Outline

- Introduction
- Background
- Distributed DBMS Architecture
- Distributed Database Design
- Semantic Data Control
- Distributed Query Processing
  - Distributed Transaction Management
    - Transaction Concepts and Models
    - Distributed Concurrency Control
    - Distributed Reliability
  - Parallel Database Systems
  - Distributed Object DBMS
  - Database Interoperability
  - Concluding Remarks

Transaction

A transaction is a collection of actions that make consistent transformations of system states while preserving system consistency.

- concurrency transparency
- failure transparency

```
Begin Transaction
```

```
Database in a consistent state
```

```
Database may be temporarily in an inconsistent state during execution
```

```
Execution of Transaction
```

```
End Transaction
```

```
Database in a consistent state
```

Transaction Example – A Simple SQL Query

Transaction BUDGET_UPDATE
begin
  EXEC SQL UPDATE PROJ
  SET      BUDGET = BUDGET * 1.1
  WHERE   PNAME = “CAD/CAM”
end.

Example Database

Consider an airline reservation example with the relations:

FLIGHT(FNO, DATE, SRC, DEST, STSOLD, CAP)
CUST(CNAME, ADDR, BAL)
FC(FNO, DATE, CNAME, SPECIAL)
Example Transaction – SQL Version

```sql
BEGIN_TRANSACTION Reservation
begin
    input(flight_no, date, customer_name);
    EXEC SQL UPDATE FLIGHT
    SET STSOLD = STSOLD + 1
    WHERE FNO = flight_no AND DATE = date;
    EXEC SQL INSERT INTO FC(FNO, DATE, CNAME, SPECIAL);
    VALUES (flight_no, date, customer_name, null);
    output("reservation completed")
end. {Reservation}
```

Termination of Transactions

```sql
BEGIN_TRANSACTION Reservation
begin
    input(flight_no, date, customer_name);
    EXEC SQL SELECT STSOLD,CAP
    INTO temp1,temp2
    FROM FLIGHT
    WHERE FNO = flight_no AND DATE = date;
    if temp1 = temp2 then
        output("no free seats");
        Abort
    else
        EXEC SQL UPDATE FLIGHT
        SET STSOLD = STSOLD + 1
        WHERE FNO = flight_no AND DATE = date;
        EXEC SQL INSERT INTO FC(FNO, DATE, CNAME, SPECIAL);
        VALUES (flight_no, date, customer_name, null);
        Commit
        output("reservation completed")
    endif
end. {Reservation}
```
Example Transaction – Reads & Writes

Begin_transaction Reservation
begin
  input(flight_no, date, customer_name);
  temp ← Read(flight_no(date).stsold);
  if temp = flight(date).cap then
    begin
      output(“no free seats”);
      Abort
    end
  else begin
    Write(flight(date).stsold, temp + 1);
    Write(flight(date).cname, customer_name);
    Write(flight(date).special, null);
    Commit;
    output(“reservation completed”)
  end
end. {Reservation}

Characterization

- Read set (RS)
  ⇒ The set of data items that are read by a transaction
- Write set (WS)
  ⇒ The set of data items whose values are changed by this transaction
- Base set (BS)
  ⇒ RS ∪ WS
Formalization

Let
\[ O_{ij}(x) \text{ be some operation } O_j \text{ of transaction } T_i \text{ operating on entity } x, \text{ where } O_j \in \{\text{read,write}\} \text{ and } O_j \text{ is atomic} \]

\[ OS_i = \bigcup_j O_{ij} \]

\[ N_i \in \{\text{abort,commit}\} \]

Transaction \( T_i \) is a partial order \( T_i = \{\Sigma_i, <_i\} \) where

1. \( \Sigma_i = OS_i \cup \{N_i\} \)
2. For any two operations \( O_{ij}, O_{ik} \in OS_i \), if \( O_{ij} = R(x) \) and \( O_{ik} = W(x) \) for any data item \( x \), then either \( O_{ij} <_i O_{ik} \) or \( O_{ik} <_i O_{ij} \)
3. \( \forall O_{ij} \in OS_i, O_{ij} <_i N_i \)

Example

Consider a transaction \( T \):

Read(x)
Read(y)
\( x \leftarrow x + y \)
Write(x)
Commit

Then
\[ \Sigma = \{R(x), R(y), W(x), C\} \]
\[ <= \{(R(x), W(x)), (R(y), W(x)), (W(x), C), (R(x), C), (R(y), C)\} \]
DAG Representation

Assume
\[ \leq \{(R(x), W(x)), (R(y), W(x)), (R(x), C), (R(y), C), (W(x), C)\} \]

Properties of Transactions

**Atomicity**
- all or nothing

**Consistency**
- no violation of integrity constraints

**Isolation**
- concurrent changes invisible \( \implies \) serializable

**Durability**
- committed updates persist
Atomicity

- Either all or none of the transaction's operations are performed.
- Atomicity requires that if a transaction is interrupted by a failure, its partial results must be undone.
- The activity of preserving the transaction's atomicity in presence of transaction aborts due to input errors, system overloads, or deadlocks is called transaction recovery.
- The activity of ensuring atomicity in the presence of system crashes is called crash recovery.

Consistency

- Internal consistency
  - A transaction which executes alone against a consistent database leaves it in a consistent state.
  - Transactions do not violate database integrity constraints.
- Transactions are correct programs
Consistency Degrees

- **Degree 0**
  - $T$ does not overwrite dirty data of other transactions
  - Dirty data refers to data values that have been updated by a transaction prior to its commitment

- **Degree 1**
  - $T$ does not overwrite dirty data of other transactions
  - $T$ does not commit any writes before EOT

Consistency Degrees (cont’d)

- **Degree 2**
  - $T$ does not overwrite dirty data of other transactions
  - $T$ does not commit any writes before EOT
  - $T$ does not read dirty data from other transactions

- **Degree 3**
  - $T$ does not overwrite dirty data of other transactions
  - $T$ does not commit any writes before EOT
  - $T$ does not read dirty data from other transactions
  - Other transactions do not dirty any data read by $T$ before $T$ completes.
Isolation

- Serializability
  - If several transactions are executed concurrently, the results must be the same as if they were executed serially in some order.

- Incomplete results
  - An incomplete transaction cannot reveal its results to other transactions before its commitment.
  - Necessary to avoid cascading aborts.

