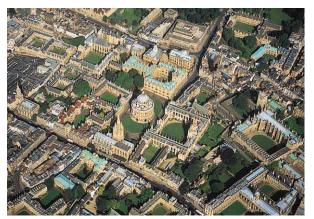
abc - the AspectBench Compiler for AspectJ



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Outline

- AspectJ introduction for compiler writers
- Challenges of building a compiler for AspectJ
- **abc** as an extensible and optimizing compiler
- How **abc** tackles performance issues
- Future Work



AspectJ Programming Language

- a seamless *aspect-oriented* extension to Java
- originally developed at Xerox PARC
- tools for AspectJ now developed and supported by the Eclipse AspectJ project
 - **ajc** compiler for the AspectJ language

(http://eclipse.org/aspectj)



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(http://eclipse.org/aspectj)

 abc, the AspectBench Compiler, is a new, alternative compiler for the AspectJ language, designed for *extensibility* and *optimization* (http://aspectbench.org)



AspectJ Introduction

- introduce a small Java program, a little expression interpreter
- illustrate three main uses of AspectJ by applying it to this small example
 - aspects for additional static checking at compile time
 - adding fields/classes/constructors to classes via aspects
 - dynamic aspects for applying advice (code) at specified run-time events



Example Java Program - expression interpreter

Consider a small interpreter for an expression language, consisting of:

- SableCC-generated files for scanner, parser and tree utilities in four packages: parser, lexer, node and analysis.
- main driver class, tiny/Main.java, which reads the input, invokes parser, evaluates resulting expression tree, prints input expression and result.
- expression evaluator class, tiny/Evaluator.java

```
> java tiny.Main
Type in a tiny exp followed by Ctrl-d :
3 + 4 * 6 - 7
The result of evaluating: 3 + 4 * 6 - 7
is: 20
```



AspectJ for Static (compile-time) Checking

- Programmer specifies a pattern describing a static program property to look for and a string with the warning text.
- An AspectJ compiler must check where the pattern matches in the program, and issue a compile-time warning (string) for each match.

```
public aspect StyleChecker {
   declare warning :
      set(!final !private * *) &&
      !withincode(void set*(..)) :
      "Recommend use of a set method.";
```



Using the StyleChecker aspect

The compilation:

```
abc StyleChecker.java */*.java
```

produces the compile-time output:

```
parser/TokenIndex.java:34:
Warning -- Recommend use of a set method.
index = 4;
^-----^
```



AspectJ for Intertype Declarations

- Programmer specifies, in a separate aspect, new fields/methods/constructors to be added to existing classes/interfaces.
- An AspectJ compiler must weave in code to implement these additions.
- Other classes in the application can use the added fields/members/constructors.
- In our example, we can use an aspect to add fields and accessors to the code generated by SableCC, without touching the generated classes.



Intertype Declarations - example

All AST nodes generated by SableCC are subclasses of **node.Node**. We must **not** directly modify the code generated by SableCC.

```
public aspect AddValue {
    int node.Node.value; // a new field
    public void node.Node.setValue(int v)
        { value = v; }
    public int node.Node.getValue()
        { return value; }
}
```



```
abc AddValue.java */*.java
```

where, the evaluator visitor can be now written using the value field to store intermediate values.

instead of the "old" way of storing intermediate values in a hash table. The aspect-oriented method is more efficient because fewer objects are created during the evaluation.



AspectJ for Dynamic Advice

- Programmer specifies a pattern describing run time events, and some extra code (advice) to execute before/after/around those events.
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- Programmer specifies a pattern describing run time events, and some extra code (advice) to execute before/after/around those events.
- An AspectJ Compiler must weave the advice into the base program for all potentially matching events.
- Since events can depend on dynamic information:
 - some execution state may need to be tracked, and
 - some advice may be conditional on the result of a dynamic residue test.



