

bipartite preference and stable matchings

- see also https://en.wikipedia.org/wiki/Stable_marriage_problem

[C B D A] 0 • • A [3 2 0 1]

[A D B C] 1 • • B [2 3 1 0]

[A B C D] 2 • • C [2 0 1 3]

[D A C B] 3 • • D [2 0 1 3]

- *bipartite preference system*

- hospitals $H = \{0, 1, 2, 3\}$ residents $R = \{A, B, C, D\}$

- for each hospital, a *total preference order* of residents

e.g. hospital 0 prefers $C > B > D > A$

- for each resident, a **TPO** of doctors

e.g. resident A prefers $3 > 2 > 0 > 1$

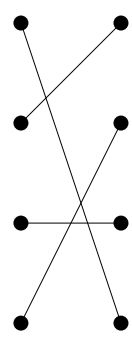
- a matching $m : H \leftrightarrow R$

[C B D A] 0 • • A [3 2 0 1]

[A D B C] 1 • • B [2 3 1 0]

[A B C D] 2 • • C [2 0 1 3]

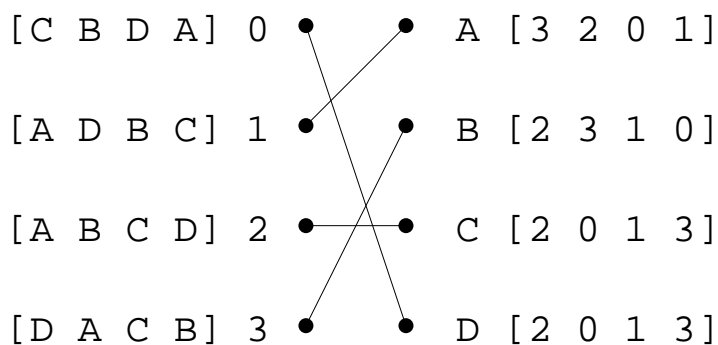
[D A C B] 3 • • D [2 0 1 3]



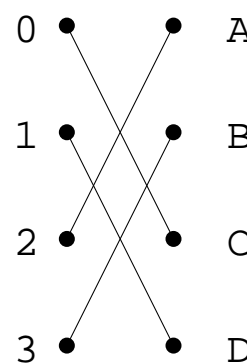
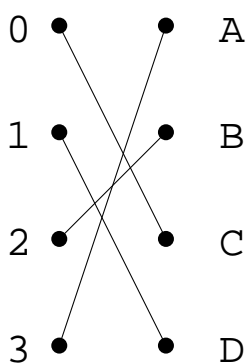
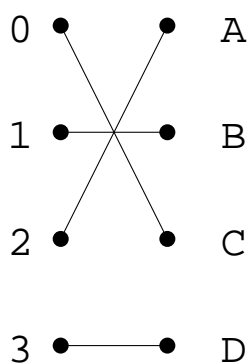
- in m , couple $\{2, A\}$ *unhappy*: prefer each other to their matches

2 prefers A to $m(2) = C$ A prefers 2 to $m^{-1}(A) = 1$

- this m *unstable*: it has an unhappy couple



- which matchings below are stable? (use above prefs)



- problem: find stable matching, if one exists
- brute force?
- for each possible matching (there are $n!$):
 - check whether matching stable in $O(n^2 \lg n)$ time (how?)
- worstcase runtime $O(n! n^2 \lg n)$
- can we do better?

Gayle-Shapley propose-maybe-reject alg'm

[C B D A] 0 • • A [3 2 0 1]

[A D B C] 1 • • B [2 3 1 0]

[A B C D] 2 • • C [2 0 1 3]

[D A C B] 3 • • D [2 0 1 3]

