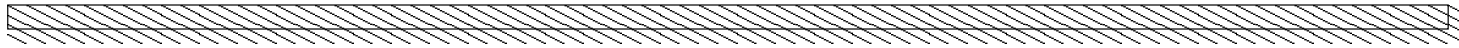


Proof for the Equivalence Between Some Best-First Algorithms and Depth-First Algorithms for AND/OR Trees

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Overview



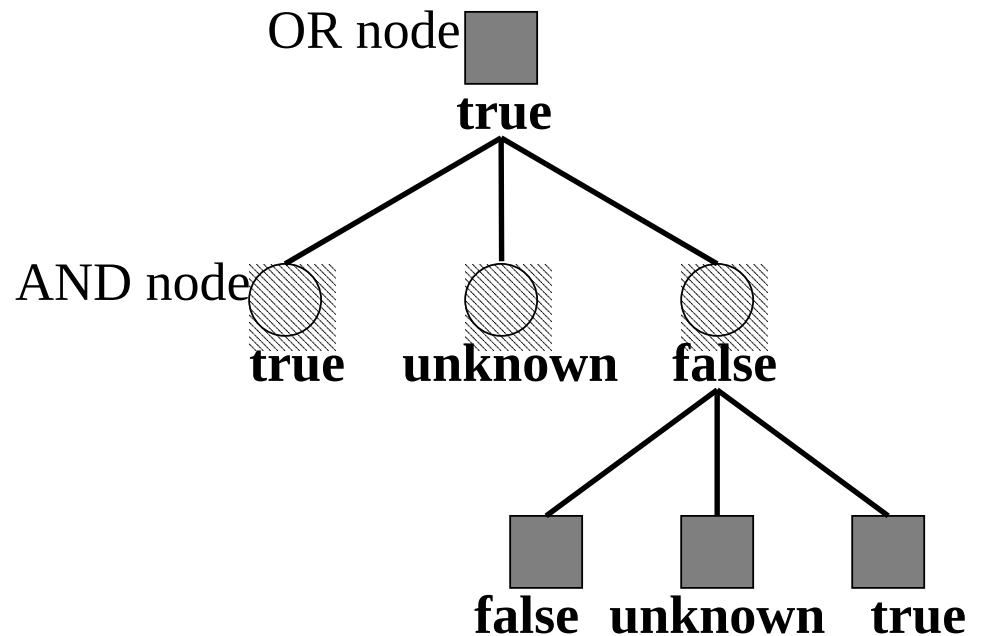
- Explanation of an AND/OR tree
 - Proof number and disproof number
 - Most-proving node
- Introduction of some algorithms
 - Pn-search
 - Df-pn-search
- Proof the outline of the equivalence
 - In the meaning of expanding always a most-proving node
- Experimental results
- Conclusion

An AND/OR Tree

- Two players ... first player and second player
- Three values ... false, unknown, and true
- Ordering (at OR nodes) ... $\text{false} < \text{unknown} < \text{true}$
- Ordering (at AND nodes) ... $\text{true} < \text{unknown} < \text{false}$

The final value of the root

- $\text{true} \Rightarrow$ proof solution
(proved)
- $\text{false} \Rightarrow$ disproof solution
(disproved)