Isolation Example

- Consider the following two transactions:
  
  \[ T_1: \text{Read}(x) \quad T_2: \text{Read}(x) \]
  
  \[ x \leftarrow x+1 \quad x \leftarrow x+1 \]
  
  \[ \text{Write}(x) \quad \text{Write}(x) \]
  
  \[ \text{Commit} \quad \text{Commit} \]

- Possible execution sequences:

  \[ T_1: \text{Read}(x) \quad T_1: \text{Read}(x) \]
  
  \[ T_1: x \leftarrow x+1 \quad T_1: x \leftarrow x+1 \]
  
  \[ T_1: \text{Write}(x) \quad T_2: \text{Read}(x) \]
  
  \[ T_1: \text{Commit} \quad T_1: \text{Write}(x) \]
  
  \[ T_2: \text{Read}(x) \quad T_2: x \leftarrow x+1 \]
  
  \[ T_2: x \leftarrow x+1 \quad T_2: \text{Write}(x) \]
  
  \[ T_2: \text{Write}(x) \quad T_1: \text{Commit} \]
  
  \[ T_2: \text{Commit} \quad T_2: \text{Commit} \]
SQL-92 Isolation Levels

Phenomena:

- Dirty read
  - $T_1$ modifies $x$ which is then read by $T_2$ before $T_1$ terminates; $T_1$ aborts $\Rightarrow T_2$ has read value which never exists in the database.

- Non-repeatable (fuzzy) read
  - $T_1$ reads $x$; $T_2$ then modifies or deletes $x$ and commits. $T_1$ tries to read $x$ again but reads a different value or can’t find it.

- Phantom
  - $T_1$ searches the database according to a predicate while $T_2$ inserts new tuples that satisfy the predicate.

SQL-92 Isolation Levels (cont’d)

- Read Uncommitted
  - For transactions operating at this level, all three phenomena are possible.

- Read Committed
  - Fuzzy reads and phantoms are possible, but dirty reads are not.

- Repeatable Read
  - Only phantoms possible.

- Anomaly Serializable
  - None of the phenomena are possible.
Durability

- Once a transaction commits, the system must guarantee that the results of its operations will never be lost, in spite of subsequent failures.
- Database recovery

Characterization of Transactions

Based on

- Application areas
  - non-distributed vs. distributed
  - compensating transactions
  - heterogeneous transactions
- Timing
  - on-line (short-life) vs batch (long-life)
- Organization of read and write actions
  - two-step
  - restricted
  - action model
- Structure
  - flat (or simple) transactions
  - nested transactions
  - workflows
Transaction Structure

- Flat transaction
  - Consists of a sequence of primitive operations embraced between a `begin` and `end` markers.
  - `Begin_transaction Reservation ...
  - `end`.
- Nested transaction
  - The operations of a transaction may themselves be transactions.
  - `Begin_transaction Reservation ...
  -  `Begin_transaction Airline ...
  -  `end. `{Airline}
  -  `Begin_transaction Hotel ...
  -  `end. `{Hotel}
  - `end. `{Reservation}

Nested Transactions

- Have the same properties as their parents ⇒ may themselves have other nested transactions.
- Introduces concurrency control and recovery concepts to within the transaction.
- Types
  - Closed nesting
    - Subtransactions begin after their parents and finish before them.
    - Commitment of a subtransaction is conditional upon the commitment of the parent (commitment through the root).
  - Open nesting
    - Subtransactions can execute and commit independently.
    - Compensation may be necessary.
Workflows

- “A collection of tasks organized to accomplish some business process.” [D. Georgakopoulos]

- Types
  - Human-oriented workflows
    - Involve humans in performing the tasks.
    - System support for collaboration and coordination; but no system-wide consistency definition
  - System-oriented workflows
    - Computation-intensive & specialized tasks that can be executed by a computer
    - System support for concurrency control and recovery, automatic task execution, notification, etc.
  - Transactional workflows
    - In between the previous two; may involve humans, require access to heterogeneous, autonomous and/or distributed systems, and support selective use of ACID properties

Workflow Example

T₁: Customer request obtained
T₂: Airline reservation performed
T₃: Hotel reservation performed
T₄: Auto reservation performed
T₅: Bill generated
Transactions Provide...

- Atomic and reliable execution in the presence of failures
- Correct execution in the presence of multiple user accesses
- Correct management of replicas (if they support it)

Transaction Processing Issues

- Transaction structure (usually called transaction model)
  - Flat (simple), nested
- Internal database consistency
  - Semantic data control (integrity enforcement) algorithms
- Reliability protocols
  - Atomicity & Durability
  - Local recovery protocols
  - Global commit protocols
Transaction Processing Issues

- Concurrency control algorithms
  - How to synchronize concurrent transaction executions (correctness criterion)
  - Intra-transaction consistency, Isolation
- Replica control protocols
  - How to control the mutual consistency of replicated data
  - One copy equivalence and ROWA

Architecture Revisited
Centralized Transaction Execution

```
User Application ... User Application

Begin_Transaction, Read, Write, Abort, EOT

Transaction Manager (TM)

Read, Write, Abort, EOT

Scheduler (SC)

Scheduled Operations

Recovery Manager (RM)

Results

User Application

Results & User Notifications
```

Distributed Transaction Execution

```
User application

Begin_transaction, Read, Write, EOT, Abort

Results & User notifications

TM

Read, Write, EOT, Abort

SC

RM

Distributed Transaction Execution Model

Replica Control Protocol

Distributed Concurrency Control Protocol

Local Recovery Protocol
```

Page 16
Concurrency Control

- The problem of synchronizing concurrent transactions such that the consistency of the database is maintained while, at the same time, maximum degree of concurrency is achieved.

- Anomalies:
  - Lost updates
    - The effects of some transactions are not reflected on the database.
  - Inconsistent retrievals
    - A transaction, if it reads the same data item more than once, should always read the same value.

Execution Schedule (or History)

- An order in which the operations of a set of transactions are executed.
- A schedule (history) can be defined as a partial order over the operations of a set of transactions.

\[
T_1: \text{Read}(x) \quad T_2: \text{Write}(x) \quad T_3: \text{Read}(x) \\
\quad \text{Write}(x) \quad \quad \text{Write}(y) \quad \text{Read}(y) \\
\quad \text{Commit} \quad \text{Read}(z) \quad \text{Read}(z) \\
\quad \quad \text{Commit} \quad \text{Commit}
\]

\[
H_1 = \{W_2(x), R_1(x), R_2(x), W_3(x), C_1, W_2(y), R_3(y), R_2(z), C_2, R_3(z), C_3\}
\]
**Formalization of Schedule**

A complete schedule $SC(T)$ over a set of transactions $T = \{T_1, \ldots, T_n\}$ is a partial order $SC(T) = \{\sum_{T_i} < T\}$ where

1. $\sum_T = \cup_i \sum_i$, for $i = 1, 2, \ldots, n$
2. $< T \supseteq \cup_i < i'$, for $i = 1, 2, \ldots, n$
3. For any two conflicting operations $O_{ij}, O_{kl} \in \sum_T$, either $O_{ij} < T O_{kl}$ or $O_{kl} < T O_{ij}$

