

## A Proofs

**Proof of Theorem 1:** As the set  $\mathcal{BN}_{\Theta \geq \gamma}(G)$  is uncountably infinite, we cannot simply apply the standard techniques for PAC-learning a finite hypothesis set. We can, however, partition this uncountable space into a finite number  $L = L(K, \gamma, \epsilon)$  of sets, such that any two BNs within a partition have similar conditional log-likelihood scores. We can then, in essence, simultaneously estimate the scores of all members of  $\mathcal{BN}_{\Theta \geq \gamma}(G)$  if we collect enough query instances to estimate the score for one representative of each partition.

Now for the details: We prove below that, if the CPTables for two BNs  $\Theta^{(1)}, \Theta^{(2)} \in \mathcal{BN}_{\Theta \geq \gamma}(G)$  have similar CPTables  $\Theta^{(1)} = \{\theta_{d_i|\mathbf{f}_i}^{(1)}\}_i$  and  $\Theta^{(2)} = \{\theta_{d_i|\mathbf{f}_i}^{(2)}\}_i$ , then they will have similar LCL-scores wrt any query; *i.e.*,

$$\text{if } \left| \theta_{d_i|\mathbf{f}_i}^{(1)} - \theta_{d_i|\mathbf{f}_i}^{(2)} \right| \leq \frac{\gamma \epsilon}{6K} \text{ then } \forall c, \mathbf{e} \left| \ln(P_{\Theta^{(1)}}(c|\mathbf{e})) - \ln(P_{\Theta^{(2)}}(c|\mathbf{e})) \right| \leq \frac{\epsilon}{6}. \quad (1)$$

This of course implies the same bound on the difference between their overall LCL-scores

$$|\text{LCL}_k(\Theta^{(1)}) - \text{LCL}_k(\Theta^{(2)})| \leq \frac{\epsilon}{6}$$

for any distribution  $\text{LCL}_k(\cdot)$  — both for the “true” query distribution  $\text{LCL}(\cdot)$ , and for the distribution associated with any empirical sample  $\widehat{\text{LCL}}(\cdot)$ .

We therefore partition the  $\mathcal{BN}_{\Theta \geq \gamma}(G)$  space into  $L = \left(\frac{6K}{\gamma \epsilon}\right)^K$  disjoint sets (where any two BNs from any partition will have similar CPTable values), then define the set  $R = \{\Theta_i\}_i$  to contain one representative from each partition. We prove below that a sample  $S$  of size

$$M\left(\frac{\epsilon}{6}, \frac{\delta}{L}\right) = 2 \left(\frac{3N \log \gamma}{\epsilon}\right)^2 \ln \frac{2L}{\delta} \quad (2)$$

is sufficient to estimate each of these single representatives to within  $\epsilon/6$  of correct, with probability of error at most  $\delta/L$ ; *i.e.*, such that, for each  $i$ ,

$$P \left[ \left| \widehat{\text{LCL}}^{(S)}(\Theta_i) - \text{LCL}(B_i) \right| > \frac{\epsilon}{6} \right] < \frac{\delta}{L}.$$

As there are  $L$  representatives, we have a total probability of at most  $L \frac{\delta}{L} = \delta$  that *any* of the representative’s scores are mis-estimated by more than  $\epsilon/6$ .

This means we have, in effect, estimated the scores on *any*  $\Theta \in \mathcal{BN}_{\Theta \geq \gamma}(G)$  to within  $\epsilon/2$ : For any  $\Theta \in \mathcal{BN}_{\Theta \geq \gamma}(G)$ , let  $\Theta' \in R$  be the representative in  $\Theta$ ’s partition. Observe

$$\begin{aligned} |\widehat{\text{LCL}}(\Theta) - \text{LCL}(\Theta)| &\leq |\widehat{\text{LCL}}(\Theta) - \widehat{\text{LCL}}(\Theta')| + |\widehat{\text{LCL}}(\Theta') - \text{LCL}(\Theta')| + |\text{LCL}(\Theta') - \text{LCL}(\Theta)| \\ &\leq \frac{\epsilon}{6} + \frac{\epsilon}{6} + \frac{\epsilon}{6} \\ &= \frac{\epsilon}{2}. \end{aligned}$$