Proof Number and Disproof Number

- If e is proved



a is proved

⇒ Proof number of $a = 1$

- If e and i is disproved

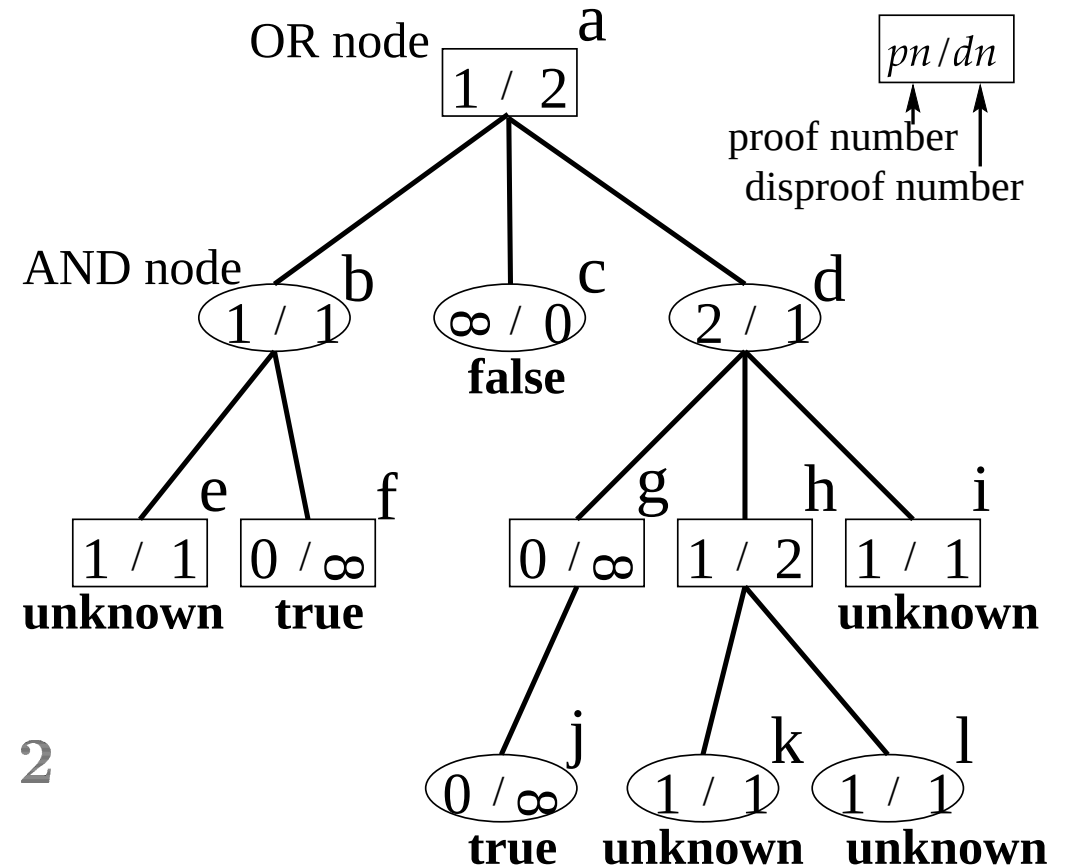


a is disproved

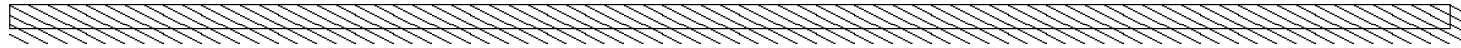
⇒ Disproof number of $a = 2$

- e is the most-proving node

⇒ Affecting to both proof number and disproof number



Calculation of Proof Number and Disproof Number



- At an OR node n

- One of its children is proved $\implies$ proved
- All the children is disproved $\implies$ disproved

$$\begin{aligned} n.pn &= \min_{n_{\text{child}} \in \text{children of } n} n_{\text{child}.pn} && \Leftarrow \\ n.dn &= \sum_{n_{\text{child}} \in \text{children of } n} n_{\text{child}.dn} \end{aligned}$$