**Complete Schedule – Example**

Given three transactions

<table>
<thead>
<tr>
<th>$T_1$</th>
<th>$T_2$</th>
<th>$T_3$</th>
</tr>
</thead>
<tbody>
<tr>
<td>Read(x)</td>
<td>Write(x)</td>
<td>Read(x)</td>
</tr>
<tr>
<td>Write(x)</td>
<td>Write(y)</td>
<td>Read(y)</td>
</tr>
<tr>
<td>Commit</td>
<td>Read(z)</td>
<td>Read(z)</td>
</tr>
<tr>
<td>Commit</td>
<td>Commit</td>
<td>Commit</td>
</tr>
</tbody>
</table>

A possible complete schedule is given as the DAG

```
  R_1(x) -> W_2(x) -> R_3(x) -> C_3
  |     |     |         |
  W_1(x) -> W_3(y) -> R_5(y) -> C_2
  |     |     |         |
  C_1 -> R_6(z) -> R_5(z) -> C_3
```
Schedule Definition

A schedule is a prefix of a complete schedule such that only some of the operations and only some of the ordering relationships are included.

\[
\begin{align*}
T_1: & \text{ Read}(x) & T_2: & \text{ Write}(x) & T_3: & \text{ Read}(x) \\
& \text{ Write}(x) & & \text{ Write}(y) & & \text{ Read}(y) \\
& \text{ Commit} & & \text{ Read}(z) & & \text{ Read}(z) \\
& & \text{ Commit} & & \text{ Commit} \\
\end{align*}
\]

Serial History

- All the actions of a transaction occur consecutively.
- No interleaving of transaction operations.
- If each transaction is consistent (obeys integrity rules), then the database is guaranteed to be consistent at the end of executing a serial history.

\[
\begin{align*}
T_1: & \text{ Read}(x) & T_2: & \text{ Write}(x) & T_3: & \text{ Read}(x) \\
& \text{ Write}(x) & & \text{ Write}(y) & & \text{ Read}(y) \\
& \text{ Commit} & & \text{ Read}(z) & & \text{ Read}(z) \\
& & \text{ Commit} & & \text{ Commit} \\
\end{align*}
\]

\[
H = \{ W_2(x), W_2(y), R_2(z), C_2, R_1(x), W_1(x), C_1, R_3(x), R_3(y), R_3(z), C_3 \} \]
Serializable History

Transactions execute concurrently, but the net effect of the resulting history upon the database is equivalent to some serial history.

Equivalent with respect to what?

- **Conflict equivalence**: the relative order of execution of the conflicting operations belonging to unaborted transactions in two histories are the same.
- **Conflicting operations**: two incompatible operations (e.g., Read and Write) conflict if they both access the same data item.
  - Incompatible operations of each transaction is assumed to conflict; do not change their execution orders.
  - If two operations from two different transactions conflict, the corresponding transactions are also said to conflict.

The following are not conflict equivalent

\[
H_s = \{(W_2(x), W_2(y), R_2(x), C_2, R_1(x), W_1(x), C_1, R_3(x), R_3(y), R_3(z), C_3) \}
\]

\[
H_1 = \{(W_2(x), R_1(x), R_3(x), W_1(x), C_1, W_3(y), R_2(y), R_3(x), C_2, R_3(z), C_3) \}
\]

The following are conflict equivalent; therefore \(H_2\) is serializable.

\[
H_s = \{(W_2(x), W_2(y), R_2(x), C_2, R_1(x), W_1(x), C_1, R_3(x), R_3(y), R_3(z), C_3) \}
\]

\[
H_2 = \{(W_2(x), R_1(x), W_1(x), C_1, R_3(x), W_3(y), R_2(y), R_3(z), C_2, R_3(z), C_3) \}
\]
Serializability in Distributed DBMS

- Somewhat more involved. Two histories have to be considered:
  - local histories
  - global history

- For global transactions (i.e., global history) to be serializable, two conditions are necessary:
  - Each local history should be serializable.
  - Two conflicting operations should be in the same relative order in all of the local histories where they appear together.

Global Non-serializability

\[ T_1: \text{Read}(x) \quad T_2: \text{Read}(x) \]
\[ x \leftarrow x+5 \quad x \leftarrow x \times 15 \]
Write(x) \quad Write(x)
Commit \quad Commit

The following two local histories are individually serializable (in fact serial), but the two transactions are not globally serializable.

\[ LH_1 = \{R_1(x), W_1(x), C_1, R_2(x), W_2(x), C_2\} \]
\[ LH_2 = \{R_2(x), W_2(x), C_2, R_1(x), W_1(x), C_1\} \]
Concurrency Control Algorithms

- Pessimistic
  - Two-Phase Locking-based (2PL)
    - Centralized (primary site) 2PL
    - Primary copy 2PL
    - Distributed 2PL
  - Timestamp Ordering (TO)
    - Basic TO
    - Multiversion TO
    - Conservative TO
- Hybrid

- Optimistic
  - Locking-based
  - Timestamp ordering-based

Locking-Based Algorithms

- Transactions indicate their intentions by requesting locks from the scheduler (called lock manager).
- Locks are either read lock (rl) [also called shared lock] or write lock (wl) [also called exclusive lock]
- Read locks and write locks conflict (because Read and Write operations are incompatible)
  
<p>| | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>rl</td>
<td>wl</td>
</tr>
<tr>
<td>rl</td>
<td>yes</td>
</tr>
<tr>
<td>wl</td>
<td>no</td>
</tr>
</tbody>
</table>
- Locking works nicely to allow concurrent processing of transactions.
Two-Phase Locking (2PL)

1. A Transaction locks an object before using it.
2. When an object is locked by another transaction, the requesting transaction must wait.
3. When a transaction releases a lock, it may not request another lock.

Strict 2PL

Hold locks until the end.
Centralized 2PL

- There is only one 2PL scheduler in the distributed system.
- Lock requests are issued to the central scheduler.

Distributed 2PL

- 2PL schedulers are placed at each site. Each scheduler handles lock requests for data at that site.
- A transaction may read any of the replicated copies of item $x$, by obtaining a read lock on one of the copies of $x$. Writing into $x$ requires obtaining write locks for all copies of $x$. 

