## Abstraction in Pathfinding

with a focus on commercial video games
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## Abstraction

- Most search problems can be represented by a graph
- Build a smaller graph which retains most relevant information in the original graph
- Similar to a low-resolution image
- (Holte 96; Bulitko et. al. 07)


## Bridging contexts...

- Sven discussed many techniques for enhancing $A^{*}$ search
- There is another dimension along which we can optimize the performance of pathfinding algorithms



## Why abstraction?

- Abstract graph is smaller
- Search is cheaper
- Defines subgoals in search
- Can be used for optimal or suboptimal solutions
- Pattern databases


## How is abstraction used?

- Almost all commercial video games use some type of abstraction
- Unreal engine has Kynapse A.I. plug-in
- Automatically builds high-level graph
- Units can only walk on the graph
- Will describe the system built for BioWare Corp for their upcoming title Dragon Age (Sturtevant, 2007, AIIDE)




## Motivation

- Games have tight memory budgets
- ~5MB total memory for map data
- 1024x1024 or larger maps
- 1MB per byte per grid cell
- Can we use build an abstraction which minimizes memory usage?
- Total memory usage by abstraction
- Memory used during planning


## Motivation

- Don't computers have lots of memory now?
- Develop / design for low end
- Models / graphics expensive
- New gaming platforms
- Nintendo DS
- iPhone


## Motivation

■ Need to speed up search

- Previous pathfinding engine was taking up to 100 ms to plan
- Ideally should plan in 1 ms or less
- 3 -5ms for all planning per frame
- May need to handle many units in the same time frame


## Assumptions

- Grid world
- No true 3-d movement
- Cells can be blocked/free/weighted
- May be height difference between cells
- Units can move across real-valued space


## Solution

- Build abstract graph from low-level data
- Divide world into sectors
- Sub-divide into regions
- Maintain connectivity information between sectors


## Sectors / Regions

- Divide world into large sectors
- Fixed size
- Index implicitly
- Divide sectors into regions
- Regions entirely connected
- Regions have a center point


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## Abstract Graph

- Original Map:
- 32x32 = 1024 cells
- Abstract Graph:
- 9 nodes
- 10 edges



## Edges

- Look at borders of regions to determine edges



## Memory Usage

- Sector data
- Fixed size (32 bits)
- Region data

| Sector Data |
| :---: | :---: |
| \# Regions |
| Memory Address |
| unused |

- Variable sized
- 8 bit region center
- Edge count
- 8 bits per edge


## Memory Usage

- Can save more memory:
- 16 bits for sectors

- "Default" regions
- Edges stored twice
- Other optimizations



## Find Abstract Region

- Begin with $x / y$ location in real world
- Sector implicit
- If sector only has 1 region, done
- Otherwise do BFS to find region center
- Extra bits in grid can store region info
- Pointers not needed


## How is abstraction used?

- Need to find path between two points in the actual map
- Find abstract region
- Find abstract path
- Refine



## Find Abstract Path

- Given sector/region for start and goal:
- Use A* to find a complete abstract path
- Use Manhattan/octile distance between region centers as both heuristic and edge cost


## Refine

- Many different ways to use abstract path
- Simplest method:
- Find path from start to first region
- Find path to successive region centers
- Find path from last region to goal


## Analysis

- Total pathfinding cost
- Optimizations
- Region center placement
- Reducing suboptimality
- Experimental verification


## Usage Example

- Find abstract parents
- Find abstract path
- Find real path



## Total Pathfinding Cost

- Abstract planning
- Depends on path length, sector size
- Refinement
- Depends on path length, edge refinement (region centers, suboptimality)
- Maximize sector size


## Optimizing Region Centers



Optimizing
Region Centers

- How to determine the region centers?
- Some locations are much better than others
- Harder with larger sector sizes


Optimizing

## Region Centers



Optimizing Region Centers

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## Optimizing <br> Region Centers

- In a sector, for each cell in a region:
- Measure A* cost to plan a path to each neighboring region from that cell
- Choose the region center which minimizes the maximum cost
- Can optimize any cost function


## Pathfinding Optimization

- Refinement at
start/goal can be inefficient
- Trim path segments
- Skip to next
- Skip to next
region at start/ goal



## Sector-Related Errors

- All points within a sector/region are treated equally
- Adjust abstraction when performing search



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## Experimental Results

- 93,000 paths over 120 maps
- Maps scaled to $512 \times 512$
- Paths length
1... 512
- Evaluate:
- Memory
- Region center optimization
- Optimality
- Total Work


## Memory Usage

- How does the memory usage scale with sector size?
- How much memory can be saved with simple compression?
- Don't store "default" sectors with