This means, in particular, that our estimate of the scores of both  $\widehat{\Theta}$  and  $\Theta^*$  are within  $\epsilon/2$ , and so

$$\begin{aligned} \text{LCL}(\widehat{\Theta}) - \text{LCL}(\Theta^*) &\leq |\text{LCL}(\widehat{\Theta}) - \widehat{\text{LCL}}(\widehat{\Theta})| + \widehat{\text{LCL}}(\widehat{\Theta}) - \widehat{\text{LCL}}(\Theta^*) + |\widehat{\text{LCL}}(\Theta^*) - \text{LCL}(\Theta^*)| \\ &\leq \frac{\epsilon}{2} + 0 + \frac{\epsilon}{2} \end{aligned}$$

To complete the proof, we need only prove Equations 1 and 2. For Equation 1: Consider the sequence of BNs  $\Theta_0, \Theta_1, \dots, \Theta_K$  where the first  $i$  of  $\Theta_i$ ’s CPTables come from  $\Theta^{(1)}$ , and the remaining from  $\Theta^{(2)}$  — *i.e.*,

$$\Theta_i \sim \{\theta_{d_1|\mathbf{f}_1}^{(1)}, \dots, \theta_{d_i|\mathbf{f}_i}^{(1)}, \theta_{d_{i+1}|\mathbf{f}_{i+1}}^{(2)}, \dots, \theta_{d_K|\mathbf{f}_K}^{(2)}\}.$$

Now observe

$$\left| \ln(P_{\Theta^{(1)}}(c|\mathbf{e})) - \ln(P_{\Theta^{(2)}}(c|\mathbf{e})) \right| \leq \sum_{i=1}^K \left| \ln(P_{\Theta_i}(c|\mathbf{e})) - \ln(P_{\Theta_{i-1}}(c|\mathbf{e})) \right|,$$

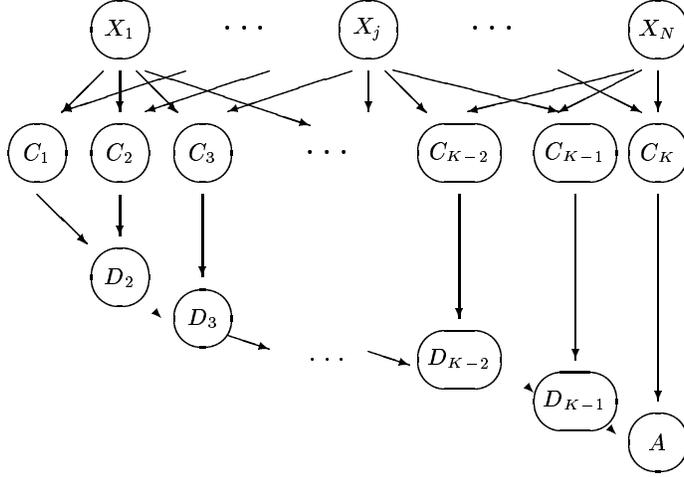


Figure 1: Belief Net structure corresponding to arbitrary SAT problem [Coo90]

and each  $|\ln(P_{\Theta_i}(c|\mathbf{e})) - \ln(P_{\Theta_{i-1}}(c|\mathbf{e}))|$  is based on changing a single CPTable entry. We therefore need only show  $|\ln(P_{\Theta_i}(c|\mathbf{e})) - \ln(P_{\Theta_{i-1}}(c|\mathbf{e}))| \leq \frac{\epsilon}{6K}$ . For any value of  $z = \theta_{d_i|\mathbf{f}_i}$ , let  $f(z) = \ln(P_{\Theta[z]}(c|\mathbf{e}))$ , where  $\Theta[z]$  be the BN whose first  $i-1$  CPTable entries come from  $\Theta^{(1)}$ , whose final  $K-i-1$  entries come from  $\Theta^{(2)}$ , and whose  $i^{\text{th}}$  CPTable entries is  $z$ ; hence  $f(\theta_{d_i|\mathbf{f}_i}^{(1)}) = \ln(P_{\Theta_i}(c|\mathbf{e}))$ , and  $f(\theta_{d_i|\mathbf{f}_i}^{(2)}) = \ln(P_{\Theta_{i+1}}(c|\mathbf{e}))$ . As this function is continuous, we know that

