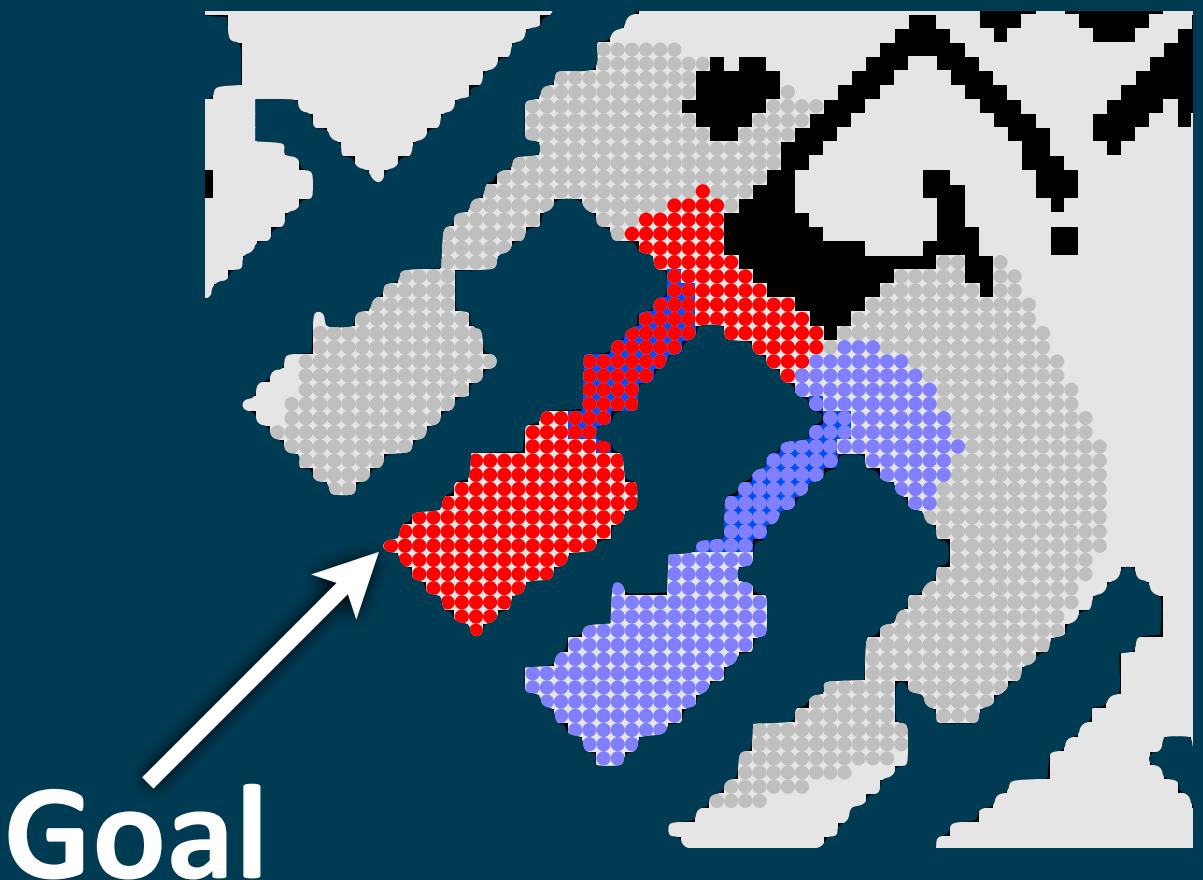
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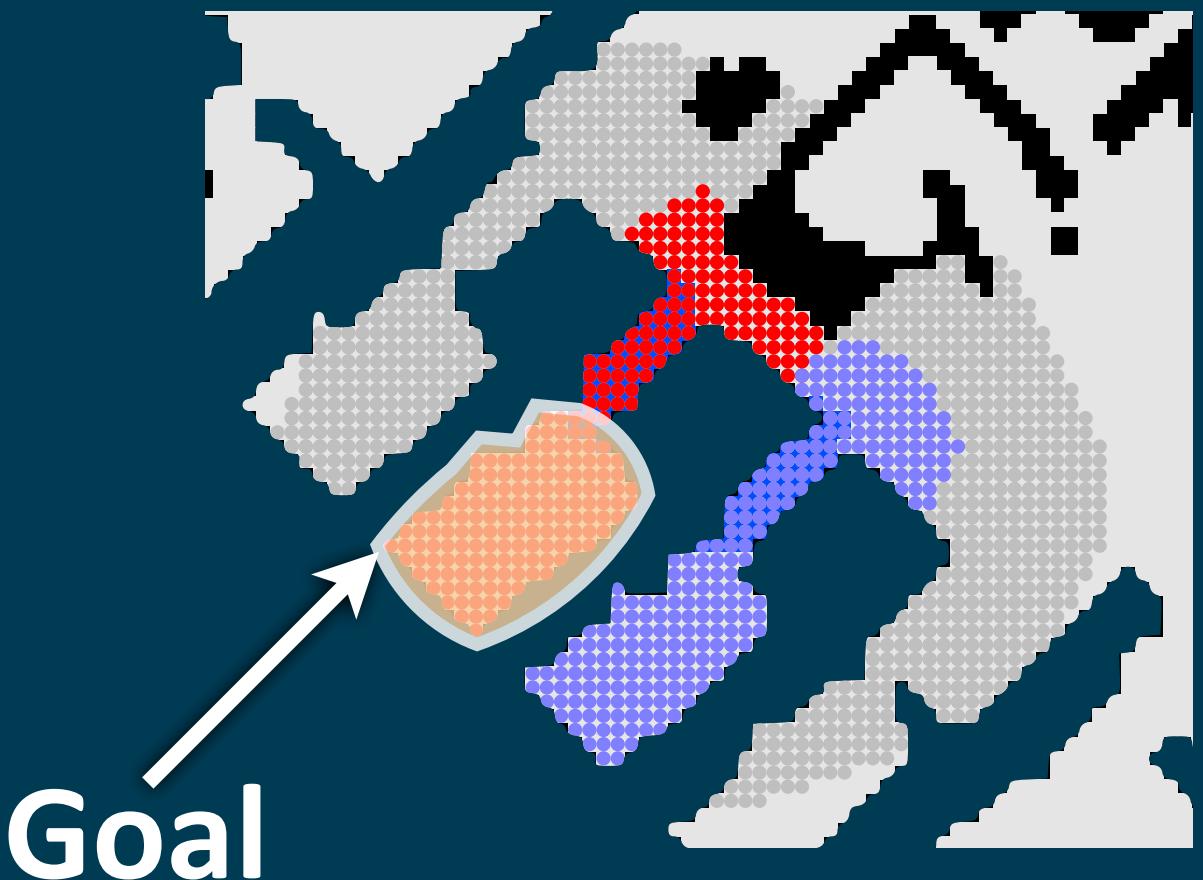
CRITICAL STATES