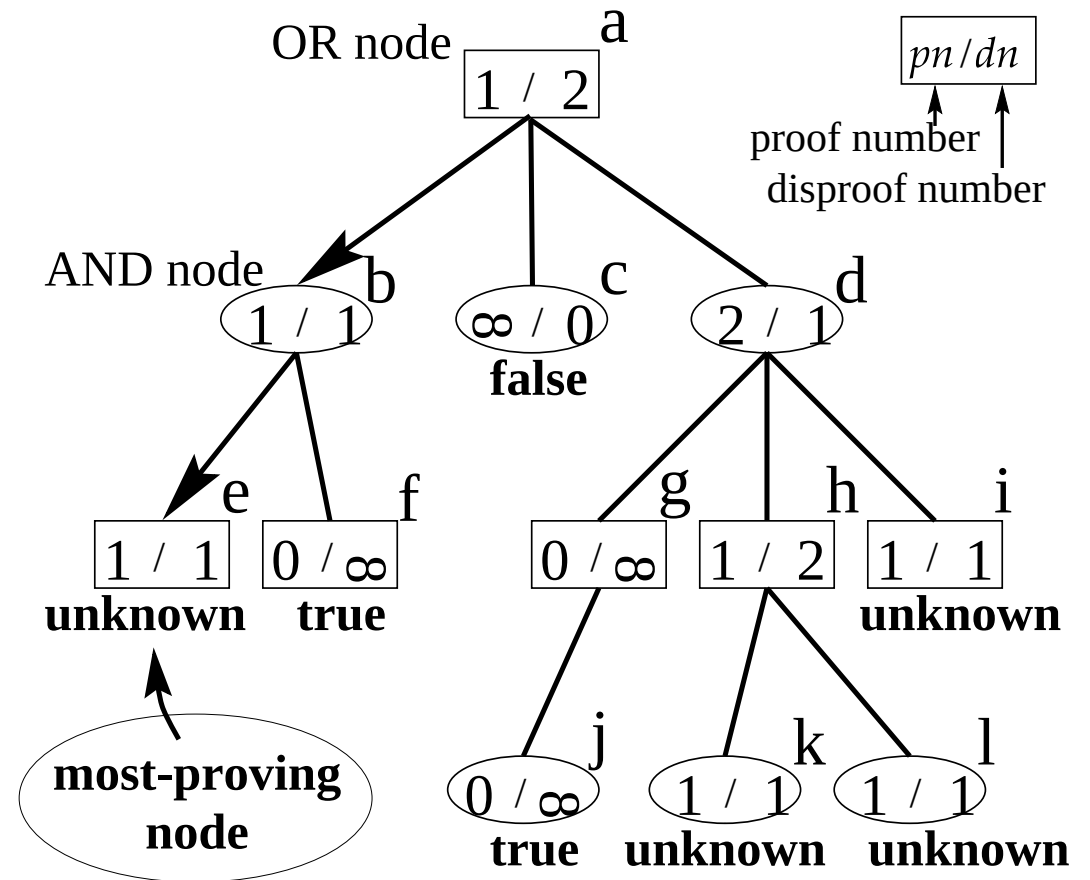
- At an AND node n

- All the children is proved $\implies$ proved
- One of its children is disproved $\implies$ disproved

$$\begin{aligned} n.pn &= \sum_{n_{\text{child}} \in \text{children of } n} n_{\text{child}.pn} \\ n.dn &= \min_{n_{\text{child}} \in \text{children of } n} n_{\text{child}.dn} \end{aligned}$$

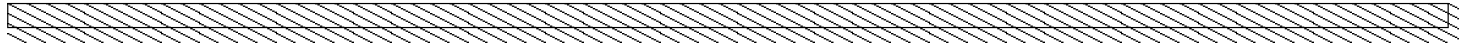
A Most-Proving Node

- Starting from the root
- Selection among the child
 - At an OR node
 - ... child of least proof number
 - At an AND node
 - ... child of least disproof number



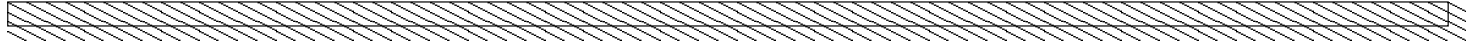
⇒ Finally reaches to the most-proving node

Related Work and its Classification



	criteria of evaluation used	
	only proof number	proof number and disproof number
best-first	AO*[Nilson,1980] (Elkan[1989])	pn-search[Allis,1994]
depth-first	Seo's Algorithm [1995]	df-pn-search[this paper] (PDS[Nagai,1998])