1 region, 8 neighbors

## Memory Usage

Maps Size 512x512


## Dynamic Region Centers

- Is there a gain to dynamically optimizing region centers?
- Measure 95\% work done in one-step path refinement


## Dynamic Centers



## Optimality

- Paths will not be optimal
- Special cases for start/goal help
- Smoothing is applied as a post-processing step (not measured)


## Optimality



## Total Work

- Compare total work by sector size
- Find abstract path
- Refine low-level path
- Compare to $\mathrm{A}^{*}$

Total Work


## Implementation

- Custom implementation for upcoming title Dragon Age ${ }^{\text {TM }}$ (BioWare Corp ${ }^{\circledR}$ )
- Worked in-game during parts of 4 months
- Initial implementation took two weeks
- Rebuilt pathfinding core
- Spent $\sim 4$ weeks optimizing code, adding smoothing, control structures, etc


## Total Work v. A*





## Final Performance

- About $0.1 \mathrm{~ms}(100 \mu \mathrm{~s})$ per step of planning
- Find abstract path
- Refine 1 edge from abstract path
- Interleave planning and acting
- Can plan for 30-50 units every frame
- Units do not need to plan every frame
- Can "gracefully" degrade performance
- Units offscreen don't need to smooth


## Memory Usage

- Memory usage is well within requirements
- Very little memory needed on a per-unit basis for planning
- Abstract path
- Current path
- State of planning/smoothing


## Data Structures

- $A^{*}$ uses:
- Open list -- usually a heap
- Closed list -- hash table
- Back pointers -- reconstruct path
- Can't store these on the map
- Simple implementation occasionally slow
- Allocate small closed list for each sector
- Can quickly be cleared; no deallocation



## Summary

- Units walk on real-space
- Abstract into a high-resolution 2-d grid
- Abstract again into coarse graph
- Units pretend to live on high-resolution grid
- Michael will talk about getting rid of the 2-d grid


## That's great but...

- In many domains, pathfinding involves multiple units
- How can units cooperate when planning?
- Ignore each other and replan
- Using 'flocking' methods to avoid other units
- Explicitly cooperate
- (Dresner and Stone, 2008, JAIR)
- (Silver, 2005, AIIDE)
- (Sturtevant and Buro, 2006, AIIDE)


## Dresner \& Stone

- Traffic management problem
- Can cooperative cars increase traffic throughput?
- Centralized system manages reservations
- Can a car get through the intersection safely?
- Tries several different speeds
- Forces cars to wait until the can get a reservation


## Data Structure




## Generalized Cooperative Pathfinding

- Goal: Multiple agents cooperate during path planning and execution
- Generalized travel (eg no lanes)
- Centralized reservation system
- Use abstraction to reduce costs






## Possible Strategies

- Plan all units simultaneously
- Computationally intractable
- (unitsactions) depth
- Plan individual units
- Not complete
- A lot of techniques needed to be practical


## Overview

- Why problem is hard
- What techniques simplify the problem
- Improving performance with abstraction
- Evaluation


## WHCA*

- Use a hash table to store time-space indexed reservations
- Constant time acces
- Is a space/time cell free?
- Reserve a space/time cell
- Free a space/time cell


## WHCA* ${ }^{*}$ w)

- Windowed Hierarchical Cooperative A*
- Cooperative A*
- Hierarchical Heuristic
- Windowed cooperation
- Silver, 2005


## A*

- A* relies on a heuristic to guide search
- Poor heuristics cause extra node expansions
- Cost is the area in which the heuristic is poor


## Cooperative A*

- 3-dimensional search problem
- x-location, y-location, time
- Still need a heuristic
- Cost is the area in which the heuristic is poor times the time to get out of that area = volume



|  | 10 | 10 | 10 | 10 | 10 | 10 |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 10 | 10 | 10 |  | 10 | 10 | 10 |  |
| 10 | S | 10 |  | 10 | 10 | 10 |  |
| 10 | 10 | 10 |  |  |  | G |  |
| 10 |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |

## Heuristics

- Need a very accurate heuristic
- Where can we get a heuristic?
- Run $\mathrm{A}^{*}$ from the goal to the start state to get h() value for many states


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|  | $8+2$ | $7+3$ | $6+4$ | $5+5$ | $4+6$ | $3+7$ |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- |
|  | $9+1$ | $8+2$ |  | $4+4$ | $3+5$ | $2+6$ |
|  | S |  |  | $3+7$ | $2+4$ | $1+5$ |
|  |  |  | $2+4$ | $1+5$ | G | $1+7$ |
|  |  |  |  | $3+5$ | $2+6$ | $1+7$ |
|  |  |  |  | $2+8$ |  |  |
| g+h |  |  |  | $4+6$ | $3+7$ | $2+8$ |