- Sample states near goal and find heuristic inaccuracies



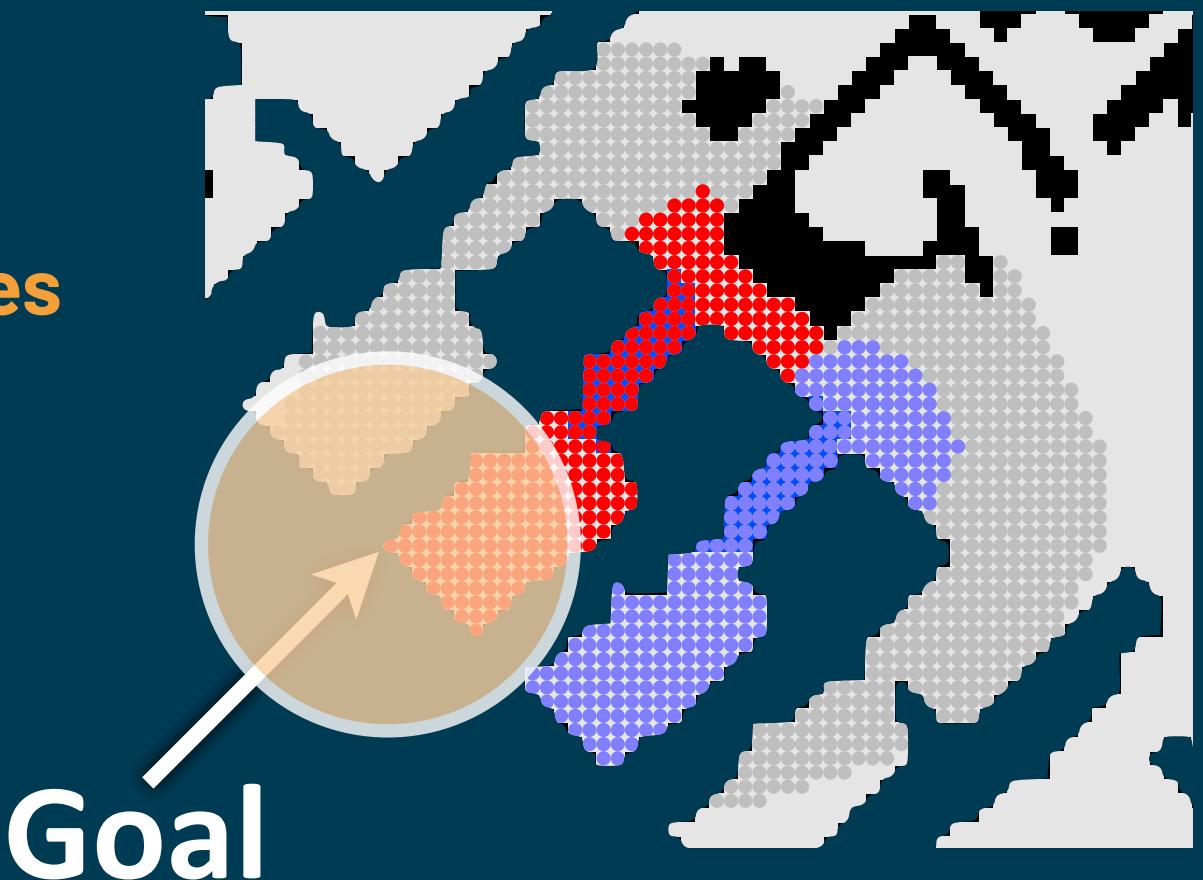
