

MCGS: A Minimax-Based Combinatorial Game Solver

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Motivation

- MCGS: A framework for solving (sums of) *short* 2 player games
 - “Short” games are finite and loop-free
 - Given a sum of games and first player, find the winner
 - “Normal play” convention: last player to move wins
- Existing CGT tools: CGSuite
 - Lots of features
 - Based on canonical forms. Very slow!
- Example: empty 1x16 NoGo board
 - CGSuite: finds canonical form in almost 8 minutes
 - MCGS: finds outcome class in ~1.6 seconds
 - Without transposition table or database enabled!

Motivation (Continued)

> G1



> G2



> G3



> G4



> G5



> G6



*

$+-\{*,^{\wedge}\}$

0

$+-\{^{\wedge}[2],\{^{\wedge*},^{\wedge\wedge}|0,^{\wedge*},+-\{0,^{\wedge*}\}\}\}$

$+-\{^{\wedge}<2>*,^{\wedge},\{0|^{\wedge},+-\{*,^{\wedge}\}\},\{^{\wedge\wedge*}|v,+-\{*,^{\wedge}\}\}\}$

$+-\{0,^{\wedge*}\}$

Motivation (Continued)

> G7



$$\begin{aligned}
 &+ - \{ \{ \wedge [2]^*, \wedge^*, \{ 0 | \{ 0 | \wedge^*, + - \{ 0, \wedge^* \} \}, \{ \wedge^*, \wedge^* | 0, \wedge^*, + - \{ 0, \wedge^* \} \} \}, \{ \wedge^* 3^* | 0, \{ \wedge^*, \wedge^* | 0, \wedge^*, + - \{ 0, \wedge^* \} \} \} | 0, \{ \wedge^* | 0, \wedge^* \}, \{ 0, \{ \wedge^*, \wedge^* | 0, \wedge^*, + - \{ 0, \wedge^* \} \} | v^*, + - \{ 0, \wedge^* \} | 0 \}, + - \{ \wedge < 2 >, \wedge^*, \{ 0 | \wedge^*, + - \{ 0, \wedge^* \} \}, \{ \wedge^* | v^*, + - \{ 0, \wedge^* \} \} \}, \{ \wedge^*, \{ \wedge^*, \wedge^* | 0, \wedge^*, + - \{ 0, \wedge^* \} \} | 0, v^* \} \}, \{ \{ \wedge^* | \wedge^* \}, \{ \wedge^* 3^* | 0, \{ \wedge^*, \wedge^* | 0, \wedge^*, + - \{ 0, \wedge^* \} \} \} | 0, *, \{ 0, *, \{ \wedge^*, \wedge^* | 0, v^* \} | v v, v^* \} \}, \{ 0 | \{ 0 | \wedge^*, + - \{ 0, \wedge^* \} \}, \{ \wedge^*, \wedge^* | 0, \wedge^*, + - \{ 0, \wedge^* \} \} | + - \{ \wedge [2], \{ \wedge^*, \wedge^* | 0, \wedge^*, + - \{ 0, \wedge^* \} \} \}, \{ \wedge^*, \wedge^* | *, \wedge^*, + - \{ *, \wedge^* \} | \{ 0, v^*, + - \{ 0, \wedge^* \} | v v, v^* \}, + - \{ 0, \wedge^* \} \} \}, \{ 0 | \{ 0 | \wedge^*, + - \{ 0, \wedge^* \} \}, \{ \wedge^*, \wedge^* | 0, \wedge^*, + - \{ 0, \wedge^* \} \} | \{ 0 | \wedge^*, + - \{ *, \wedge^* \} | *, \{ *, v, + - \{ *, \wedge^* \} | v v^*, v \} \}, \{ 0 | \wedge^*, + - \{ 0, \wedge^* \} | 0, \{ 0, v^*, + - \{ 0, \wedge^* \} | v v, v^* \} \}, \{ + - \{ 0, \wedge^* \}, \{ \wedge^*, \wedge^* | 0, \wedge^*, + - \{ 0, \wedge^* \} \} | *, v, + - \{ *, \wedge^* \} | v v^*, v \} \}, \{ \{ 0 | + - \{ 0, \wedge^* \} \}, \{ \wedge^*, \{ \wedge^* | \wedge^* \} | \wedge^*, \{ \wedge^*, \wedge^* | 0, *, \{ 0, \wedge^* | v, v^* \} \}, \{ \wedge^*, \wedge^* | 0, *, \{ 0, \wedge^* | v, v^* \} \} | + - \{ 0, \wedge^* \}, \{ *, \wedge^* < 2 >, \{ \wedge^*, \wedge^* | 0, *, \{ 0, \wedge^* | v, v^* \} | v, v^*, \{ 0, v^*, \{ 0, \wedge^* | v, v^* \} | v v, v v^* \} \}, + - \{ *, \wedge^* \}, \{ + - \{ 0, \wedge^* \}, \{ \wedge^*, \wedge^* | 0, \wedge^*, + - \{ 0, \wedge^* \} \} | *, v, + - \{ *, \wedge^* \} | v v^*, v \} \} \}
 \end{aligned}$$