Distributed DBMS © 1998 M. Tamer Özsu & Patrick Valduriez
Distributed 2PL Execution

<table>
<thead>
<tr>
<th>Coordinating TM</th>
<th>Participating LMs</th>
<th>Participating DPs</th>
</tr>
</thead>
<tbody>
<tr>
<td><img src="image1" alt="Diagram" /></td>
<td><img src="image2" alt="Diagram" /></td>
<td><img src="image3" alt="Diagram" /></td>
</tr>
</tbody>
</table>

Timestamp Ordering

1. Transaction \( T_i \) is assigned a globally unique timestamp \( ts(T_i) \).
2. Transaction manager attaches the timestamp to all operations issued by the transaction.
3. Each data item is assigned a write timestamp \( wts(x) \) and a read timestamp \( rts(x) \):
   - \( rts(x) = \) largest timestamp of any read on \( x \)
   - \( wts(x) = \) largest timestamp of any read on \( x \)
4. Conflicting operations are resolved by timestamp order.

Basic T/O:

<table>
<thead>
<tr>
<th>Operation</th>
<th>Condition</th>
<th>Action</th>
</tr>
</thead>
<tbody>
<tr>
<td>( R_i(x) )</td>
<td>if ( ts(T_i) &lt; wts(x) )</td>
<td>reject ( R_i(x) )</td>
</tr>
<tr>
<td></td>
<td>if ( ts(T_i) &lt; rts(x) ) and ( ts(T_i) &lt; wts(x) )</td>
<td>reject ( W_i(x) )</td>
</tr>
<tr>
<td>( W_i(x) )</td>
<td>then</td>
<td>reject ( W_i(x) )</td>
</tr>
<tr>
<td></td>
<td>else</td>
<td>accept ( W_i(x) )</td>
</tr>
<tr>
<td></td>
<td>( rts(x) \leftarrow ts(T_i) )</td>
<td>( wts(x) \leftarrow ts(T_i) )</td>
</tr>
</tbody>
</table>
**Conservative Timestamp Ordering**

- Basic timestamp ordering tries to execute an operation as soon as it receives it
  - progressive
  - too many restarts since there is no delaying
- Conservative timestamping delays each operation until there is an assurance that it will not be restarted
- Assurance?
  - No other operation with a smaller timestamp can arrive at the scheduler
  - Note that the delay may result in the formation of deadlocks

**Multiversion Timestamp Ordering**

- Do not modify the values in the database, create new values.
- A \( R_i(\alpha) \) is translated into a read on one version of \( x \).
  - Find a version of \( x \) (say \( x_r \)) such that \( ts(x_r) \) is the largest timestamp less than \( ts(T_i) \).
- A \( W_j(\alpha) \) is translated into \( W_j(x_w) \) and accepted if the scheduler has not yet processed any \( R_j(\alpha) \) such that

\[
\text{ts}(T_i) < \text{ts}(x_r) < \text{ts}(T_j)
\]
Optimistic Concurrency Control Algorithms

Pessimistic execution

- Validate
- Read
- Compute
- Write

Optimistic execution

- Read
- Compute
- Validate
- Write

Transaction execution model: divide into subtransactions each of which execute at a site

- \( T_{ij} \): transaction \( T_i \) that executes at site \( j \)

- Transactions run independently at each site until they reach the end of their read phases

- All subtransactions are assigned a timestamp at the end of their read phase

- Validation test performed during validation phase. If one fails, all rejected.
### Optimistic CC Validation Test

1. If all transactions $T_k$ where $ts(T_k) < ts(T_{ij})$ have completed their write phase before $T_{ij}$ has started its read phase, then validation succeeds.

   ➔ Transaction executions in serial order

   ![Diagram](image1)

   ![Diagram](image2)

2. If there is any transaction $T_k$ such that $ts(T_k) < ts(T_{ij})$ and which completes its write phase while $T_{ij}$ is in its read phase, then validation succeeds if $WS(T_k) \cap RS(T_{ij}) = \emptyset$.

   ➔ Read and write phases overlap, but $T_{ij}$ does not read data items written by $T_k$

   ![Diagram](image1)

   ![Diagram](image2)
Optimistic CC Validation Test

If there is any transaction \( T_k \) such that \( ts(T_k) < ts(T_{ij}) \) and which completes its read phase before \( T_{ij} \) completes its read phase, then validation succeeds if \( WS(T_k) \cap RS(T_{ij}) = \emptyset \) and \( WS(T_k) \cap WS(T_{ij}) = \emptyset \)

- They overlap, but don't access any common data items.

\[
\begin{array}{c|c|c|c|}
T_k & R & V & W \\
\hline
T_{ij} & R & V & W \\
\end{array}
\]

Deadlock

- A transaction is deadlocked if it is blocked and will remain blocked until there is intervention.
- Locking-based CC algorithms may cause deadlocks.
- TO-based algorithms that involve waiting may cause deadlocks.
- Wait-for graph
  - If transaction \( T_i \) waits for another transaction \( T_j \) to release a lock on an entity, then \( T_i \rightarrow T_j \) in WFG.
Local versus Global WFG

Assume \( T_1 \) and \( T_2 \) run at site 1, \( T_3 \) and \( T_4 \) run at site 2. Also assume \( T_3 \) waits for a lock held by \( T_1 \) which waits for a lock held by \( T_2 \) which waits for a lock held by \( T_1 \) which waits for a lock held by \( T_3 \).

Local WFG

\[
\begin{array}{c|c}
\text{Site 1} & \text{Site 2} \\
\hline
T_1 & T_4 \\
T_2 & T_3 \\
\end{array}
\]

Global WFG

\[
\begin{array}{c|c}
\text{Site 1} & \text{Site 2} \\
\hline
T_1 & T_4 \\
T_2 & T_3 \\
\end{array}
\]

Deadlock Management

- Ignore
  - Let the application programmer deal with it, or restart the system
- Prevention
  - Guaranteeing that deadlocks can never occur in the first place. Check transaction when it is initiated. Requires no run time support.
- Avoidance
  - Detecting potential deadlocks in advance and taking action to insure that deadlock will not occur. Requires run time support.
- Detection and Recovery
  - Allowing deadlocks to form and then finding and breaking them. As in the avoidance scheme, this requires run time support.
Deadlock Prevention

- All resources which may be needed by a transaction must be predeclared.
  - The system must guarantee that none of the resources will be needed by an ongoing transaction.
  - Resources must only be reserved, but not necessarily allocated a priori
  - Unsuitability of the scheme in database environment
  - Suitable for systems that have no provisions for undoing processes.

Evaluation:
- Reduced concurrency due to preallocation
- Evaluating whether an allocation is safe leads to added overhead.
- Difficult to determine (partial order)
+ No transaction rollback or restart is involved.