Purpose



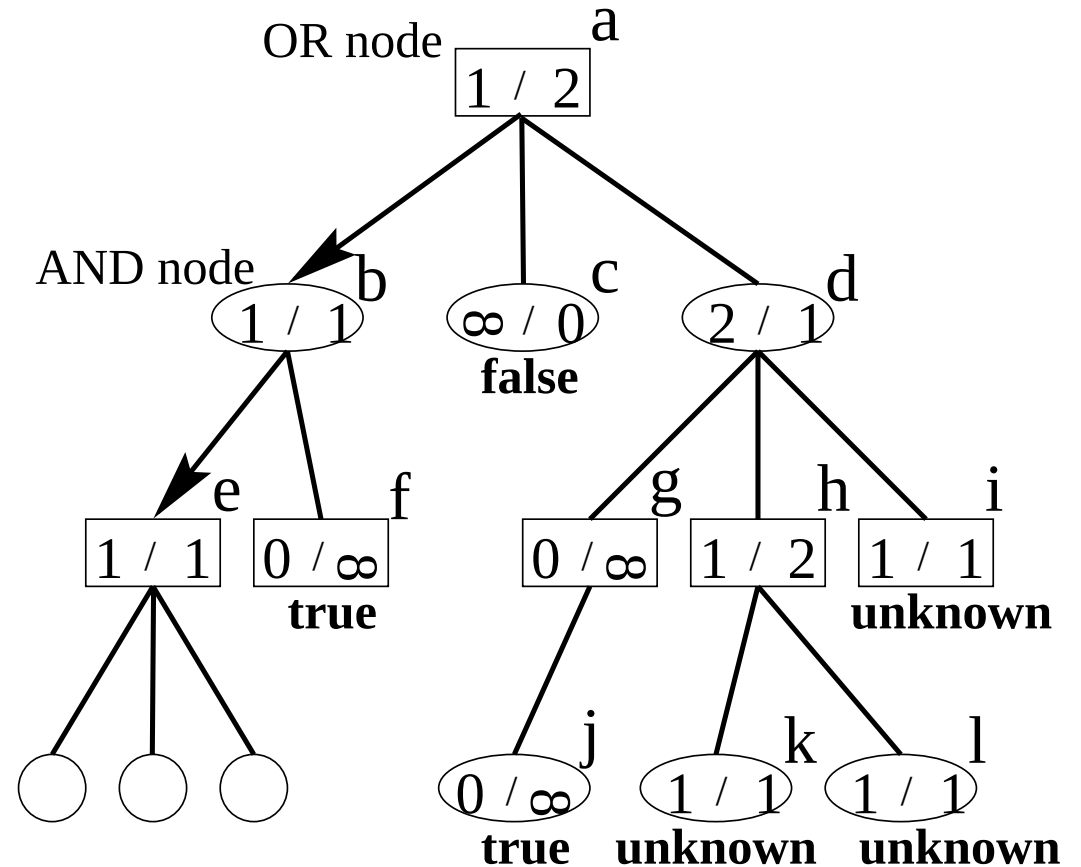
- Proof of the equivalence
 - between AO* and Seo's Algorithm
 - between pn-search and df-pn-search $\Leftarrow$ focus

Pn-search

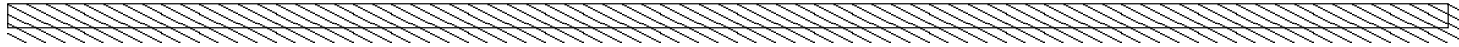
1. Select a most-proving node
2. Generate all the children
3. Propagate toward the root

Repeat until
the solution is found

⇒ Goes back and forth



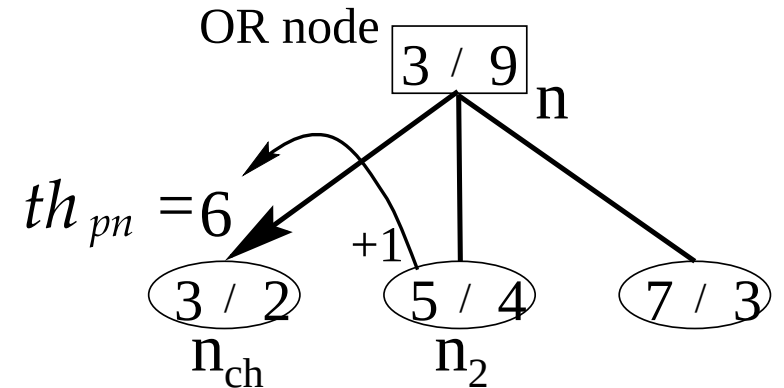
Df-pn-search



- At each node n , search below n

while $n.pn < n.th_{pn}$

and $n.dn < n.th_{dn}$



- At each OR node n ,

– search the child n_{ch} with minimum proof number

– assign

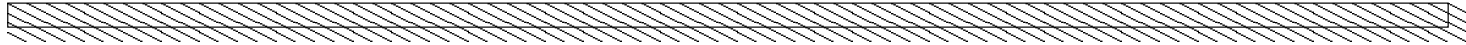
$$n_{ch}.th_{pn} = \min(n.th_{pn}, \underline{\underline{n_2.pn + 1}})$$

$$n_{ch}.th_{dn} = n.th_{dn} + n_{ch}.dn - \Sigma n_{child}.dn$$

$n_2 \dots$ is the child with second minimum proof number

Proof of the equivalence (1)

(! (! Course of the Proof (! (!