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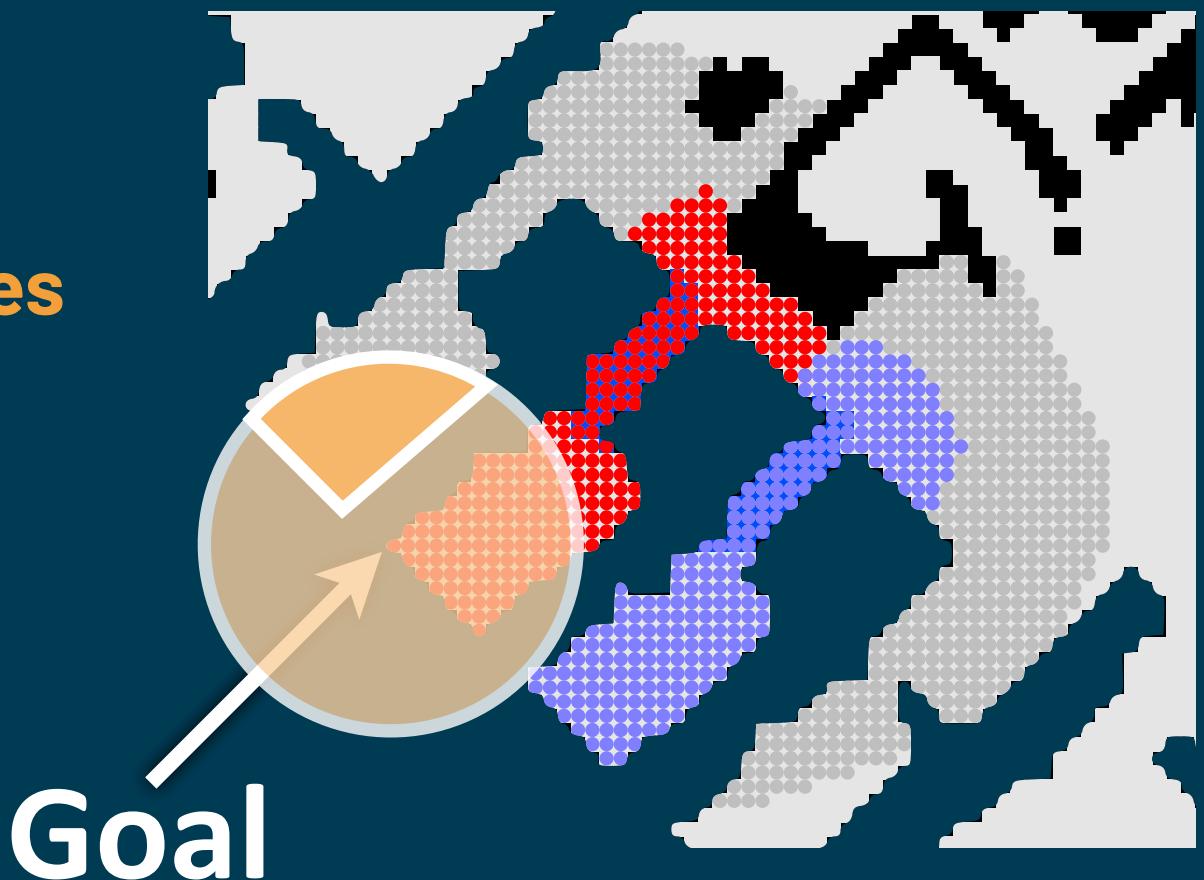
CRITICAL STATES