Deadlock Avoidance

- Transactions are not required to request resources a priori.
- Transactions are allowed to proceed unless a requested resource is unavailable.
- In case of conflict, transactions may be allowed to wait for a fixed time interval.
- Order either the data items or the sites and always request locks in that order.
- More attractive than prevention in a database environment.
Deadlock Avoidance –
Wait-Die & Wound-Wait Algorithms

WAIT-DIE Rule: If $T_i$ requests a lock on a data item which is already locked by $T_j$, then $T_i$ is permitted to wait iff $ts(T_i) < ts(T_j)$. If $ts(T_i) > ts(T_j)$, then $T_i$ is aborted and restarted with the same timestamp.
- if $ts(T_i) < ts(T_j)$ then $T_i$ waits else $T_i$ dies
- non-preemptive: $T_i$ never preempts $T_j$
- prefers younger transactions

WOUND-WAIT Rule: If $T_i$ requests a lock on a data item which is already locked by $T_j$, then $T_i$ is permitted to wait iff $ts(T_i) > ts(T_j)$. If $ts(T_i) < ts(T_j)$, then $T_j$ is aborted and the lock is granted to $T_i$.
- if $ts(T_i) < ts(T_j)$ then $T_j$ is wounded else $T_i$ waits
- preemptive: $T_i$ preempts $T_j$ if it is younger
- prefers older transactions

Deadlock Detection

- Transactions are allowed to wait freely.
- Wait-for graphs and cycles.
- Topologies for deadlock detection algorithms
  - Centralized
  - Distributed
  - Hierarchical
**Centralized Deadlock Detection**

- One site is designated as the deadlock detector for the system. Each scheduler periodically sends its local WFG to the central site which merges them to a global WFG to determine cycles.
- How often to transmit?
  - Too often ⇒ higher communication cost but lower delays due to undetected deadlocks
  - Too late ⇒ higher delays due to deadlocks, but lower communication cost
- Would be a reasonable choice if the concurrency control algorithm is also centralized.
- Proposed for Distributed INGRES

**Hierarchical Deadlock Detection**

Build a hierarchy of detectors

```
         DD_{cx}
        /   \  
DD_{11}  DD_{14} 
|     / \       |
Site 1  Site 2 Site 3 Site 4 
DD_{21} DD_{22} DD_{23} DD_{24}
```
Distributed Deadlock Detection

- Sites cooperate in detection of deadlocks.
- One example:
  - The local WFGs are formed at each site and passed on to other sites. Each local WFG is modified as follows:
    - Since each site receives the potential deadlock cycles from other sites, these edges are added to the local WFGs
    - The edges in the local WFG which show that local transactions are waiting for transactions at other sites are joined with edges in the local WFGs which show that remote transactions are waiting for local ones.
  - Each local deadlock detector:
    - looks for a cycle that does not involve the external edge. If it exists, there is a local deadlock which can be handled locally.
    - looks for a cycle involving the external edge. If it exists, it indicates a potential global deadlock. Pass on the information to the next site.

Reliability

Problem:

How to maintain
atomicity
durability
properties of transactions
Fundamental Definitions

- **Reliability**
  - A measure of success with which a system conforms to some authoritative specification of its behavior.
  - Probability that the system has not experienced any failures within a given time period.
  - Typically used to describe systems that cannot be repaired or where the continuous operation of the system is critical.

- **Availability**
  - The fraction of the time that a system meets its specification.
  - The probability that the system is operational at a given time $t$.

Basic System Concepts

[Diagram showing ENVIRONMENT, SYSTEM, Component 1, Component 2, Component 3, Stimuli, Responses, External state, Internal state]
Fundamental Definitions

- **Failure**
  - The deviation of a system from the behavior that is described in its specification.

- **Erroneous state**
  - The internal state of a system such that there exist circumstances in which further processing, by the normal algorithms of the system, will lead to a failure which is not attributed to a subsequent fault.

- **Error**
  - The part of the state which is incorrect.

- **Fault**
  - An error in the internal states of the components of a system or in the design of a system.

Faults to Failures

Fault \(\rightarrow\) Error \(\rightarrow\) Failure
Types of Faults

- Hard faults
  - Permanent
  - Resulting failures are called hard failures

- Soft faults
  - Transient or intermittent
  - Account for more than 90% of all failures
  - Resulting failures are called soft failures

Fault Classification
Failures

Fault occurs
Error caused
Detection of error
Repair
Fault occurs
Error caused

Multiple errors can occur during this period

Fault Tolerance Measures

Reliability

\[ R(t) = \Pr\{0 \text{ failures in time } [0,t] \mid \text{ no failures at } t=0\} \]

If occurrence of failures is Poisson

\[ R(t) = \Pr\{0 \text{ failures in time } [0,t]\} \]

Then

\[ \Pr(k \text{ failures in time } [0,t]) = \frac{e^{-m(t)}[m(t)]^k}{k!} \]

where \( m(t) \) is known as the hazard function which gives the time-dependent failure rate of the component and is defined as

\[ m(t) = \int_0^t z(x)dx \]
Fault-Tolerance Measures

Reliability

The mean number of failures in time \([0, t]\) can be computed as

\[
E[k] = \sum_{k=0}^{\infty} k \frac{e^{-m(t)} (m(t))^k}{k!} = m(t)
\]

and the variance can be computed as

\[
Var[k] = E[k^2] - (E[k])^2 = m(t)
\]

Thus, reliability of a single component is

\[
R(t) = e^{-m(t)}
\]

and of a system consisting of \(n\) non-redundant components as

\[
R_{sys}(t) = \prod_{i=1}^{n} R_i(t)
\]

Fault-Tolerance Measures

Availability

\[
A(t) = \text{Pr}\{\text{system is operational at time } t\}
\]

Assume

- Poisson failures with rate \(\lambda\)
- Repair time is exponentially distributed with mean \(1/\mu\)

Then, steady-state availability

\[
A = \lim_{t \to \infty} A(t) = \frac{\mu}{\lambda + \mu}
\]
Fault-Tolerance Measures

**MTBF**
Mean time between failures

\[ \text{MTBF} = \int_{0}^{\infty} R(t) dt \]

**MTTR**
Mean time to repair

**Availability**

\[ \frac{\text{MTBF}}{\text{MTBF} + \text{MTTR}} \]

Sources of Failure – SLAC Data (1985)

Sources of Failure – Japanese Data (1986)


Sources of Failure – 5ESS Switch (1987)

Sources of Failures – Tandem Data (1985)


Types of Failures

- **Transaction failures**
  - Transaction aborts (unilaterally or due to deadlock)
  - Avg. 3% of transactions abort abnormally
- **System (site) failures**
  - Failure of processor, main memory, power supply, ...
  - Main memory contents are lost, but secondary storage contents are safe
  - Partial vs. total failure
- **Media failures**
  - Failure of secondary storage devices such that the stored data is lost
  - Head crash/controller failure (?)
- **Communication failures**
  - Lost/undeliverable messages
  - Network partitioning
Local Recovery Management – Architecture

- **Volatile storage**
  - Consists of the main memory of the computer system (RAM).