> G8



$$\begin{aligned}
 &+ - \{ \{ \wedge^*, \wedge^* | *, \wedge^*, + - \{ *, \wedge^* \} \}, \{ \wedge^* | + - \{ \wedge < 2 >, \wedge^*, \{ 0 | \wedge^*, + - \{ *, \wedge^* \} \}, \{ \wedge^* | v, + - \{ *, \wedge^* \} \}, \{ \wedge^* [2]^*, \{ 0 | \wedge^*, + - \{ *, \wedge^* \} \}, \{ \wedge^*, \wedge^* | *, \wedge^*, \{ \wedge^*, \wedge^* | *, v \} \} | *, \{ *, v, + - \{ *, \wedge^* \} | v v^*, v \} \} \}
 \end{aligned}$$

69



> G10


$$* , v \} | \{ \{ 0, v^*, + - \{ 0, \wedge^* \} | v v, v^* \}, \{ v^*, + - \{ 0, \wedge^* \} | 0 \} | 0 \} \} \}$$

MCGS

- Goal: create an efficient, general framework for solving games
 - Apply lessons learned from previous specialized solvers
 - Linear NoGo and NoGo (Du et al. 2023, 2024)
 - Linear Clobber (Tavakoli et al. 2022, Folkersen et al. 2022, 2025)
 - Implement rules of a game, benefit from MCGS's search engine
 - Easier than writing a game-specific solver from scratch
 - C++ source code
- Support for both *partizan* and *impartial* games
 - Partizan games: boolean minimax search to find winner
 - Impartial games: specialized algorithms to find nim-value
 - Wrapper creates impartial versions of any implemented game

MCGS (Continued)

- Test framework
 - “.test” input files: specify positions as sums of subgames, with first players and expected results
 - .csv/.html output
 - Search results
 - Search statistics: time, search node count, etc.
 - Thousands of test cases
 - Many verified by external solvers (i.e. CGSuite, game-specific solvers)
- Documentation
 - For general users
 - For developers
 - Instructions for adding new games
 - Comprehensive discussion of implementation details

Supported Games

- “Basic” CGT games:

- Integer
- Dyadic Rational
- Nimber
- Switch
- Up/Star

* Denotes games coming in future releases

- Strip (1D) games:

- Linear Clobber
- Linear NoGo
- Elephants and Rhinos
- * Generalized Toads and Frogs
- * Toppling Dominoes

- Grid (2D) games:

- Clobber
- NoGo
- * Amazons
- * Fission
- * Domineering
- * Battle Sheep

- Others:

- Kayles

Hashing

- 64 bit Zobrist hashes
 - Used to index into transposition table and endgame database
- Two-part hashing scheme
 - Local hash h (one subgame)
 - Global hash H (entire sum)
 - Computed from local hashes (next slide)
- Zobrist table Z
 - Table of random 64 bit integers
 - Indexed by pair: (location, color)
 - Specialized solver may only represent 3 colors (Clobber, NoGo)
 - MCGS's Zobrist table represents all 128 bit color values

$h(\text{Clobber:} \blacksquare \square _ \blacksquare) := Z(0, \blacksquare) \oplus Z(1, \square) \oplus Z(2, _) \oplus Z(3, \blacksquare) \oplus \text{gametype}(\text{Clobber})$

Global Hash

- Compute global hash $H(S, P)$ for sum S and player P :
 - a. Sort games of sum: $S = g_0 + g_1 \dots + g_n$
 - Sums differing only in order should have the same global hash
 - $H(g_0 + g_1, P) = H(g_1 + g_0, P)$
 - b. Compute list of local hashes: $[h(g_0), h(g_1), \dots, h(g_n)] = [h_0, h_1, \dots, h_n]$
 - c. $H(S, P) := Z_sum(0, h_0) \oplus Z_sum(1, h_1) \oplus \dots \oplus Z_sum(n, h_n) \oplus player(P)$
 - Transform local hashes using another Zobrist table Z_sum
 - Two-level re-coding scheme prevents contents of local hashes cancelling out
 - Yields different hashes even for sums with thousands of copies of one game