## Windowed Search

- We now have a perfect heuristic
- With a perfect heuristic only 1-step lookahead is needed
- Stop search at any time and be guaranteed to be on a path to the goal
■ Do $k$-step lookahead in cooperative space


## WHCA* ${ }^{*}(k)$

- Do single $A^{*}$ search from goal to start
- Do $k$-step forward cooperative search
- Expand original search if new heuristic values needed


## WHCA* Drawbacks

- First step is expensive
- Compute complete reverse A* search
- Compute forward CA* search
- Memory per unit is expensive
- Keep whole search frontier in memory
- Goal State can't change


## Improving WHCA*

- Abstraction
- Widely used idea (eg Holte, 1996)
- Two possible usages
- WHCA* $(w, a)$
- CPRA* $(k)$


## Abstraction

- Use fine-grained map abstraction
- Dragon Age abstraction abstractions 16x16 sectors in one step
- Instead abstract $2 \times 2$ sectors in one step
- Or: abstract small cliques (4 nodes) in the map
- Theoretical work suggests this minimizes pathfinding computation
- (Holte, 96; Sturtevant and Jansen, 07)




## WHCA* $(k, a)$

- Same as WHCA* $(k)$ but do reverse A* search at abstract level a
- Keep smaller $\mathrm{A}^{*}$ open/closed list in memory
- Faster A* computation
- Eventually less accurate


## PRA*

- Partial-Refinement $\mathrm{A}^{*}$
- Use multiple abstraction levels
- Refine abstract paths using A*


## Pathfinding

- Given abstract path:
- Path defines a corridor in the lower level of abstraction
- Run A* in this corridor to find next path
- Repeat until done

PRA $^{*}(\infty)$
PRA* ${ }^{*}(k)$


## CPRA* ${ }^{*}$ )

- Same as $\operatorname{PRA}^{*}(k)$, but do $\operatorname{WHCA}^{*}(k, 1)$ at last refinement level
- Only plan part of total path
- Much lower first-step cost
- Repeated WHCA* calls after executing each path


## Experiments

- Run algorithms on 256x256 map
- Place units on opposite sides of map and ask them to cross sides
- Report 95\%


Memory Usage


Nodes
First second

Nodes
Average per second


Nodes
Average per second


## Generalizing

- General technique for $n$-dimensional pathfinding problems
- Solve problem in $n$-1 dimensional space
- Use as heuristic in $n$-dimensional search
- If possible use "lower resolution" version of $n-1$ dimensional problem


## But...

- How well does it work with lots of units in open space?
- Not as well as one might expect



## But...

- How well does it work with lots of units in open space?
- Not as well as one might expect
- Units are searching for the shortest path
- Prefer shorter paths over paths which have a higher probability of success


Real crowd movement


Simulated crowd movement

## Some perspective

- Static 2-d search is cheaper than 3-d search
- Static information about other units isn't very useful
- Is there any other static information that we can retain?

Static information about motion.

## Retained Information

- Direction Vector
- Associated with a location on a map
- Which direction units travel through the location
- Updated dynamically as units move
- Direction Map
- Direction vectors for every location on the map
- Similar to flow fields used for flocking


Now what...?

- What do we do with all these arrows?
- During planning:
- Traveling in the same direction of an arrow is cheaper
- Traveling in the opposite direction is more expensive



Uncoordinated Unit Behavior


Coordinated Unit Behavior


Uncoordinated
Unit Behavior

Coordinated
Unit Behavior


Uncoordinated Unit Behavior

Coordinated Unit Behavior

## Other considerations

- Where do the initial weights come from?
- How much memory does it take to store the weights?
- What is the additional planning cost?



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## Planning Cost

- Planning using direction maps is more expensive
- Weighted $A^{*}$ can reduce the cost
- Use abstraction to reduce planning length
- Can maintain direction maps for classes of units, or only in congested areas of the map

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## Summary

- Abstraction is orthogonal to many other search enhancements
- Everything Sven talked about could be used on one or more levels of abstraction
- Rich toolbox for balancing performance in any particular domain


## Summary

- Abstraction techniques very effective across a variety of problems in reducing planning costs
- Used for defining subgoals in search
- Dragon Age
- Used for heuristics in search
- Cooperative pathfinding
- Many different ways of applying abstraction
- Best method depends on problem constraints


## Thanks!

- Comments, questions?
- Co-collaborators:
- Markus Enzenberger
- Renee Jansen
- Michael Buro
- Vadim Bulitko