Dynamic Advice - counting runtime events

```
public aspect CountEvalAllocs {
  int allocs; // counter
  before () : call(* *.eval(..)) &&
              within(*.Main)
    \{ allocs = 0; \}
  after () : call(* *.eval(..)) &&
             within(*.Main)
    { System.out.println(
          "*** Eval allocs: " + allocs); }
  before () : call(*.new(..)) &&
              cflow(call(* *.eval(..)))
    \{ allocs ++; \}
```



Using the CountEvalAllocs aspect

 Using the interpreter with the CountEvalAllocs aspect included.

```
The result of evaluating:
3 + 4 * 6 + 9 / 3
*** Eval allocations: 17
is: 30
```

 Using the interpreter with the CountEvalAllocs aspect, and the improved evaluator enabled by the addvalue aspect.

```
The result of evaluating:
3 + 4 * 6 + 9 / 3
*** Eval allocations: 2
is: 30
```



Dynamic Advice - example 2

```
public aspect ExtraParens {
   String around() :
      execution(String node.AMultFactor.toString())
      execution(String node.ADivFactor.toString())
      { String normal = proceed();
      return "(" + normal + ")";
   }
}
```

```
Compile: abc ExtraParens.java */*.java
Run: java tiny.Main
The result of evaluating:
The result of evaluating:
3 + (4 * 6) + (9 / 3)
is: 30
```

Recap: uses of AspectJ for example

- Static (compile-time) check: Check that accessor methods are always used to set non-private non-final fields.
- Intertype declaration: Add a new field and associated accessor methods to the SableCC-generated node.Node class.
- Dynamic advice:
 - Count the number of allocations peformed during a an expression evaluation.
 - Intercept calls to tostring() for factors and add surrounding parentheses, if they are not already there.



Challenges: front-end

- AspectJ-specific language features, including relatively complex pointcut (patterns) language.
- Intertype declarations, need to be able to extend the type system in non-trivial ways.



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- AspectJ-specific language features, including relatively complex pointcut (patterns) language.
- Intertype declarations, need to be able to extend the type system in non-trivial ways.
- abc's solution:
 - use Polyglot, an extensible framework for Java compilers (Cornell)
 - express AspectJ language via LALR(1) grammar: base Java grammar + additional grammar rules for AspectJ
 - use Polyglot's extension mechanisms to override key points in type system to handle intertype declarations.



Challenges: back-end

- Need to handle input from .java and .class files.
- AspectJ compilers need additional modules: matcher, weaver
- need to produce **efficient** woven code (.class files)



Challenges: back-end

- Need to handle input from .java and .class files.
- AspectJ compilers need additional modules: matcher, weaver
- need to produce **efficient** woven code (.class files)
- abc's solution:
 - clean design of matcher and weaver using a simplified and factored pointcut language
 - use Soot, which provides Jimple IR (typed 3-addr), standard optimizations, and an optimization framework



The abc approach

abc has been designed to be an:

- extensible compiler:
 - easy to implement language extensions
 - build on two extensible frameworks, Polyglot and Soot
 - see AOSD 2005 submission at http://aspectbench.org/techreports



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optimizing compiler:

- convenient IR
- good weaving strategies
- standard compiler optimizations
- AspectJ-specific optimizations



Does the weaving strategy matter?

- Studied the code produced by **ajc** by tagging instructions that are introduced by the **ajc** weaver and using *J tool to measure dynamic metrics. (OOPSLA 2004)
- When there is **not** a lot of overhead:
 - very simple before and after advice
 - when the aspect only applies to a small, cold, part of the program
 - when the aspect body is a large computation
- When there can be overhead:
 - frequent (hot) aspects with small bodies
 - frequent (hot) use of cflow and/or around advice



How abc reduces overhead

- use Soot in back-end, so can optimize generated code
- new around weaving strategy
- new cflow implementation



Reducing overhead by using Soot

- the abc backend uses Jimple, a typed 3-address IR
 (ajc use stack-based Java bytecode)
 - abc weaver does not need to save implicit values on the stack, leads to fewer locals in generated code
 - abc weaver can use def-use and variable types to generate better code
- abc uses the Soot basic optimizations to clean up generated code
- abc can use Soots intra- and inter-procedural analysis frameworks to implement AspectJ-specific optimizations.