Game-specific Functions

- Abstract base type *game*
- Required functions (each new game must implement):
 - Play and undo move
 - Compute local hash
 - Print (produces textual representation of the game)
 - Create move generator
 - Move generators produce all moves for a subgame

Game-specific Functions (Continued)

- Optional functions:
 - Split
 - Break game into subgames
 - Clobber: $\blacksquare \square _ \blacksquare \blacksquare \square = \text{Clobber:} \blacksquare \square + \text{Clobber:} \blacksquare \blacksquare \square$
 - Clobber: $\blacksquare \square _ \blacksquare \blacksquare \square = * + \uparrow$
 - Takes advantage of “basic” CGT game optimizations
 - Normalize
 - Convert equivalent games to one common, unique representation
 - Clobber: $\blacksquare \blacksquare \square _ _ = \text{Clobber:} \square \blacksquare \blacksquare$
 - Split and normalize functions increase transposition table hits
- Optional incremental hash update
 - Update local hash when changing a subgame's state
 - Avoids recomputing local hash in its entirety

Database (DB)

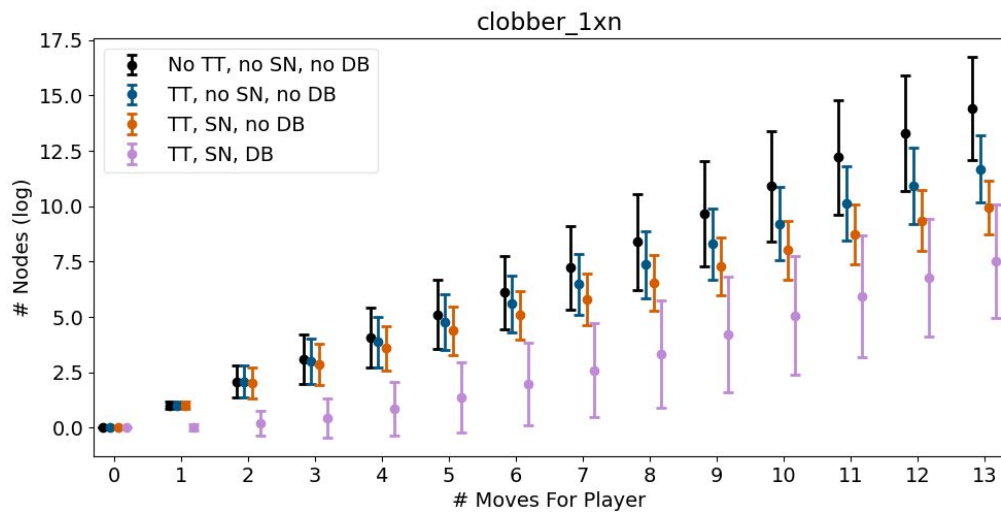
- Stores pre-computed information about subgames
 - DB entries contain outcome classes
 - Assists with sum simplification, and earlier search termination
 - Currently: contains single subgames, which are normalized and don't split
 - Reduces size of DB
 - No sums yet
 - Queried using local hashes
- DB creation uses abstract base type *database game generator*
 - Generates subgames of some game type, for the DB to create their entries
 - Generates in order of increasing “size”
 - Ensures DB creation is fast: after 1 move on generated game, resulting subgames are in DB
 - Modular design with reusable components
 - Easy to add database support for a game type

