A Search-based Approach for Solving Sum Games

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Overview

- Combinatorial games - mathematics vs computing science

- Goals of this research - Why another CGT solver?

- Work on single games

- MCGS - A **M**inimax-based **C**ombinatorial **G**ame **S**olver

  - Framework

  - Current state and future plans

Combinatorial Games Research

in Mathematics vs Computing Science

- In Mathematics:

  - Build theories of CGT, deeply understand games with specific structure

  - Prove theorems about whole (infinite) **classes** of games

- In Computing Science:

  - Solve **specific** game positions

  - Find efficient and **general** algorithms to solve games

Mathematics + Computing Science for CGT

- **Computing Science helps** mathematicians:

  - Explore (large) numbers of games in a class

  - Help to find patterns, verify or refute conjectures

- **Math helps** computing scientists:

  - Improve efficiency of computation

  - Rules to simplify games

  - Basis to exploit sum structure, and partial order of games

Project Background and (Short) History

- 2017 and 2022: taught graduate courses - Algorithms in Combinatorial Game Theory

- Several course assignments and research projects: solve Clobber by search

- Strongest solver by Taylor Folkersen's team. Later in 2022: Taylor improves it further

- 2023/24: Henry Du creates two solvers for NoGo. The newer one is based on CGT

- Goal: build general search-based solver system

  - MCGS, Minimax-based Combinatorial Game Solver

- Fall 2024: started implementation.

- December 2024: hired Taylor Folkersen

Single Game Results - Linear Clobber

- Conjecture (Albert et al 2005, [1]): $(BW)^n$ is an N position for n != 3

  - Previously [1]: proven up to n=19

  - Now proven up to **n = 33** (!) (Folkersen et al 2022, Folkersen 2022)
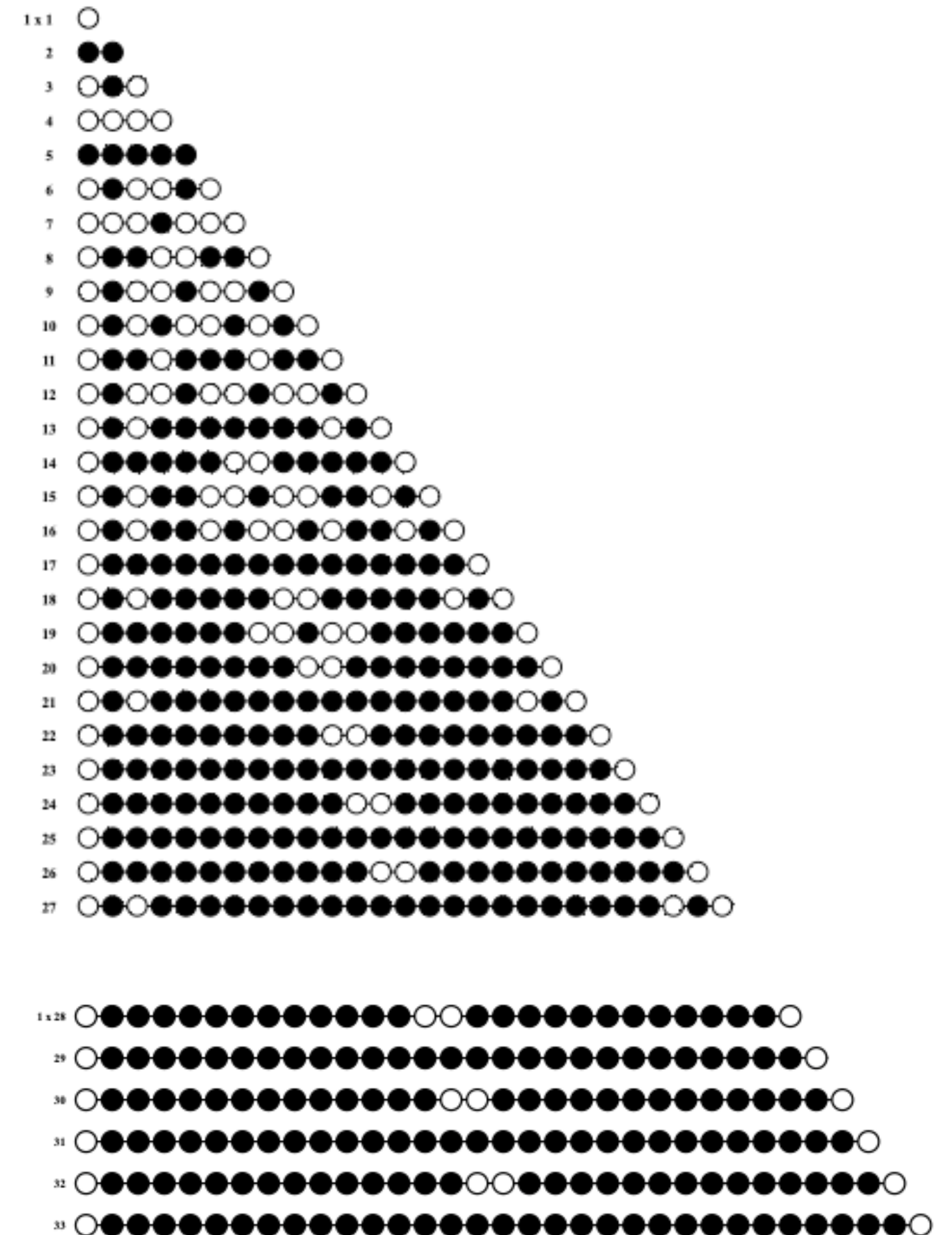