Weaving in bytecode (ajc)

	c int foo(int x,	int y, int z)
0:	aload_0	
1:	iload_1	
2:	iload_2	
3:	iload_3	
4:	istore	% 4
6:	istore	% 5
8:	istore	% <mark>6</mark>
10:	astore	%7
12:	invokestatic	A.aspectOf ()LA; (52)
15:	aload	87
17:	invokevirtual	A.ajc\$before\$A\$124 (LFoo;)
20:	aload	%7
22:	iload	% <mark>6</mark>
24:	iload	% 5
26:	iload	% 4
28:	invokevirtual	Foo.bar (III)I (37)
31:	ireturn	

Weaving in Jimple (abc)

```
public int foo(int, int, int)
{ Foo this;
   int x, y, z, $i0;
   A theAspect;
   this := @this;
   x := @parameter0;
   y := @parameter1;
   z := @parameter2;
   theAspect = A.aspectOf();
   theAspect.before$0(this);
   i0 = this.bar(x, y, z);
   return $i0;
```



Sascha's strategy for around weaving

- ajc has two strategies, inlining around advice, and using closures
 - the inlining method will work well for small advice bodies, or advice the applies in few places
 - the closure strategy is very inefficient, but must be used in some situations (i.e. when an advice applies to itself)
- abc has another strategy (http://aspectbench.org/theses):
 - doesn't inline (no code bloat), but uses generic advice methods
 - replaces polymorphism with lookup tables
 - avoids object creation
 - no closures in the general case
 - uses closures only at specific points, degrades gracefully



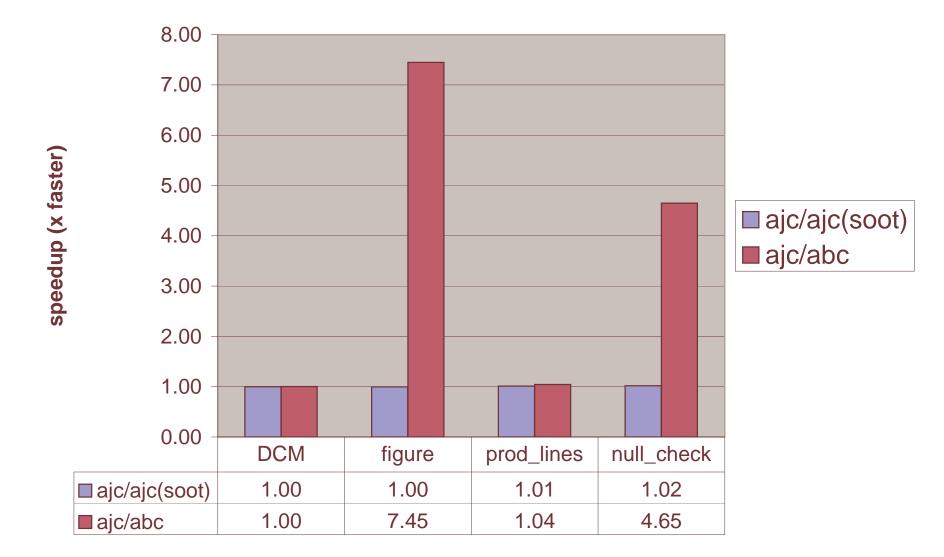
Improving the implementation of cflow

- to track if a runtime computation is within the cflow of some event, the compiler has to generate code to track when that event begins and when it ends
- in general the event may have some state, but most often it does not
- **ajc** uses a stack of states that must be thread-safe
- **abc** improves upon this by:
 - recognizing when there is no state and using a counter instead of a stack of empty states
 - recognizing when counters (or stacks) are equivalent and can be shared
 - only peforming thread-specific operations once per method body
 - will soon use interprocedural analysis to determine if the cflow can be decided statically (and thus no runtime book-keeping in necessary)



Peformance Improvement(1)

