## Solving $8 \times 8$ Hex

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## Hex Rules and Properties



## Rules

- Two players alternate turns playing on any empty cell
- Stones are permanent (no moving or capturing)
- Goal is to connect your two sides of the board


## Hex Rules and Properties



## Properties

- Extra P-stones never disadvantageous for player P
- Draws are impossible
- First player wins: strategy-stealing argument
- Determining winner is PSPACE-complete


## Previously Solved States



- Last milestone for automated Hex solvers in 2004
- All $7 \times 7$ openings solved in two weeks (Hayward et al)
- By hand, humans have solved centre opening on $9 \times 9$ (Yang) and a few openings on $8 \times 8$ (Mishima et al, Yang)


## H-Search



- H-Search: algorithm that deduces existing connection strategies in a given Hex position (Anshelevich)
- Virtual connections (VC): 2nd-player connection strategy
- Semi-connections (SC): 1st-player connection strategy
- Carrier: empty cells required for a connection strategy


## Mustplay



- Identifying a winning VC terminates search
- Identifying winning SCs immediately prunes losing moves
- Mustplay: intersection of winning opponent SC carriers


## Inferior Cell Analysis



- Graph-theoretic properties and combinatorial game theory
- Fill-in: can add stones to the board without changing its win/loss value
- Reversible and dominated moves: can be pruned from consideration


## Opposite Color Bridges



- If a P-chain is adjacent to both $\bar{P}$ edges, then splits board into two independent regions
- Easy to detect these decompositions, but very rare
- Opposite-color bridges: can treat the two carrier cells as non-adjacent


## Split Decompositions



- Two chains touch if they are adjacent or form an opposite-color bridge
- Split decomposition: when a P-chain touches both $\bar{P}$-edges


## Four-Sided Decompositions



- Four-sided decomposition: a 4-cycle of touching Black and White chains
- If player P has a VC connecting the two P -chains of a four-sided decomposition, the region can be filled-in with P-stones


## Proof Set Pruning



- During search we identify previously-unknown winning SCs
- Can use discovered SCs to further reduce mustplay
- The smaller the SC carrier, the more moves can be pruned


## Proof Set Reduction



- Given a discovered SC, we try to shrink its carrier
- Cells outside the SC for player P can be assigned to $\bar{P}$
- Inferior cell analysis may identify $\bar{P}$-fill-in
- These cells can be deleted from the SC's carrier


## Proof Set Transpositions



- While solving states we track the winning strategy's carrier
- The losing player's stones can be any combination of cells outside of this carrier
- We can store the result for all these combinations as well


## Player Exchange Transpositions



- Want to translate a solved state to equivalent ones with players reversed
- Mirroring stones and reversing their colors is not adequate
- Stone must be added or removed; depends on player to move and who won


## Current Results



- $7 \times 7$ : 10 minutes
- $8 \times 8$ : 300 hours and $10^{8}$ internal nodes
- $9 \times 9$ : Cannot solve any opening in two weeks time


## Feature Contributions on $7 \times 7$

| feature $f$ | only $f$ off |  | only $f$ on |  |
| :---: | :---: | :---: | :---: | :---: |
|  | time | nodes | time | nodes |
| rotation/transposition deduction | 2.17 | 2.22 | 0.43 | 0.43 |
| decompositions | 1.29 | 1.51 | 0.68 | 0.61 |
| proof set reduction | 0.98 | 1.01 | 1.03 | 0.87 |

## Summary

- New: decompositions, proof set reduction, transposition deductions
- Enhanced H-search, inferior cell analysis
- First automated solver for $8 \times 8$ Hex openings

Future Work

- $9 \times 9$ at least 3 magnitudes more difficult
- Depth-first proof-number search (parallelized)
- Further improve inferior cell analysis, decompositions, etc


## Any Questions?

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