$$|f(a) - f(b)| = \frac{\partial f(z)}{\partial z} [b - a]$$

for some  $z \in [a, b]$ . As  $f(z) = \ln(P_{\Theta[z]}(c, \mathbf{e})) - \ln(P_{\Theta[z]}(\mathbf{e}))$ , we see that

$$\begin{aligned} \frac{\partial f(z)}{\partial z} &= \frac{1}{P_{\Theta[z]}(c, \mathbf{e})} P_{\Theta[z]}(c, \mathbf{e} | d_i, \mathbf{f}_i) \times P_{\Theta[z]}(\mathbf{f}_i) - \frac{1}{P_{\Theta[z]}(\mathbf{e})} P_{\Theta[z]}(\mathbf{e} | d_i, \mathbf{f}_i) \times P_{\Theta[z]}(\mathbf{f}_i) \\ &= \frac{1}{z} [P_{\Theta[z]}(d_i, \mathbf{f}_i | c, \mathbf{e}) - P_{\Theta[z]}(d_i, \mathbf{f}_i | \mathbf{e})] \end{aligned}$$

which means that  $|\frac{\partial f(z)}{\partial z}| \leq 1/z \leq 1/\gamma$ . (The second inequality follows from the assumption that we are only considering  $\Theta \in \mathcal{BN}_{\Theta \succeq \gamma}(G)$ .) Hence,

$$\begin{aligned} |\ln(P_{\Theta_{i+1}}(c|\mathbf{e})) - \ln(P_{\Theta_i}(c|\mathbf{e}))| &= |f(\theta_{d_i|\mathbf{f}_i}^{(2)}) - f(\theta_{d_i|\mathbf{f}_i}^{(1)})| \\ &\leq \frac{1}{\gamma} \times |\theta_{d_i|\mathbf{f}_i}^{(2)} - \theta_{d_i|\mathbf{f}_i}^{(1)}| \leq \frac{1}{\gamma} \times \frac{\gamma\epsilon}{6K} = \frac{\epsilon}{6K}. \end{aligned}$$

To prove Equation 2: Observe first that the probability of any event must be at least the product of  $N$  CPTable entries, and hence  $P_{\Theta}(c) \geq \gamma^N$  for any  $c$  and any  $\Theta \in \mathcal{BN}_{\Theta \succeq \gamma}(G)$ . This means the value of  $-\ln(P_{\Theta}(c|\mathbf{e}))$ , and hence  $\text{LCL}_{sq}(\Theta)$  for any distribution  $sq$ , is between 0 and  $-N \ln \gamma$ .

As the queries  $q = P(c, \mathbf{e})$  are drawn at random from a stationary distribution, we can view the quantity  $\ln P_{\Theta}(q)$  as an iid random value, whose range is  $[0, -N \ln \gamma]$  and whose expected value is  $\text{LCL}(\Theta)$ . Hoeffding's Inequality bounds the chance that the empirical average score after  $M$  iid examples (here  $\widehat{\text{LCL}}^{(S)}(\Theta)$ ) will be far away from the true mean  $\text{LCL}(\Theta)$ :

$$P(|\widehat{\text{LCL}}^{(S)}(\Theta) - \text{LCL}(\Theta)| > \frac{\epsilon}{6}) < 2 \exp[-2M((\epsilon/6)/N \ln \gamma)^2]. \quad (3)$$

Here, we want the right-hand-side to be under  $\delta/L$ , which requires  $M = M(\epsilon, \delta) = 2 \left( \frac{3N \ln \gamma}{\epsilon} \right)^2 \ln \left( \frac{2L}{\delta} \right)$ . ■

**Proof of Theorem 2:** We reduce 3SAT to our task, using a construction similar to the one in [Coo90]: Given any 3-CNF formula  $\varphi \equiv \bigwedge C_i$ , where each  $C_i \equiv \bigvee \pm X_{ij}$ , we construct the network shown in Figure 1, with one node for each variable  $X_i$  and one for each clause  $C_j$ , with an arc from  $X_i$  to  $C_j$  whenever  $C_j$  involves  $X_i$  — e.g., if

Table 1: Queries used in proof of Theorem 2

$X_1$	$X_2$	$X_3$	$X_4$	$\dots$	$X_n$	$A$
0	1	0				0
0		0	1			0
	$\vdots$					$\vdots$
0		1		1		0
						1

$C_1 = x_1 \vee \neg x_2 \vee x_3$  and  $C_2 = \neg x_1 \vee \neg x_3 \vee x_4$ , then there are links to  $C_1$  from each of  $X_1, X_2$  and  $X_3$ , and to  $C_2$  from  $X_1, X_3$  and  $X_4$ . In addition, we include  $K - 1$  other boolean nodes,  $\{D_2, \dots, D_{K-1}, A\}$ , where  $D_j$  is the child of  $D_{j-1}$  and  $C_j$ , where  $D_1$  is identified with  $C_1$ , and  $A$  is used for  $D_K$ .