- Pn-search always expands most-proving node



self-evident

- Df-pn-search always expands most-proving node



proof outline

$\Rightarrow$ pn-search and df-pn-search is equivalent

Proof of the equivalence (2)

(!! Expansion of Most-Proving Node during Df-pn-search (!!)



A leaf node is now to be expanded

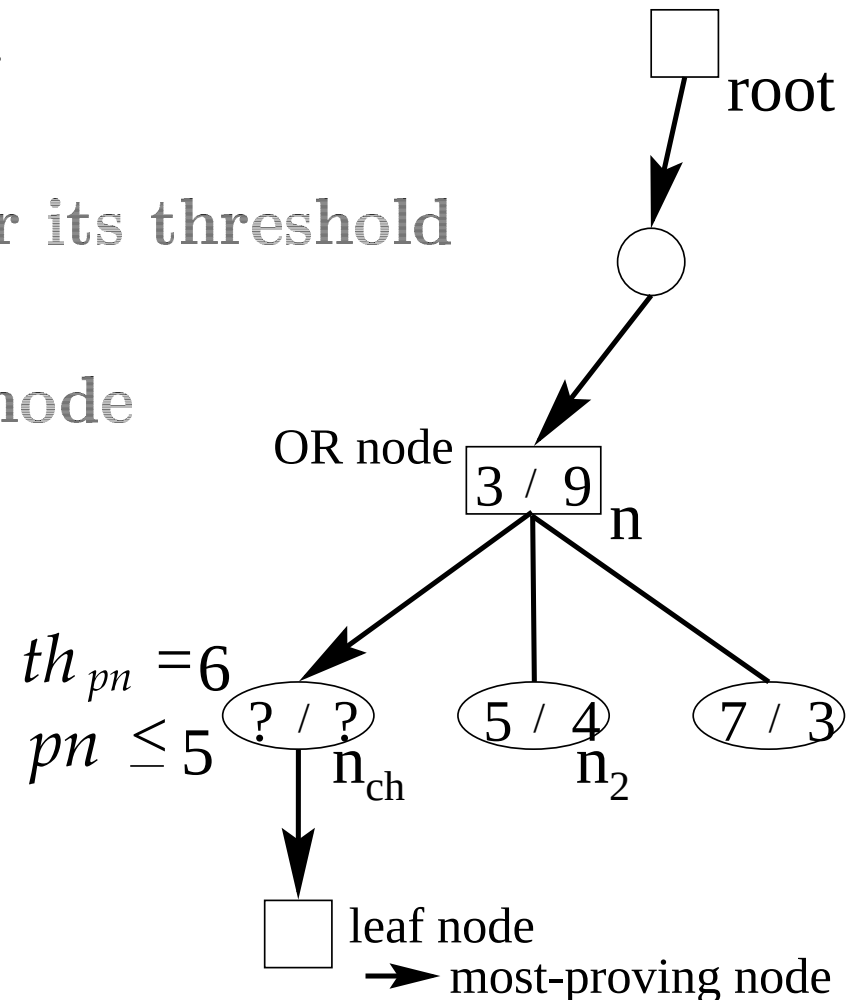
- While (dis)proof number is under its threshold



Most-proving node is below the node

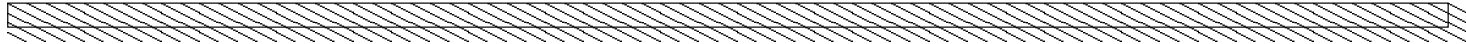
- The same thing can be said
with all the active nodes

$\Rightarrow$ Df-pn-search always expands
most-proving node



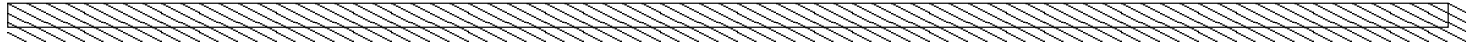
Proof of the equivalence (3)

(!! More Strong Equivalence !!)