  - Efficient solver uses many CGT-inspired techniques

  [1] M. Albert, J.P. Grossman, R. Nowakowski, and D. Wolfe. An introduction to Clobber. INTEGERS: The Electronic Journal of Combinatorial Number Theory, 5(2), 2005.

  Note: at CGTC V, Hayward et all announced a general proof of the conjecture!

Single Game Results - Linear NoGo

- Solved all empty boards **up to 1x39** (Du et al 2023, Du et al 2024)

  - Only 1x1 and 1x4 are P positions

  - All others are N positions

- Solved all opening moves as win or loss (N or P) up to 1x33 - wins = white stones in figure, losses = black stones

- Up to 1x27 (Du et al 2023) used traditional (full board) minimax search

- Larger boards: CGT-based techniques (Du et al 2024)

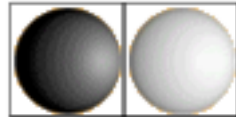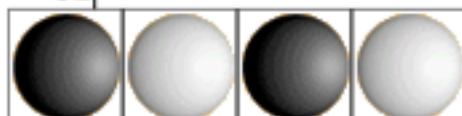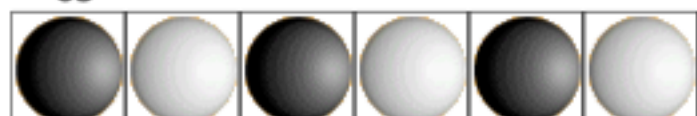 - Reduces the search by about two orders of magnitude for large boards

Goals of This Research

- **General** combinatorial game solver using the techniques learned from solving single games

- Answer the basic **question: Who wins game G**, if player p goes first?

- For short combinatorial games (finite, normal play)

- Efficient minimax search-based solver

  - Use **sum** game structure, **equality** and **partial ordering** of games for simplification

  - Work from first principles

  - **Avoid** computing canonical forms

Why Avoid Canonical Forms? (1)
Canonical form of 1xn Clobber boards (computed with CGSuite)

> G1	*
> G2	+-{*,^}
> G3	0
> G4	+-{^[2],{^*,^^|0,^*,+-{0,^*}}}
> G5	+-{^<2>*,^,{0|^,+-{*,^}},{^^^*|v,+-{*,^}}}
> G6	+-{0,^*}

Why Avoid Canonical Forms? (2)