A B C D

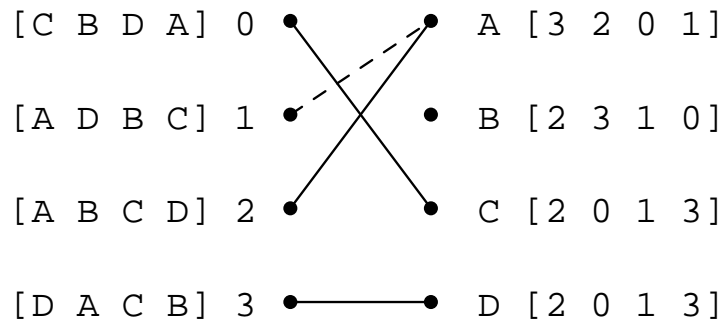
0

1

2

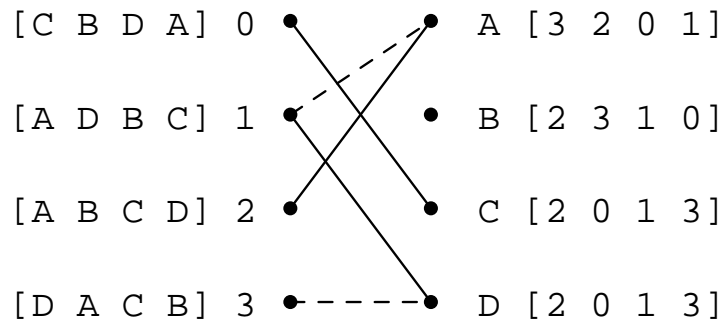
3

Gayle-Shapley: after round 1



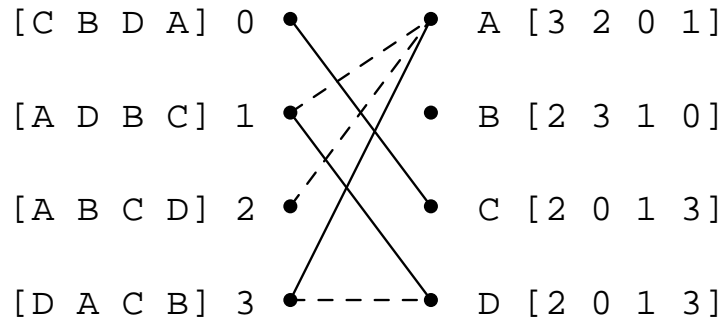
	A	B	C	D
0			?	
1	x			
2	?			
3				?

Gayle-Shapley: after round 2



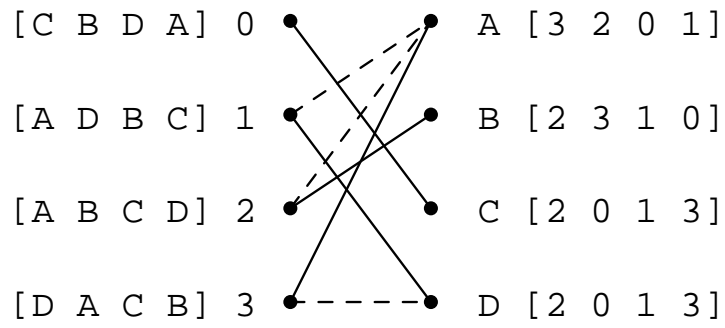
	A	B	C	D
0			?	
1	x			?
2	?			
3				x

Gayle-Shapley: after round 3



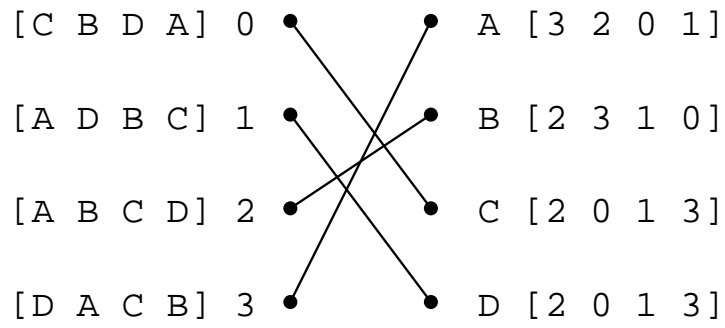
	A	B	C	D
0			?	
1	x			?
2	x			
3	?			x

Gayle-Shapley: after round 4



	A	B	C	D
0			?	
1	x			?
2	x	?		
3	?			x

Gayle-Shapley: matching



	A	B	C	D
0			?	
1				?
2		?		
3	?			