- **Stable storage**
  - Resilient to failures and loses its contents only in the presence of media failures (e.g., head crashes on disks).
  - Implemented via a combination of hardware (non-volatile storage) and software (stable-write, stable-read, clean-up) components.

Update Strategies

- **In-place update**
  - Each update causes a change in one or more data values on pages in the database buffers

- **Out-of-place update**
  - Each update causes the new value(s) of data item(s) to be stored separate from the old value(s)
In-Place Update Recovery
Information

Database Log
Every action of a transaction must not only perform
the action, but must also write a log record to an
append-only file.

Logging
The log contains information used by the
recovery process to restore the consistency of a
system. This information may include

- transaction identifier
- type of operation (action)
- items accessed by the transaction to perform the
  action
- old value (state) of item (before image)
- new value (state) of item (after image)
...
Why Logging?

Upon recovery:
- all of $T_1$'s effects should be reflected in the database (REDO if necessary due to a failure)
- none of $T_2$'s effects should be reflected in the database (UNDO if necessary)

REDO Protocol

- REDO'ing an action means performing it again.
- The REDO operation uses the log information and performs the action that might have been done before, or not done due to failures.
- The REDO operation generates the new image.
UNDO Protocol

- UNDO'ing an action means to restore the object to its before image.
- The UNDO operation uses the log information and restores the old value of the object.

When to Write Log Records Into Stable Store

Assume a transaction $T$ updates a page $P$

- **Fortunate case**
  - System writes $P$ in stable database
  - System updates stable log for this update
  - SYSTEM FAILURE OCCURS!... (before $T$ commits)

  We can recover (undo) by restoring $P$ to its old state by using the log

- **Unfortunate case**
  - System writes $P$ in stable database
  - SYSTEM FAILURE OCCURS!... (before stable log is updated)

  We cannot recover from this failure because there is no log record to restore the old value.

- **Solution:** Write-Ahead Log (WAL) protocol
Write–Ahead Log Protocol

- **Notice:**
  - If a system crashes before a transaction is committed, then all the operations must be undone. Only need the before images *(undo portion of the log).*
  - Once a transaction is committed, some of its actions might have to be redone. Need the after images *(redo portion of the log).*

- **WAL protocol:**
  1. Before a stable database is updated, the undo portion of the log should be written to the stable log
  2. When a transaction commits, the redo portion of the log must be written to stable log prior to the updating of the stable database.

Logging Interface
Out-of-Place Update
Recovery Information

- Shadowing
  - When an update occurs, don’t change the old page, but create a shadow page with the new values and write it into the stable database.
  - Update the access paths so that subsequent accesses are to the new shadow page.
  - The old page retained for recovery.

- Differential files
  - For each file F maintain
    - a read only part FR
    - a differential file consisting of insertions part DF+ and deletions part DF-
    - Thus, \( F = (FR \cup DF^+) – DF^- \)
  - Updates treated as delete old value, insert new value

Execution of Commands

Commands to consider:

```
begin_transaction
read
write
commit
abort
recover
```

Independent of execution strategy for LRM
Execution Strategies

- **Dependent upon**
  - Can the buffer manager decide to write some of the buffer pages being accessed by a transaction into stable storage or does it wait for LRM to instruct it?
    - fix/no-fix decision
  - Does the LRM force the buffer manager to write certain buffer pages into stable database at the end of a transaction's execution?
    - flush/no-flush decision

- **Possible execution strategies:**
  - no-fix/no-flush
  - no-fix/flush
  - fix/no-flush
  - fix/flush

---

No-Fix/No-Flush

- **Abort**
  - Buffer manager may have written some of the updated pages into stable database
  - LRM performs transaction undo (or partial undo)

- **Commit**
  - LRM writes an “end_of_transaction” record into the log.

- **Recover**
  - For those transactions that have both a “begin_transaction” and an “end_of_transaction” record in the log, a partial redo is initiated by LRM
  - For those transactions that only have a “begin_transaction” in the log, a global undo is executed by LRM
No-Fix/Flush

- **Abort**
  - Buffer manager may have written some of the updated pages into stable database
  - LRM performs transaction undo (or partial undo)
- **Commit**
  - LRM issues a `flush` command to the buffer manager for all updated pages
  - LRM writes an “end_of_transaction” record into the log.
- **Recover**
  - No need to perform redo
  - Perform global undo

Fix/No-Flush

- **Abort**
  - None of the updated pages have been written into stable database
  - Release the fixed pages
- **Commit**
  - LRM writes an “end_of_transaction” record into the log.
  - LRM sends an `unfix` command to the buffer manager for all pages that were previously fixed
- **Recover**
  - Perform partial redo
  - No need to perform global undo
Fix/Flush

- **Abort**
  - None of the updated pages have been written into stable database
  - Release the fixed pages

- **Commit (the following have to be done atomically)**
  - LRM issues a flush command to the buffer manager for all updated pages
  - LRM sends an unfixed command to the buffer manager for all pages that were previously fixed
  - LRM writes an “end_of_transaction” record into the log

- **Recover**
  - No need to do anything

---

Checkpoints

- Simplifies the task of determining actions of transactions that need to be undone or redone when a failure occurs.

- A checkpoint record contains a list of active transactions.

- **Steps:**
  - Write a begin_checkpoint record into the log
  - Collect the checkpoint data into the stable storage
  - Write an end_checkpoint record into the log
Media Failures – Full Architecture

- **Secondary storage**
  - **Stable log**
  - **Stable database**

- **Main memory**
  - **Local Recovery Manager**
    - Read
    - Write
    - Fetch, Flush

- **Database Buffer Manager**
  - Read
  - Write
  - Database buffers (Volatile database)
  - Log buffers

- **Archive database**
- **Archive log**

Distributed Reliability Protocols

- **Commit protocols**
  - How to execute commit command for distributed transactions.
  - Issue: how to ensure atomicity and durability?

- **Termination protocols**
  - If a failure occurs, how can the remaining operational sites deal with it.
  - Non-blocking: the occurrence of failures should not force the sites to wait until the failure is repaired to terminate the transaction.

- **Recovery protocols**
  - When a failure occurs, how do the sites where the failure occurred deal with it.
  - Independent: a failed site can determine the outcome of a transaction without having to obtain remote information.