> G7

+-{{^[2]*,^^,{0|{0|^,+-{0,^*}},{^*,^^|0,^*,+-{0,^*}}},{^3*|0,{^*,^^|0,^*,+-{0,^*}}}|0,{^^|0,^*},{0,{^*,^^|0,^*,+-{0,^*}}||v*,+-{0,^*}|0},+-{^<2>,^*,{0|^*,+-{0,^*}},{^^|v*,+-{0,^*}}},{^^,{^*,^^|0,^*,+-{0,^*}}||0,v*}},{{^^*|^},{^3*|0,{^*,^^|0,^*,+-{0,^*}}}|0,*,{0,*,{^,^*| 0,v*}|vv,v*}},{0|{0|^,+-{0,^*}},{^*,^^|0,^*,+-{0,^*}}||+-{^[2],{^*,^^|0,^*,+-{0,^*}}},{^,^^*|*,^,+-{*,^}||{0,v*,+-{0,^*}|vv,v*},+-{0,^*}}},{0|{0|^,+-{0,^*}},{^*,^^|0,^*,+-{0,^*}}||{0|^,+-{*,^}||*,{*,v,+-{*,^}|vv*,v}},{0|^,+-{0,^*}||0,{0,v*,+-{0,^*}|vv,v*}},{+-{0,^*}, {^*,^^|0,^*,+-{0,^*}}||*,v,+-{*,^}|vv*,v}},{{0|+-{0,^*}},{^,^*,{^^|^*}|^,^*,{^,^[2]*|0,*,{0,^*|v,v*}},{^*,^^|0,*,{0,^*|v,v*}}||+-{0,^*}, {*,^<2>,{^,^[2]*|0,*,{0,^*|v,v*}}|v,v*,{0,v*,{0,^*|v,v*}|vv,vv*}},+-{*,^},{+-{0,^*},{^*,^^|0,^*,+-{0,^*}}||*,v,+-{*,^}|vv*,v}}}

> G8

+-{{^,^^^*|*,^,+-{*,^}},{^^^*|+-{^<2>*,^,{0|^,+-{*,^}},{^^^*|v,+-{*,^}}},{^[2]*,{0|^,+-{*,^}},{^^,^^^*|*,^,{^,^*|*,v}}|*,{*,v,+-{*,^}|vv*,v}}}}

> G9

> G10

Why Avoid Canonical Forms? (3)

- We just want to solve for win/loss

- Computing canonical forms of all intermediate positions incurs a huge overhead

- Example: 1 ×16 empty NoGo board

  - Result in CGSuite: massive canonical form, 1201194 stops, takes several minutes

  - 1x17 in CGSuite: Java heapspace overflow

  - Solver from (Du et al 2024), from scratch without using precomputed database:

    - Solves 1x16 in less than 30 milliseconds, scales much further

Why use Minimax Search ?
Relation to Outcome Classes in CGT

- Two searches, one with each player (Left, Right) going first

- Solve yes/no question: Can the first player win?

- The total 2x2=4 results determine the outcome class

  - Black can always win, no matter who goes first: G > 0, in class L

  - White can always win: G < 0, in class R

  - First player wins, G $\parallel$ 0, in class N

  - Second player wins: G = 0, in class P

The Many Uses of Search in Solving Games

- Not just to directly solve games

- Many other uses make overall solution more efficient

- Prove if games are equal: $G = H$ iff $G - H = 0$, is second—player win

- Prove that two subgames are inverses and cancel: $G + H = 0$

- Dominated options: to show $G \geq H$,
  solve $G - H \geq 0$ by a **single** search: Right going first loses

- Replace subgame in sum by simplest equal game

- **Main Challenge**: **when** to try such searches for simplification?

MCGS - A Minimax-based Combinatorial Game Solver

- **Version 1 released!**

- Game-independent, for all short games

- Basic framework, not many optimisations so far

- Sample games implemented: linear Clobber, linear NoGo, Nim, Elephants and Rhinos

- Abstract games: integer, dyadic rational, nimber, up-star, switch

- Documentation and extensive test framework, simple file format

- Version 2 will add support for building and using databases of small games

Approach and Techniques Used..

..(or planned)..

- Abstract **game** class, specific games extend game, e.g. implement rules, decomposition

- Stack-based **sum game** data structure

- Efficient **play, undo move** on both subgame and sum level

- Supports game-specific rules to **split** game into more subgames

- (future) Simplify sumgame

- (future) Remove zeroes, and pairs of game+inverse

- (future) Replace game by simplest equal game

- (future) Static evaluation for early win/loss detection, use bounds on game values

Example - How to Add your Own Game to MCGS

```cpp
class elephants : public strip

    elephants(game_as_string);

    void play(move m, bw to_play);

    void undo_move();

    split_result split_implementation();

    move_generator* create_move_generator(bw to_play);

    void print(ostream);

    game* inverse();
```

Future Plans

- Version 2: add database, optimisations

- Future Versions and extensions

  - Many more optimisations

  - Support thermographs

  - Add specialised algorithms for impartial games

  - Implement more games including Clobber, NoGo on a grid

Thank you!