- Sample states near goal and find heuristic inaccuracies
- Sample the heuristic to find how many low heuristic values there are



CRITICAL STATES

- Sample states near goal and find heuristic inaccuracies
- Sample the heuristic to find how many low heuristic values there are



ASYMMETRY

- Sample states near goal and find heuristic inaccuracies
- Sample the heuristic to find how many low heuristic values there are
- Look at the asymmetry of the state space



EXAMPLE: TOH WITH PDB HEURISTIC

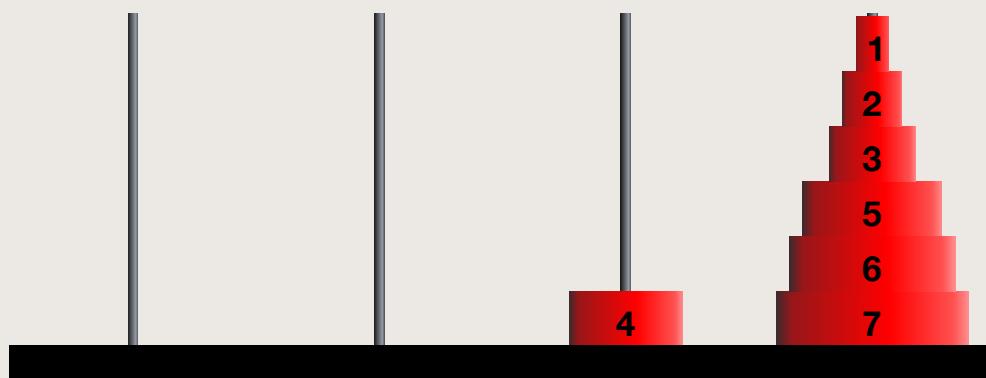
$\phi(s)$ is close to the goal and has a low heuristic value