- **Independent recovery ⇒ non-blocking termination**
Two-Phase Commit (2PC)

*Phase 1*: The coordinator gets the participants ready to write the results into the database

*Phase 2*: Everybody writes the results into the database

- **Coordinator**: The process at the site where the transaction originates and which controls the execution
- **Participant**: The process at the other sites that participate in executing the transaction

**Global Commit Rule:**
1. The coordinator aborts a transaction if and only if at least one participant votes to abort it.
2. The coordinator commits a transaction if and only if all of the participants vote to commit it.

Centralized 2PC

```
<p>| | | | | | | | | | | | | | | | | | | |
|   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |</p>
<table>
<thead>
<tr>
<th>ready?</th>
<th>yes/no</th>
<th>commit/abort?</th>
<th>committed/aborted</th>
</tr>
</thead>
<tbody>
<tr>
<td>Phase 1</td>
<td></td>
<td></td>
<td>Phase 2</td>
</tr>
</tbody>
</table>
```
2PC Protocol Actions

<table>
<thead>
<tr>
<th>Coordinator</th>
<th>Participant</th>
</tr>
</thead>
<tbody>
<tr>
<td>INITIAL</td>
<td>INITIAL</td>
</tr>
<tr>
<td>PREPARE</td>
<td>WRITE begin_commit in log</td>
</tr>
<tr>
<td>VOTE-ABORT</td>
<td>WRITE abort in log</td>
</tr>
<tr>
<td>GLOB-ABORT</td>
<td>WRITE commit in log</td>
</tr>
<tr>
<td>READY</td>
<td>WRITE end_of_transaction in log</td>
</tr>
<tr>
<td>ANY?</td>
<td>WRITE abort in log</td>
</tr>
<tr>
<td>NO</td>
<td>WRITE commit in log</td>
</tr>
<tr>
<td>COMMIT</td>
<td>WRITE commit in log</td>
</tr>
<tr>
<td>ABORT</td>
<td>WRITE commit in log</td>
</tr>
<tr>
<td>ACK</td>
<td>WRITE commit in log</td>
</tr>
<tr>
<td>ACK</td>
<td>WRITE commit in log</td>
</tr>
</tbody>
</table>

Linear 2PC

Phase 1

<table>
<thead>
<tr>
<th>Prepare</th>
<th>VC/VA</th>
<th>VC/VA</th>
<th>VC/VA</th>
<th>VC/VA</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>GC/GA</td>
<td>GC/GA</td>
<td>GC/GA</td>
<td>GC/GA</td>
</tr>
<tr>
<td>2</td>
<td>GC/GA</td>
<td>GC/GA</td>
<td>GC/GA</td>
<td>GC/GA</td>
</tr>
<tr>
<td>3</td>
<td>GC/GA</td>
<td>GC/GA</td>
<td>GC/GA</td>
<td>GC/GA</td>
</tr>
<tr>
<td>4</td>
<td>GC/GA</td>
<td>GC/GA</td>
<td>GC/GA</td>
<td>GC/GA</td>
</tr>
<tr>
<td>5</td>
<td>GC/GA</td>
<td>GC/GA</td>
<td>GC/GA</td>
<td>GC/GA</td>
</tr>
<tr>
<td>N</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Phase 2

VC: Vote-Commit, VA: Vote-Abort, GC: Global-commit, GA: Global-abort

Distributed DBMS © 1998 M. Tamer Özsu & Patrick Valduriez

Page 54
Site Failures - 2PC Termination

- Timeout in INITIAL
  - Who cares

- Timeout in WAIT
  - Cannot unilaterally commit
  - Can unilaterally abort

- Timeout in ABORT or COMMIT
  - Stay blocked and wait for the acks

Timeout in INITIAL
- Coordinator must have failed in INITIAL state
- Unilaterally abort

Timeout in READY
- Stay blocked
Site Failures - 2PC Recovery

- **Failure in INITIAL**
  - Unilaterally abort upon recovery

- **Failure in READY**
  - The coordinator has been informed about the local decision
  - Treat as timeout in READY state and invoke the termination protocol

- **Failure in ABORT or COMMIT**
  - Nothing special needs to be done

---

COORDINATOR

- **Failure in INITIAL**
  - Start the commit process upon recovery

- **Failure in WAIT**
  - Restart the commit process upon recovery

- **Failure in ABORT or COMMIT**
  - Nothing special if all the acks have been received
  - Otherwise the termination protocol is involved

---

PARTICIPANTS

- **Failure in INITIAL**
  - Unilaterally abort upon recovery

- **Failure in READY**
  - The coordinator has been informed about the local decision
  - Treat as timeout in READY state and invoke the termination protocol

- **Failure in ABORT or COMMIT**
  - Nothing special needs to be done
2PC Recovery Protocols – Additional Cases

Arise due to non-atomicity of log and message send actions

- Coordinator site fails after writing “begin_commit” log and before sending “prepare” command
  - treat it as a failure in WAIT state; send “prepare” command
- Participant site fails after writing “ready” record in log but before “vote-commit” is sent
  - treat it as failure in READY state
  - alternatively, can send “vote-commit” upon recovery
- Participant site fails after writing “abort” record in log but before “vote-abort” is sent
  - no need to do anything upon recovery

2PC Recovery Protocols – Additional Case

- Coordinator site fails after logging its final decision record but before sending its decision to the participants
  - coordinator treats it as a failure in COMMIT or ABORT state
  - participants treat it as timeout in the READY state
- Participant site fails after writing “abort” or “commit” record in log but before acknowledgement is sent
  - participant treats it as failure in COMMIT or ABORT state
  - coordinator will handle it by timeout in COMMIT or ABORT state
Problem With 2PC

- Blocking
  - Ready implies that the participant waits for the coordinator
  - If coordinator fails, site is blocked until recovery
  - Blocking reduces availability
- Independent recovery is not possible
- However, it is known that:
  - Independent recovery protocols exist only for single site failures; no independent recovery protocol exists which is resilient to multiple-site failures.
- So we search for these protocols – 3PC

Three-Phase Commit

- 3PC is non-blocking.
- A commit protocols is non-blocking iff
  - it is synchronous within one state transition, and
  - its state transition diagram contains
    - no state which is “adjacent” to both a commit and an abort state, and
    - no non-committable state which is “adjacent” to a commit state
- Adjacent: possible to go from one stat to another with a single state transition
- Committable: all sites have voted to commit a transaction
  - e.g.: COMMIT state
State Transitions in 3PC

