

# Lecture 31: NP-Completeness

## Agenda:

- $NP$ -completeness: the main ideas
- Graph representations & size of an instance
- Decision problems & instances
- Polynomial time
- Classes  $P$ ,  $NP$

## Reading:

- Textbook pages 966 – 983

## Lecture 31: NP-Completeness

The problems we have studied:

1. Sorting can be done in  $\Theta(n \log n)$
2. Longest common subsequence can be done in  $\Theta(n \times m)$
3. Minimum spanning tree can be done in  $\Theta(m \log n)$
4. Single-source shortest paths can be done in  $\Theta(m \log n)$

They can be solved in a short amount of time.

What does “short” mean?

1. Sorting  
 $n$  is the number of keys, measuring how big the sorting instance is
2. Longest common subsequence  
 $n, m$  are the lengths of the sequences, measuring how big the sorting instance is
3. Minimum spanning tree  
 $n, m$  are the numbers of vertices and the number of edges, measuring how big the sorting instance is
4. Single-source shortest paths  
 $n, m$  are the numbers of vertices and the number of edges, measuring how big the sorting instance is

“Short” is polynomial in the “size” of the instance

Some other computational problems:

- Eulerian tour — a cycle including every edge exactly once

Determine if a graph is Eulerian

can be done in  $\Theta(n^2)$

- Hamiltonian cycle — a cycle including every vertex exactly once

Determine if a graph is Hamiltonian

so far no known algorithm in  $O(n^k)$  for any  $k$

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- Shortest  $x$ -to- $y$  path

can be done in  $\Theta(n^2)$

- Longest  $x$ -to- $y$  path

so far no known algorithm in  $O(n^k)$  for any  $k$

We need/want to classify the problems into “**easy**” and “**hard**” categories ...

## Basic concepts:

- Size of an instance: in order to store the instance into the computer, how many memory units are necessary?

Adjacency list representation of a graph containing 6 vertices and 7 edges:

```
1: 2, 3, 5
2: 1, 4, 6, 3
3: 2, 1
4: 5, 2
5: 4, 1
6: 2
```

6//2,3,5/1,4,6,3/2,1/5,2/4,1/2// — 32 memory units

In this problem:  $\Theta(n + m)$

- Polynomial time — polynomial in the size(s) of the instance(s)
- Abstract problem
  - two parts:
    1. a set of instances — which are inputs
    2. a query — the question asked
  - instance solution: answer to the query — the output

Example: minimum spanning tree problem

1. instances: all edge-weighted (simple, undirected) graphs
2. query: for input  $G$ , what is the length of an MST of  $G$ ?

## Basic concepts (2):

- Decision problem: the answer to the query is **yes** or **no**

Example: minimum spanning tree problem

1. instances: all  $(G, \ell)$ :  $G$  an edge-weighted (simple, undirected) graph,  $\ell$  an integer
  2. query: for input  $(G, \ell)$ , is there a spanning tree of  $G$  of length at most  $\ell$ ?
- Optimization problem: abstract problem with an optimization goal
  - Relations between optimization problem and its decision version
    - suppose you can solve the optimization problem, then you can solve the decision problem (how?)
    - suppose you can solve the decision problem, then generally you can solve the optimization problem as well (how?)
  - Notes:
    - for some abstract problems, their decision version is the same. Example: Hamiltonian graph problem
      - \* instances: all (simple, undirected) graphs  $G$
      - \* query: for input  $G$ , does it have a Hamiltonian cycle?
    - correspondence: optimization  $\leftrightarrow$  decision
      1. minimization  $\leftrightarrow$  *at most*
      2. maximization  $\leftrightarrow$  *at least*

The classes of  $P$  and  $NP$ :

- Class  $P$ 
  - decision problem
  - there exists some algorithm solving the problem in polynomial time
  - which of the above problems are in  $P$ :
    1. sorting
    2. longest common subsequence
    3. minimum spanning tree
    4. single-source shortest paths
    5. shorest  $x$ -to- $y$  path
    6. determining Eulerian graphs
  
- Class  $NP$ 
  - decision problem
  - for every **yes**-instance, there exists a proof that the answer is yes and the proof can be verified in polynomial time
  - which of the above problems are in  $NP$ :
    1. sorting
    2. longest common subsequence
    3. minimum spanning tree
    4. single-source shortest paths
    5. shorest  $x$ -to- $y$  path
    6. determining Eulerian graphs
    7. longest  $x$ -to- $y$  path
    8. determining Hamiltonian graphs

## Lecture 31: NP-Completeness

Have you understood the lecture contents?

well	ok	not-at-all	topic
<input type="checkbox"/>	<input type="checkbox"/>	<input type="checkbox"/>	rough idea on 'easy' and 'hard'
<input type="checkbox"/>	<input type="checkbox"/>	<input type="checkbox"/>	size of an instance
<input type="checkbox"/>	<input type="checkbox"/>	<input type="checkbox"/>	polynomial time
<input type="checkbox"/>	<input type="checkbox"/>	<input type="checkbox"/>	abstract, optimization, decision problems
<input type="checkbox"/>	<input type="checkbox"/>	<input type="checkbox"/>	optimization $\leftrightarrow$ decision
<input type="checkbox"/>	<input type="checkbox"/>	<input type="checkbox"/>	$P$ and $NP$