- A narrow definition of most-proving node
 - $\Rightarrow$ Most-proving node is defined uniquely
(but dynamically)
 - $\Rightarrow$ Same most-proving nodes are expanded in the same order
as with both pn-search and df-pn-search

Difference Between Pn-Search and Df-Pn-Search



- **Pn-search** searches

$a \rightarrow b \rightarrow e \rightarrow b \rightarrow a \rightarrow b \rightarrow e \rightarrow m$

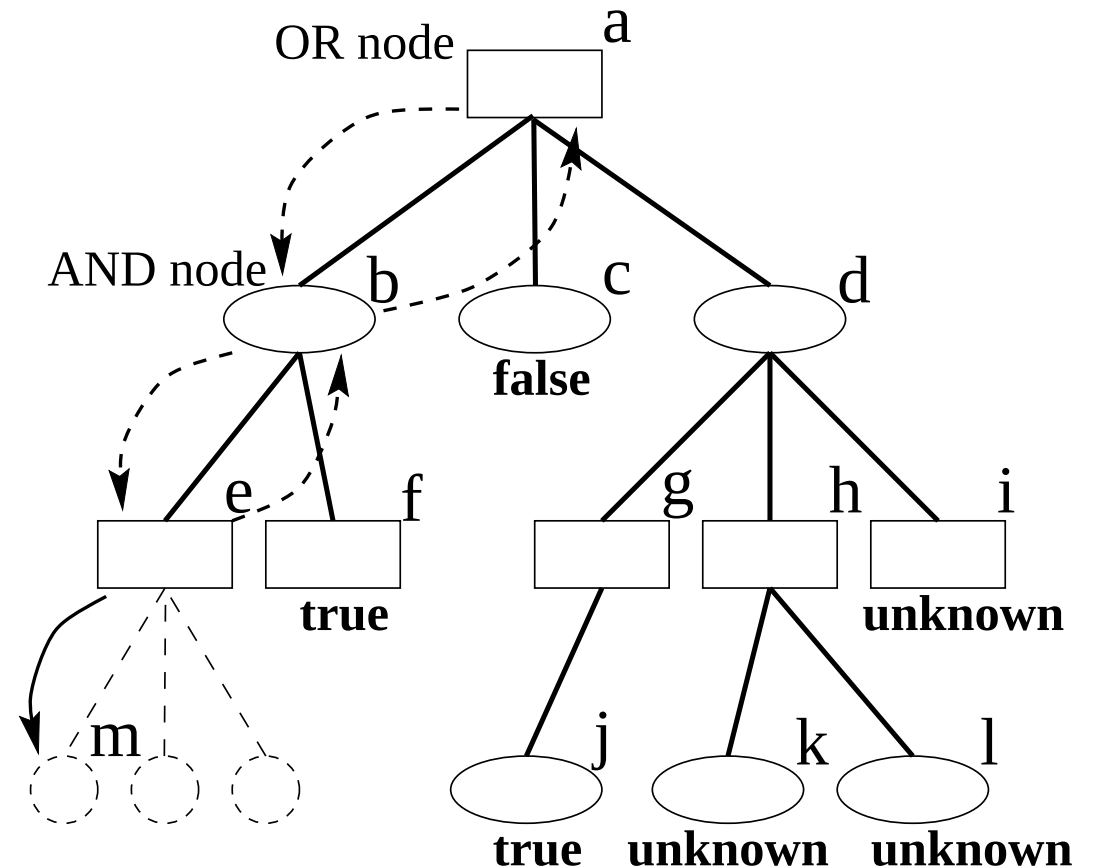
- **Df-pn-search** searches

$a \rightarrow b \rightarrow e \rightarrow m$

As with df-pn-search,

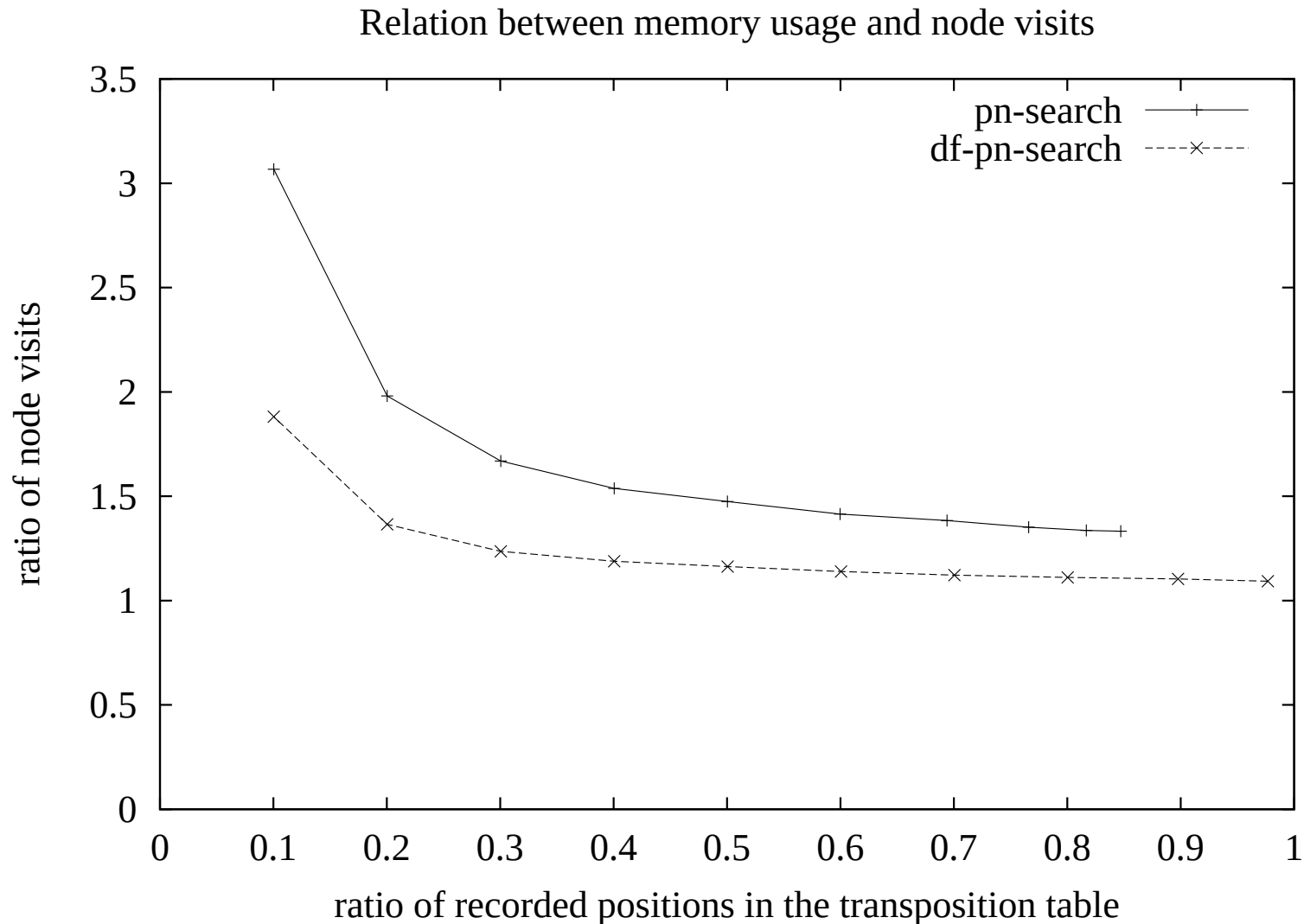
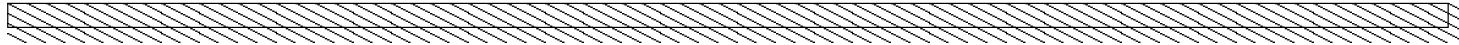
- **No unnecessary
back and forth**

↑
Two thresholds



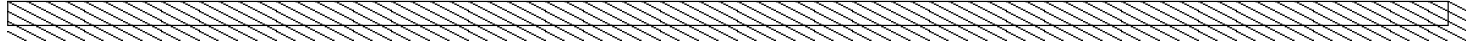
Experimental Result

(151 Othello positions with 15 vacant squares appeared at PrincetonII)



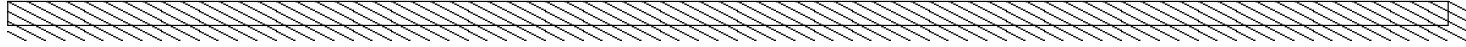
⇒ Df-pn-search needs only 60-85% node visits

Conclusion



- Proposal of a new depth-first algorithm (df-pn-search)
- Proof of the equivalence
between pn-search and df-pn-search
- Experimental result show the superiority of
df-pn-search to pn-search

Future Work



- Improvement of Seo's algorithm
- Using the information of heuristic evaluation function