Coordinator

INITIAL
Commit command
Prepare

WAIT
Vote-abort
Global-abort
Vote-commit
Prepare-to-commit

ABORT
Ready-to-commit
Global commit

PRE-COMMIT

COMMIT

Participants

INITIAL
Prepare
Vote-commit

READY
Global-abort
Ack
Prepared-to-commit

ABORT
Global commit
Ack

PRE-COMMIT

COMMIT

Communication Structure

Phase 1
C

P

C

P

P

P

P

C

C

P

P

P

P

Phase 2

ready? yes/no
pre-commit/pre-abort? yes/no
commit/abort
ack

Phase 3
### Site Failures – 3PC Termination

#### Coordinator

- **INITIAL**
  - Commit command
  - Prepare

- **WAIT**
  - Vote-abort
  - Global-abort
  - Vote-commit
  - Prepare-to-commit

- **ABORT**
  - Ready-to-commit
  - Global commit

- **PRE-COMMIT**
  - Vote-commit
  - Prepare-to-commit

- **COMMIT**

- **Timeout in INITIAL**
  - Who cares

- **Timeout in WAIT**
  - Unilaterally abort

- **Timeout in PRE-COMMIT**
  - Participants may not be in PRE-COMMIT, but at least in READY
  - Move all the participants to PRECOMMIT state
  - Terminate by globally committing

---

### Site Failures – 3PC Termination

#### Coordinator

- **INITIAL**
  - Commit command
  - Prepare

- **WAIT**
  - Vote-abort
  - Global-abort
  - Vote-commit
  - Prepare-to-commit

- **ABORT**
  - Ready-to-commit
  - Global commit

- **PRE-COMMIT**
  - Vote-commit
  - Prepare-to-commit

- **COMMIT**

- **Timeout in ABORT or COMMIT**
  - Just ignore and treat the transaction as completed
  - Participants are either in PRECOMMIT or READY state and can follow their termination protocols
Site Failures – 3PC Termination

- **Timeout in INITIAL**
  - Coordinator must have failed in INITIAL state
  - Unilaterally abort
- **Timeout in READY**
  - Voted to commit, but does not know the coordinator's decision
  - Elect a new coordinator and terminate using a special protocol
- **Timeout in PRECOMMIT**
  - Handle it the same as timeout in READY state

Termination Protocol Upon Coordinator Election

New coordinator can be in one of four states: WAIT, PRECOMMIT, COMMIT, ABORT

1. Coordinator sends its state to all of the participants asking them to assume its state.
2. Participants “back-up” and reply with appropriate messages, except those in ABORT and COMMIT states. Those in these states respond with “Ack” but stay in their states.
3. Coordinator guides the participants towards termination:
   - If the new coordinator is in the WAIT state, participants can be in INITIAL, READY, ABORT or PRECOMMIT states. New coordinator globally aborts the transaction.
   - If the new coordinator is in the PRECOMMIT state, the participants can be in READY, PRECOMMIT or COMMIT states. The new coordinator will globally commit the transaction.
   - If the new coordinator is in the ABORT or COMMIT states, at the end of the first phase, the participants will have moved to that state as well.
Failure in INITIAL
- start commit process upon recovery

Failure in WAIT
- the participants may have elected a new coordinator and terminated the transaction
- the new coordinator could be in WAIT or ABORT states ⇒ transaction aborted
- ask around for the fate of the transaction

Failure in PRECOMMIT
- ask around for the fate of the transaction

Failure in COMMIT or ABORT
- Nothing special if all the acknowledgements have been received; otherwise the termination protocol is involved
Site Failures – 3PC Recovery

- Failure in INITIAL
  - unilaterally abort upon recovery
- Failure in READY
  - the coordinator has been informed about the local decision
  - upon recovery, ask around
- Failure in PRECOMMIT
  - ask around to determine how the other participants have terminated the transaction
- Failure in COMMIT or ABORT
  - no need to do anything

Network Partitioning

- Simple partitioning
  - Only two partitions
- Multiple partitioning
  - More than two partitions
- Formal bounds (due to Skeen):
  - There exists no non-blocking protocol that is resilient to a network partition if messages are lost when partition occurs.
  - There exist non-blocking protocols which are resilient to a single network partition if all undeliverable messages are returned to sender.
  - There exists no non-blocking protocol which is resilient to a multiple partition.
Independent Recovery Protocols for Network Partitioning

- No general solution possible
  - allow one group to terminate while the other is blocked
  - improve availability
- How to determine which group to proceed?
  - The group with a majority
- How does a group know if it has majority?
  - centralized
    - whichever partitions contains the central site should terminate the transaction
  - voting-based (quorum)
    - different for replicated vs non-replicated databases

Quorum Protocols for Non-Replicated Databases

- The network partitioning problem is handled by the commit protocol.
- Every site is assigned a vote \( V_i \).
- Total number of votes in the system \( V \)
- Abort quorum \( V_a \), commit quorum \( V_c \)
  - \( V_a + V_c > V \) where \( 0 \leq V_a, V_c \leq V \)
  - Before a transaction commits, it must obtain a commit quorum \( V_c \)
  - Before a transaction aborts, it must obtain an abort quorum \( V_a \)
State Transitions in Quorum Protocols

Quorum Protocols for Replicated Databases

- Network partitioning is handled by the replica control protocol.
- One implementation:
  - Assign a vote to each copy of a replicated data item (say $V_i$) such that $\sum_i V_i = V$
  - Each operation has to obtain a read quorum ($V_r$) to read and a write quorum ($V_w$) to write a data item
  - Then the following rules have to be obeyed in determining the quorums:
    - $V_r + V_w > V$  a data item is not read and written by two transactions concurrently
    - $V_w > V/2$  two write operations from two transactions cannot occur concurrently on the same data item
Use for Network Partitioning

- Simple modification of the ROWA rule:
  - When the replica control protocol attempts to read or write a data item, it first checks if a majority of the sites are in the same partition as the site that the protocol is running on (by checking its votes). If so, execute the ROWA rule within that partition.

- Assumes that failures are “clean” which means:
  - failures that change the network’s topology are detected by all sites instantaneously
  - each site has a view of the network consisting of all the sites it can communicate with

Open Problems

- Replication protocols
  - experimental validation
  - replication of computation and communication

- Transaction models
  - changing requirements
    - cooperative sharing vs. competitive sharing
    - interactive transactions
    - longer duration
    - complex operations on complex data
  - relaxed semantics
    - non-serializable correctness criteria
Transaction Model Design Space

<table>
<thead>
<tr>
<th>Object complexity</th>
<th>Transaction structure</th>
</tr>
</thead>
<tbody>
<tr>
<td>active objects</td>
<td>○</td>
</tr>
<tr>
<td>ADT + complex objects</td>
<td>○</td>
</tr>
<tr>
<td>ADT instances</td>
<td>○</td>
</tr>
<tr>
<td>simple data</td>
<td>○</td>
</tr>
</tbody>
</table>

- flat
- closed nesting
- open nesting
- mixed