s is far from the goal

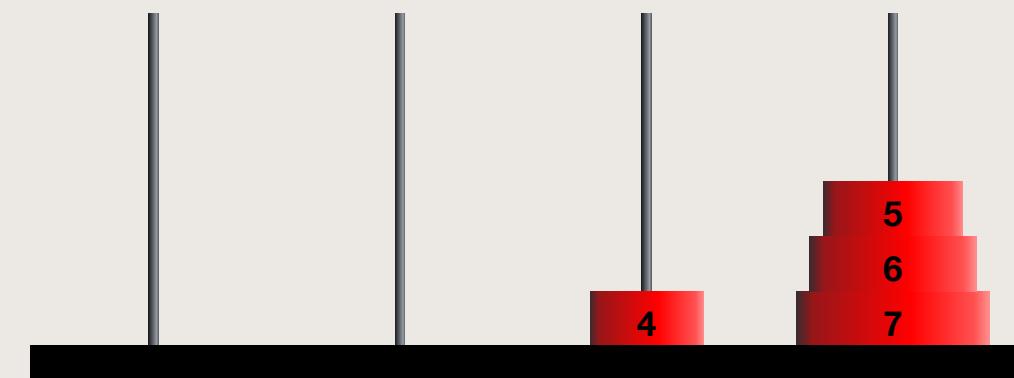
Low heuristics are inaccurate

Bidirectional heuristic search *outperforms* unidirectional search

Problem State s

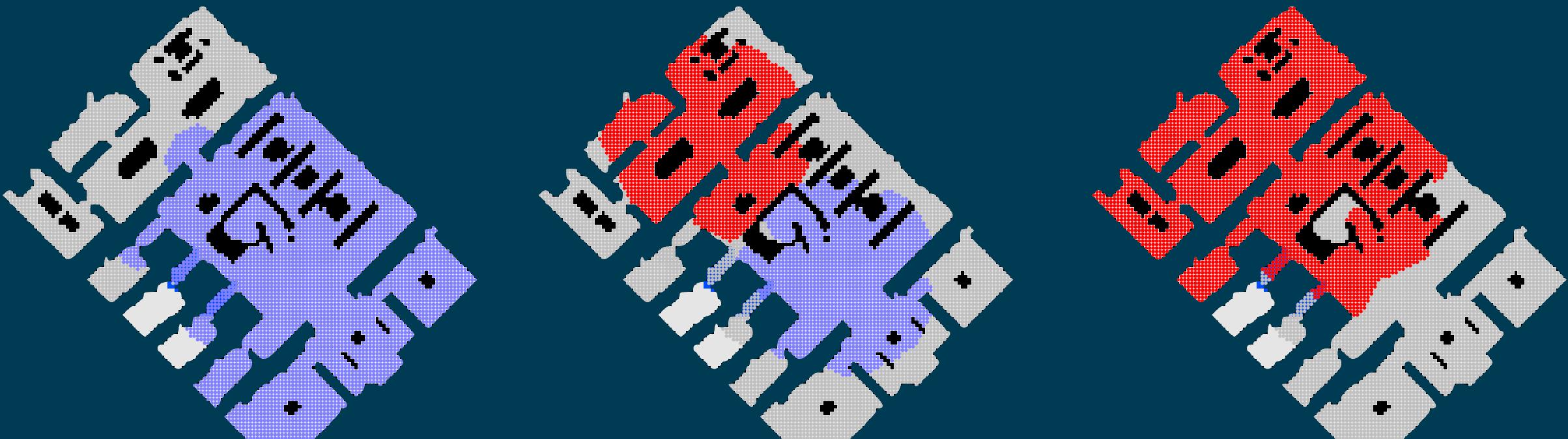


PDB Abstraction $\phi(s)$



EXAMPLE: UNINFORMED SEARCH

Low and inaccurate heuristics for almost all states



CONCLUSION

Critical states have both low and inaccurate heuristics

Need critical states for bidirectional search to perform well

More critical states → bidirectional search will do better

<https://webdocs.cs.ualberta.ca/~nathanst/papers/sturtevant2020unibidi.pdf>