

(this is one of the quiz versions, the others were similar)

1. Playing go is an AI challenge because of its huge search space and because \_\_\_\_\_  
 \_\_\_\_\_ . AlphaGo uses its policy network

to

\_\_\_\_\_ . AlphaGo beat European Go champion \_\_\_\_\_ by score \_\_\_\_  
 \_\_\_\_

and later in the city of \_\_\_\_\_ beat Lee Sedol by score \_\_\_\_ — \_\_\_\_ .

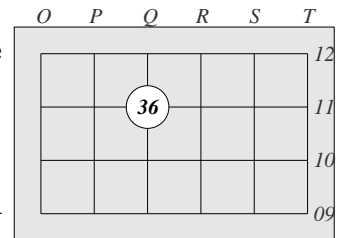
In Game 2 of the AlphaGo–Lee Sedol match, Lee left the room after  
 move 36 because \_\_\_\_\_ .

AlphaGo moved about \_\_\_\_\_ minutes later: mark move 37 on the

2. diagram.

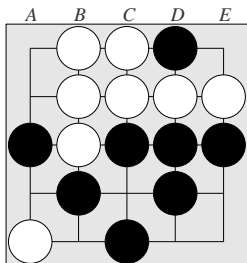
This move — called a \_\_\_\_\_

surprised Lee because \_\_\_\_\_



If this Go game ends at this position P, the Tromp-Taylor score is 8  
 stones plus \_\_\_\_\_ points for Black, 8 stones plus \_\_\_\_\_ points

3.



for White, so \_\_\_\_\_ wins by \_\_\_\_\_. From P, if White plays at

E5 then Black's best reply (assume each player wants to maximize  
 her final score) is \_\_\_\_\_ and the game will eventually end a win

for \_\_\_\_\_ by \_\_\_\_\_ points. From P, if White plays at A4 then

Black's best reply (assume as above) is \_\_\_\_\_ and the game will  
 eventually end a win for \_\_\_\_\_ by \_\_\_\_\_ points.

On this maze, label empty cells 1 to 11 in the order in which they are discovered by a recursive depth-first search that starts at + .

```

X X X X X X
X + _ _ _ X
X _ _ _ _ X
X _ _ _ _ X
X X X X X X

```

Assume that cell neighbours are checked in this order: below, right,

4. left, above.

If a breadth-first search of a graph with 100 nodes and 4950 edges takes about 10 minutes, then a BFS of a graph with 200 nodes and 19900 edges should take about \_\_\_\_\_ minutes.

5. This 2x3 sliding tile puzzle has \_\_\_\_\_ inversions and an \_\_\_\_\_ number of columns so it (**circle one**) is is not solvable. Below the puzzle, draw the next two levels of the search tree.

```

4 2 5
1 3

```

6. I based `stile_search.py` that I showed in class on (**circle one**) BFS DFS random walk because

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The number of solvable 3x3 positions is (**circle one**)  $3! \cdot 3!/2 \cdot 9! \cdot 9!/2$  . The number of solvable 3x5 positions is (**circle one**)  $5! \cdot 5!/2 \cdot 8! \cdot 8!/2 \cdot 15! \cdot 15!/2$  . The runtime of `stile_search.py` is proportional to (**circle one**) number of search-tree nodes number of search-tree nodes-plus-edges square of search-tree number of nodes square of search-tree number of nodes-plus-edges, so it will take about \_\_\_\_\_ times as long to solve a 3x5 position as a 3x3 position.

1. see the AlphaGo Nature paper abstract and wikipedia
2. Lee left for a cigarette break. AlphaGo took about 2-3 minutes for each move. shoulder hit. no top pro had played a shoulder hit this far from the edge (5th line) in a competitive match ... and as he studied the move, he realized it was unexpectedly strong
3. Black 8 stones 4 points territory, White 8 stones 0 points territory, Black wins by 4.

After this (bad) move by White, Black's most aggressive move is to play at the end of the White row of 4 stones. White can eventually kill this Black stone, but will not be able to give the White group 2 eyes. Black can kill all, and win by 25.

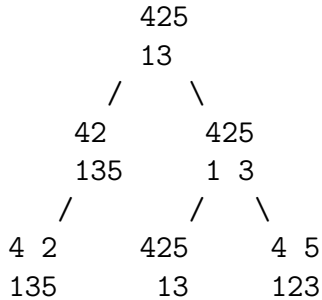
After this (good) move by White, Black should play next to the single White stone, making Black's group safe. Now White plays where Black played in the previous question, giving White's group 2 eyes. After each player has killed opponent stones inside their territory, White has 11 stones plus points, Black has 15, so Black wins by 4.

- |    |  |   |    |    |
|----|--|---|----|----|
|    |  | 9 | 10 | 11 |
| 4. | at right are dfs labels for this order: below, right, left, above. | 1 | 8  | 7  |
|    | bfs takes time proportional to nodes plus edges                    | 2 | 3  | 4  |
|    |  |   | 5  | 6  |

5. odd number of columns, so solvable if and only if even number of inversions

42513 42531 43521 have respectively 6 7 8 inversions, so yes no yes.

here is start of search tree for first.



6. bfs, because I wanted a shortest solution

number of  $n \times m$  positions is  $(n \times m)!$ , exactly half are solvable, proportional to nodes-plus-edges, so  $(15!/2)$  over  $9!/2 = 15 * 14 * 13 * 12 * 10$  times as long