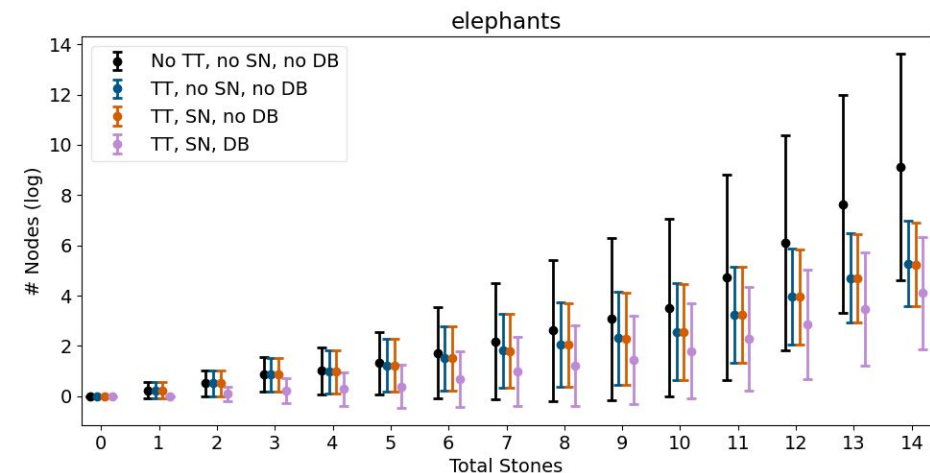
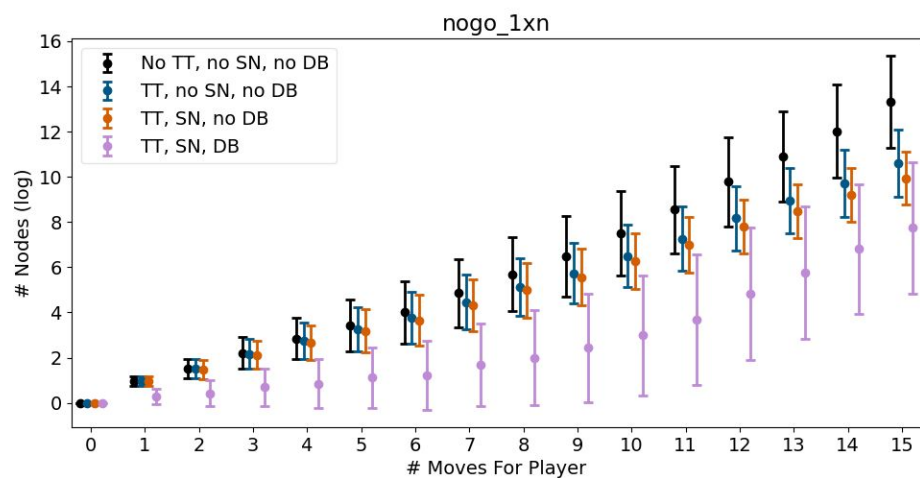
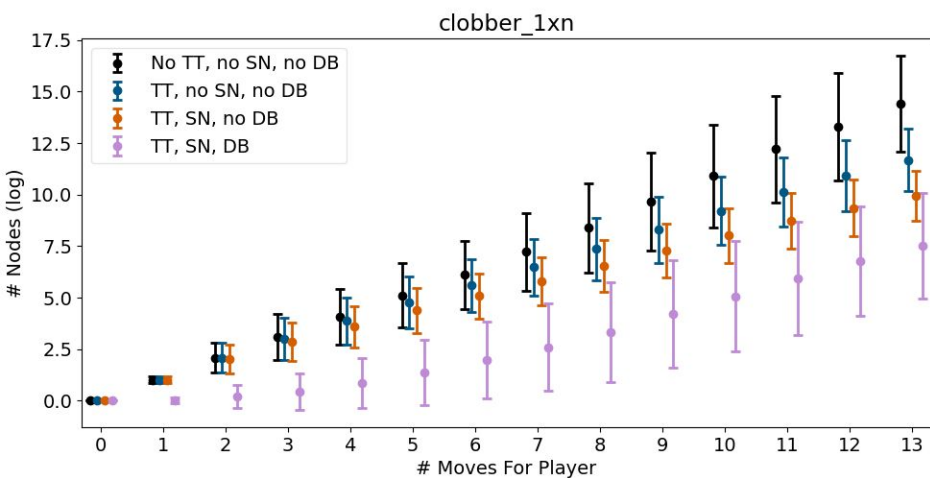
Sum Optimizations

- Basic CGT game simplification
 - Switches: $\{A \mid B\}$ for rationals A and B
 - Extract mean: $\{4 \mid -8\} = -2 + \{6 \mid -6\}$
 - Convert to number: $\{2 \mid 3\} = 5/2$
 - Sum up integers and rationals:
 - $1 + 2 + 1/2 = 7/2$
 - Nim-add numbers:
 - $*3 + *2 = *1$
 - Sum together multiples of $\uparrow$ and $\downarrow$:
 - $2\uparrow + 3\downarrow = \downarrow$
- Look up subgames in database:
 - Prune games with outcome class P
 - Sometimes: outcome classes determine result without further search
- After playing a move on a subgame, use its split and normalize functions
 - Increase transposition table hits