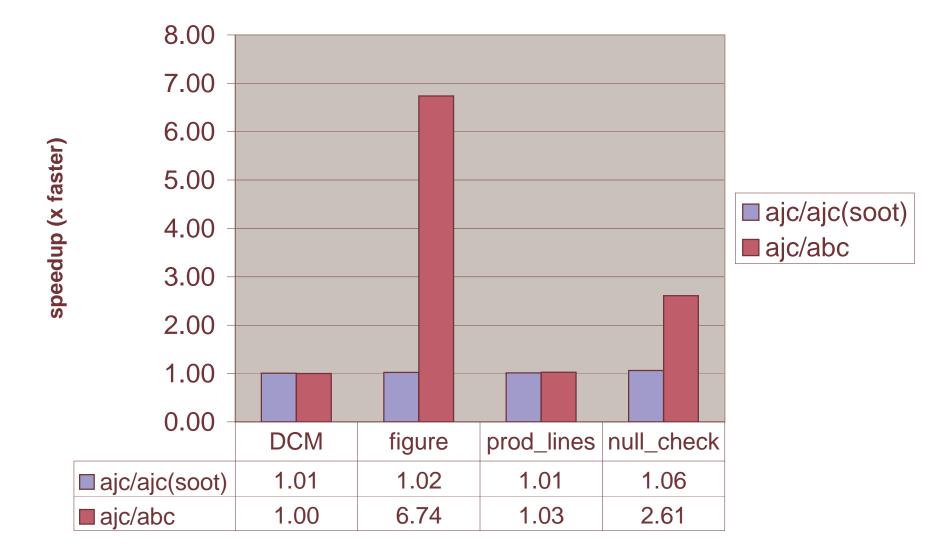
speedup (client JIT)





Peformance Improvement(2)

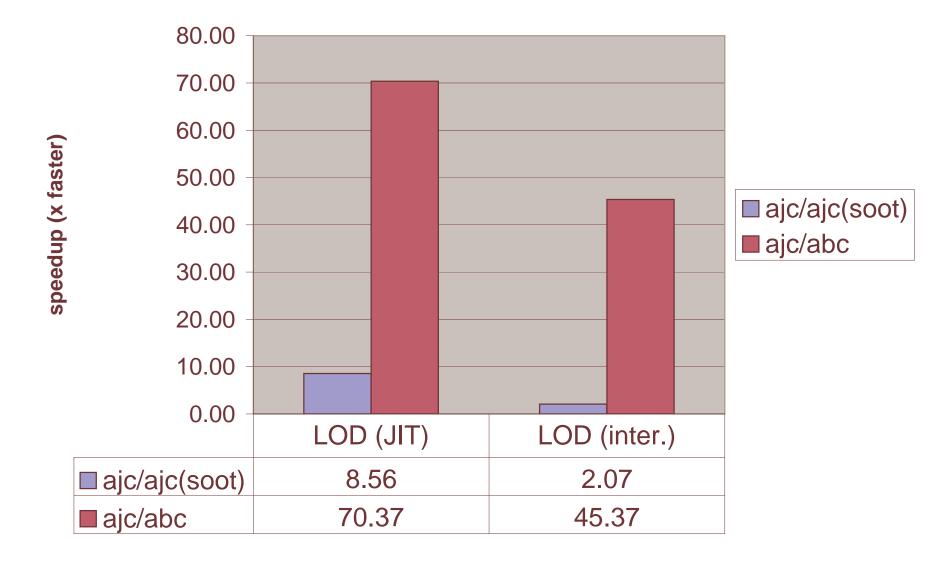
Speedup (Interpreter)





Peformance Improvement(3)

Killer Benchmark (Law of Demeter)





Conclusions

- AspectJ is very useful for many tasks illustrated with tiny interpreter
- **abc** is a new compiler for AspectJ which is extensible and optimizing.
- You can use abc as an alternative AspectJ compiler, or you can use it for research into language extensions and new optimizations.
- It is worth thinking about AspectJ-specific optimizations, and abc has already implemented some of these.
- Lots more work by the abc team to come ... we welcome users!



http://aspectbench.org