Here, we intend each  $C_i$  to be true if the assignment to the associated variables  $X_{i1}, X_{i2}, X_{i3}$  satisfies  $C_i$ ; and  $A$  corresponds is the conjunction of those  $C_i$  variables. We do this using all-but-the-final instances in Table 1. (Note only 3 of the  $X_i$  variables are specified in each of these instances; the other  $n - 3$   $X_i$ s are not, nor are any  $C_j$ s nor  $D_k$ s.) There is one such instance for each clause, with exactly the assignment (of the 3 relevant variables) that falsifies this clause. Hence, the first line corresponds to  $C_1 \equiv x_1 \vee \neg x_2 \vee x_3$ . The final instance is just stating that the prior value for  $A$  should  $P(+a) = 1.0$ . The “label” of each instance always corresponds to the single variable  $A$ .

We now prove, in particular, that

There is a set of parameters for the structure in Figure 1, producing the  $\widehat{\text{LCL}}(\cdot)$ -score, over the queries in Table 1, of 0

*iff*

there is a satisfying assignment for the associated  $\varphi$  formula.

$\Leftarrow$ : Just set the CPTable for each  $C_i$  to be the disjunction of the associated  $X_{i1}, X_{i2}, X_{i3}$  variables (its parents), with the appropriate  $\pm$  parity. *E.g.*, using  $C_1 \equiv x_1 \vee \neg x_2 \vee x_3$ , then  $C_1$ 's CPTable would be

$x_1$	$x_2$	$x_3$	$P(+c_1   x_1, x_2, x_3)$
0	0	0	1.0
0	0	1	1.0
0	1	0	0.0
0	1	1	1.0
1	0	0	1.0
1	0	1	1.0
1	1	0	1.0
1	1	1	1.0

Similarly set the CPTables for the  $D_j$  to correspond to the conjunction of its 2 parents  $D_j = D_{j-1} \wedge C_j$ ; *e.g.*,

$D_4$	$C_5$	$P(+d_5   D_4, C_5)$
0	0	0.0
0	1	0.0
1	0	0.0
1	1	1.0

Finally, set  $X_i$  to correspond to the satisfying assignment; *i.e.*, if  $X_1 = 1$ , then  $\frac{P(+x_1)}{1.0}$ ; and if *i.e.*, if  $X_4 = 0$ ,

then  $\frac{P(+x_4)}{0.0}$ . Note that these CPTable values satisfy all  $k + 1$  of the labeled instances.

$\Rightarrow$ : Here, we assume there is no satisfying assignment. Towards a contradiction, we can assume that there is a 0-LCL set of CPTable entries. This means, in particular, that  $P(+a | x_{i1}, x_{i2}, x_{i3}) = 0$ , where  $x_{i1}, x_{i2}, x_{i3}$  correspond to the assignment that violates the  $i$ th constraint. (*E.g.*, for  $C_1 \equiv x_1 \vee \neg x_2 \vee x_3$ , this would be  $X_1 = 0, X_2 = 1, X_3 = 0$ .)

Now consider the final labeled instance,  $P(a)$ . As there is no satisfying assignment, we know that each assignment  $\mathbf{x}$  violates at least one constraint. For notation, let  $\gamma^{\mathbf{x}}$  refer to one of these violations (say the one with the smallest index). So if  $\mathbf{x} = \langle 0, 1, 0, \dots \rangle$ , then  $\gamma^{\langle 0, 1, 0, \dots \rangle} = \langle X_1 = 0, X_2 = 1, X_3 = 0 \rangle$  corresponds to the violation of the first constraint  $C_1$ . We also let  $\beta^{\mathbf{x}}$  refer to the rest of the assignment.