Experiments

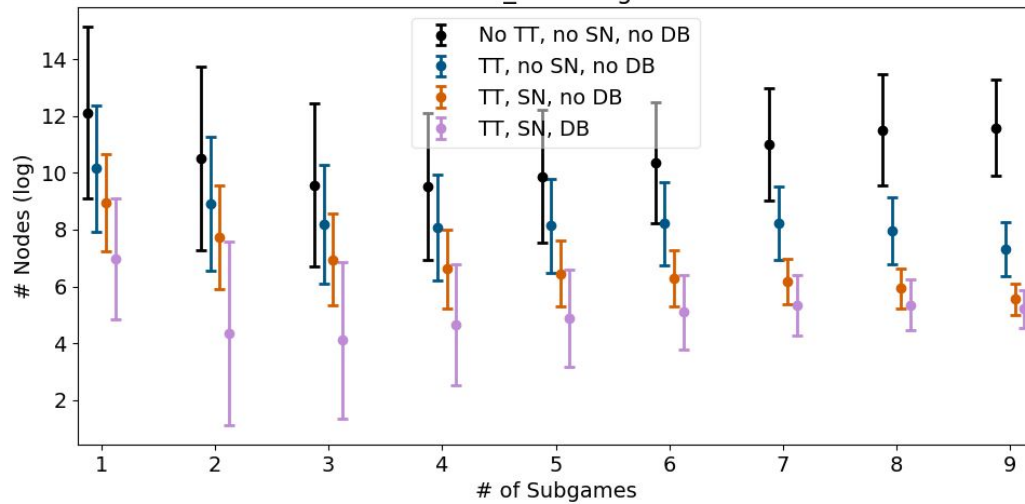
- Scaling results for several game types
 - Generate random games, sort into buckets (of up to 2000 games) along X axis
 - Y axis: (natural log of) search node counts
 - Node count: Invocations of recursive minimax search function
 - Same games across 4 ablations, changing 3 variables:
 - TT (transposition table)
 - SN (split and normalize)
 - DB (database)
 - 60 second timeout
 - A game which times out for any ablation setting is excluded from the data





- Games:
 - Linear Clobber, Linear NoGo
 - X axis: Move count for first player
 - Elephants and Rhinos
 - X axis: Stone count at start
- Performance scales best with DB
 - Effect diminishes along X axis. Due to DB only storing outcome classes?

clobber_1xn subgames



- Linear Clobber
 - X axis: Number of subgames at start
 - Again, benefit of DB diminishes with increasing number of subgames

Conclusions/Future Work

- MCGS fills the gap between:
 - CGT-based software relying on canonical forms
 - Solvers based on monolithic states
- Generic, pre-implemented optimizations are useful for a variety of games
 - Transposition table
 - Subgame structure
 - Database
- Future work:
 - Database
 - Support storing sums
 - Add more fields
 - Thermographs
 - Bounds
 - Simpler equal games (for substitution during search)
 - Dominated moves
 - More games
 - Move ordering heuristics
 - And much more!

Thanks For Listening!

- <https://github.com/ualberta-mueller-group/MCGS>
 - MIT license

* Incremental Hash Update

- Functions: Play, undo move, normalize, undo normalize

- Play example: Clobber: ■□■□ → Clobber: _■■□

- $h(\text{Clobber: } _■■□) = h(\text{Clobber: } ■□■□) \oplus \underbrace{Z(0, ■)}_{\text{Update location 0}} \oplus \underbrace{Z(0, _)}_{\text{Update location 1}} \oplus Z(1, □) \oplus Z(1, ■)$

Update location 0 Update location 1

* Database Game Generator

- Generates in order of increasing “size”
 - Ensures DB creation is fast: after 1 move on generated game, resulting subgames are in DB
 - Ordered first by grid dimensions, then by game-dependent order:
 - Grid dimension ordering: $1 \times 1 \rightarrow 1 \times 2 \rightarrow 2 \times 1 \rightarrow 2 \times 2 \rightarrow \dots \rightarrow 3 \times 3$
 - For given dimensions, Clobber generates in order of increasing stone count:
 - 1x3: $___ \rightarrow \dots \rightarrow \blacksquare \square \blacksquare \rightarrow \dots$
 - Similarly, NoGo generates in order of decreasing stone count:
 - 1x3: $\blacksquare \blacksquare \blacksquare \rightarrow \dots \rightarrow ___$