Now observe

$$\begin{aligned} P(+a) &= \sum_{\mathbf{x}} P(+a, \mathbf{x}) \\ &= \sum_{\mathbf{x}} P(+a | \gamma^{\mathbf{x}}) \cdot P(\gamma^{\mathbf{x}}) \cdot P(\beta^{\mathbf{x}} | +a, \gamma^{\mathbf{x}}) \\ &= \sum_{\mathbf{x}} 0 \cdot P(\gamma^{\mathbf{x}}) \cdot P(\beta^{\mathbf{x}} | +a, \gamma^{\mathbf{x}}) = 0, \end{aligned}$$

which shows that the final instance will be mislabeled. This proves that there can be no set of CPtable values that produce 0 LCL-score when there are no satisfying assignments. ■

**Proof of Proposition 3:** Below, we will use  $P(\chi)$  to refer to  $P_{\Theta}(\chi)$ , the value the belief net with parameters  $\Theta$  will assign to the  $\chi$  event. In general, for any assignment  $Z$ ,

$$P(Z) = \sum_{\mathbf{f}'} \sum_{d'} P(Z | D=d', \mathbf{F}=\mathbf{f}') P(D=d' | \mathbf{F}=\mathbf{f}') P(\mathbf{F}=\mathbf{f}'). \quad (4)$$

As we assume the different CPtable rows are estimated independently, and  $\mathbf{F}$  is the set of parents of  $D$ , this means

$$\frac{\partial P(Z)}{\partial \beta_{d|\mathbf{f}}} = \sum_{d'} P(Z | d', \mathbf{f}) \frac{\partial P(d' | \mathbf{f})}{\partial \beta_{d|\mathbf{f}}} P(\mathbf{f}).$$

Recalling  $\theta_{d|\mathbf{f}} = P(d | \mathbf{f}) = e^{\beta_{d|\mathbf{f}}} / \sum_{d'} e^{\beta_{d'|\mathbf{f}}}$ , observe that  $\frac{\partial P(d | \mathbf{f})}{\partial \beta_{d|\mathbf{f}}} = \theta_{d|\mathbf{f}}(1 - \theta_{d|\mathbf{f}})$ , and when  $d \neq d'$ ,  $\frac{\partial P(d' | \mathbf{f})}{\partial \beta_{d|\mathbf{f}}} = -\theta_{d|\mathbf{f}} \theta_{d'|\mathbf{f}}$ . This means  $\frac{\partial P(Z)}{\partial \beta_{d|\mathbf{f}}} = P(Z, d, \mathbf{f}) - \theta_{d|\mathbf{f}} P(Z, \mathbf{f})$ .

Hence, as  $\ln P(c | \mathbf{e}) = \ln P(c, \mathbf{e}) - \ln P(\mathbf{e})$ ,

$$\begin{aligned} \frac{\partial \ln P(c | \mathbf{e})}{\partial \beta_{d|\mathbf{f}}} &= \frac{\partial \ln P(c, \mathbf{e})}{\partial \beta_{d|\mathbf{f}}} - \frac{\partial \ln P(\mathbf{e})}{\partial \beta_{d|\mathbf{f}}} \\ &= \frac{1}{P(c, \mathbf{e})} \frac{\partial P(c, \mathbf{e})}{\partial \beta_{d|\mathbf{f}}} - \frac{1}{P(\mathbf{e})} \frac{\partial P(\mathbf{e})}{\partial \beta_{d|\mathbf{f}}} \\ &= \frac{1}{P(c, \mathbf{e})} [P(c, \mathbf{e}, d, \mathbf{f}) - \theta_{d|\mathbf{f}} P(c, \mathbf{e}, \mathbf{f})] - \frac{1}{P(\mathbf{e})} [P(\mathbf{e}, d, \mathbf{f}) - \theta_{d|\mathbf{f}} P(\mathbf{e}, \mathbf{f})] \\ &= [P(d, \mathbf{f} | c, \mathbf{e}) - P(d, \mathbf{f} | \mathbf{e})] - \theta_{d|\mathbf{f}} [P(\mathbf{f} | c, \mathbf{e}) - P(\mathbf{f} | \mathbf{e})]. \quad \blacksquare \end{aligned}$$

## References

[Coo90] G.F. Cooper. The computational complexity of probabilistic inference using Bayesian belief networks. *Artificial Intelligence*, 42(2–3):393–405, 1990.